# Homework 6: FDs, NFs and XML (due April 13<sup>\*</sup>, 2016, 4:00pm, hard-copy in-class please)

#### **Reminders**:

- a. Out of 100 points. Contains 5 pages.
- b. Rough time-estimates: ~5-7 hours.
- c. Please type your answers. Illegible handwriting may get no points, at the discretion of the grader. Only drawings may be hand-drawn, as long as they are neat and legible.
- d. There could be more than one correct answer. We shall accept them all.
- e. Whenever you are making an assumption, please state is clearly.
- f. **Important:** Unless otherwise mentioned, assume that you need to show your work e.g. if the question asks 'what are R's keys?' we do not just want a list; we want a step-by-step explanation as well.
- g. Lead TA for this homework: Sorour Amiri.

# Q1. XML [7 points]

(Courtesy Widom and Ullman) This question is based on the XML "Vehicles" document below.

```
<Vehicles>
    <Car manf="Hyundai">
        <Model>Azera</Model>
        <HorsePower>240</HorsePower>
    </Car>
    <Car manf="Toyota">
        <Model>Camry</Model>
        <HorsePower>240</HorsePower>
    </Car>
    <Truck manf="Toyota">
        <Model>Tundra</Model>
        <HorsePower>240</HorsePower>
    </Truck>
    <Car manf="Hyundai">
        <Model>Elantra</Model>
        <HorsePower>120</HorsePower>
    </Car>
    <Car manf="Toyota">
```

```
<Model>Prius</Model>
<HorsePower>120</HorsePower>
</Car>
</Vehicles>
```

- Q1.1 (3 points) Which of the following XPath expressions, when applied to the "Vehicles" document, does NOT produce a sequence of exactly three items?
  - a) /Vehicles/Car[@manf="Toyota"]/HorsePower
  - b) /Vehicles/\*[HorsePower>200]/Model
  - c) / /\*[@manf="Toyota"]/@manf

Justify your answer briefly (i.e. give the output of the query (queries) which do not produce a sequence of three items).

Q1.2 (4 points) In a DTD that the "Vehicles" document satisfies, which of the following element declarations would you definitely NOT find?

a) <!ELEMENT Vehicles (Car\*, Truck+, Car\*)>
b) <!ELEMENT Vehicles (Car+, Truck\*, Car)>
c) <!ELEMENT Vehicles ((Car | Truck)\*)>
d) <!ELEMENT Vehicles (Car\*, Truck\*, Car\*, Truck\*, Car\*)>

Again, just justify your choice(s) briefly (i.e. explain why you would not find your selected element declaration(s) in a DTD for Vehicles).

#### Q2. FDs Definition [8 points]

Suppose that we have the following four tuples in an *instance* of relation schema R with five attributes ABCDE:

Α	B	C	D	Ε
1	2	4	3	8
1	2	5	7	9
3	6	4	7	0
1	6	5	7	9

Q2.1 (4 points) Which of the following dependencies can you infer does not hold over R? Just list the ones that do not hold, giving the reason in 1 line for each.

(a) 
$$E \rightarrow D$$
, (b)  $D \rightarrow E$ , (c)  $DE \rightarrow B$ , (d)  $E \rightarrow AB$ , (e)  $BC \rightarrow ED$ 

- Q2.2 (4 points) If ABE  $\rightarrow$  ACE and ABE  $\rightarrow$  ADE hold in R, which of the following dependencies can be deduced? (No explanation is needed just mark the correct FDs)
  - (a)  $ABE \rightarrow ACDE$  (b)  $B \rightarrow C$  (c)  $ABCE \rightarrow ADE$  (d)  $ABE \rightarrow AE$

## Q3. Inferring FDs [27 points]

Consider the relations below and sets of functional dependencies that hold in those relations and answer the following questions. For each sub-part, it is enough for you to list only completely non-trivial FDs with a single attribute on the right hand side.

**Note** that 'candidate key' means just 'key' (i.e. both words are interchangeable). A candidate key (or simply key) should imply the entire relation and should be minimal. On the other hand, a 'superkey' is any super-set of a candidate key.

Q3.1 R1(A, B, C, D) with FDs AB  $\rightarrow$  C, C  $\rightarrow$  D, and D  $\rightarrow$  A.

- Q3.1.1 (3 points) What are all the non-trivial FDs that follow from the given FDs?
- Q3.1.2 (3 points) What are all the keys (i.e. candidate keys) of R1?
- Q3.1.3 (3 points) What are all the superkeys for R1 that are not keys?

Q3.2 R2(A, B, C, D) with FDs A $\rightarrow$  B, B  $\rightarrow$  C, and B  $\rightarrow$  D.

- Q3.2.1 (3 points) What are all the non-trivial FDs that follow from the given FDs?
- Q3.2.2 (3 points) What are all the keys (i.e candidate keys) of R2?
- Q3.2.3 (3 points) What are all the superkeys for R2 that are not keys?

Q3.3 R3(A, B, C, D) with FDs ABD  $\rightarrow$  C, BC  $\rightarrow$  D, CD  $\rightarrow$  A, and AD  $\rightarrow$  B.

- Q3.3.1 (3 points) What are all the non-trivial FDs that follow from the given FDs?
- Q3.3.2 (3 points) What are all the keys (i.e. candidate keys) of R3?
- Q3.3.3 (3 points) What are all the superkeys for R3 that are not keys?

#### Q4. Projection of FDs [26 points]

Consider the relation  $R = \{A, B, C, D, E, F\}$  and the following set of FDs on R:

$$F = \{AB \rightarrow C, AC \rightarrow B, AD \rightarrow E, B \rightarrow D, BC \rightarrow A, E \rightarrow F\}$$

Consider also the relations: R1 (A, B, C), R2 (A, B, C, E, F), and R3 (A, B, C, D).

- Q4.1 (2x3=6 points) Compute the projection of the FD-set *F* on each of the three relations R1, R2 and R3. Only write down a minimal cover.
- Q4.2 (2x3=6 points) What are *all* the candidate keys (i.e. the keys) for each of the three relations R1, R2 and R3?

- Q4.3 (2x2=4 points) For each of the three relations, indicate if it is (A) BCNF and (B) 3NF. Explain your answer.
- Q4.4 (6 points) Is it possible to decompose each R2 and R3 into new schemas such that they are in BCNF and the decomposition is lossless and dependency preserving? Show your work.
- Q4.5 (4 points) Decompose R into multiple relations so that they are in 3NF and the decomposition is lossless and dependency preserving. Use the standard 3NF synthesis algorithm. Compare your result with Q4.4.

### Q5. Lossy decompositions [8 points]

Consider a relation R(A, B, C, D, E). Assume R satisfies the following functional dependencies:  $C \rightarrow E$  and  $A \rightarrow C$ .

Suppose we decompose *R* into these two relations: R1(A, B, C) and R2(C, D, E). So you can think of  $R1 = \prod_{\{A,B,C\}} (r)$  and  $R2 = \prod_{\{C,D,E\}} (r)$ .

- Q5.1 (5 points) Show that this decomposition is not a lossless-join by giving an example. That is, give an instance *r* of *R* (i.e., give a set of tuples for *R*) such that both dependencies hold on *r* but  $r \neq R1 \triangleright \lhd R2$ . Give the simplest instance you can find. Also show your work (i.e. compute *R1* and *R2* and the join).
- Q5.2 (3 points) Can you show that the decomposition is lossy by using the theorem discussed in class? Again, show your work.

## Q6. FDs in English language [24 points]

Consider the following relation, which stores information of Facebook dataset.

*FacebookDB* (u\_id, u\_name, u\_location, u\_friends\_count, u\_posts\_count, u\_post\_importance, u\_age, u\_influence, p\_id, p\_text, p\_location, p\_like\_count, p\_reshare\_count, p\_parent\_id, c\_id, c\_text, c\_reply\_count, c\_like\_count)

Now consider the following constraints in English:

Statement 1 Every user has a unique user id (u\_id). Other information for users such as name (u\_name), location (u\_location), number of friends (u\_friends\_count), and posts (u\_post\_count), importance of posts (u\_post\_importance), age, and the influence value are also kept.

Statement 2 Each user can have multiple friends, and	posts.
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- Statement 3 For every post we have p\_id (which is unique), the text, the location, the number of likes, and re-shares. If the post itself is a re-share of another post, we have the parent post id.
- Statement 4 There might be several comments for each post. The comments have a unique id (c\_id), text, number of replies to it (c\_reply\_count), and number of likes (c\_like\_count).
- Statement 5 The importance of a user's posts (u\_post\_importance) is computed based on number of friends and posts (u\_friends\_count, u\_posts\_count) of the user.
- Statement 6 The user's influence (u\_influence) is defined using the age and the importance of the user's posts (u\_age, u\_post\_importance).

Clearly this is not a good relational design, as it contains redundancies and update, insertion and deletion anomalies. Hence we want to decompose it into a good schema.

- Q6.1 (6 points) List the functional dependencies for this relation that you can construct using the English description and also specifically mention the corresponding statement number(s) you used to formulate each FD. *Note*: some statement(s) might not result in a FD.
- Q6.2 (3 points) Rigorously prove (using Armstrong's axioms) that you can derive the FD: u\_friends\_count, u\_posts\_count, u\_age → u\_influence from the set of the FDs you found in Q5.1. Show your steps.
- Q6.3 (10 points) Is the above relation in BCNF using the FDs you found in Q5.1? If not, decompose it into a set of BCNF relations (using our standard algorithm). Is your resulting decomposition dependency-preserving too?
- Q6.4 (5 points) Decompose the original relation into a set of 3NF relations instead (using the standard synthesis algorithm). Show your steps. Is your resulting decomposition in BCNF as well? Explain your answer.