

**Homework 7: Transactions, Logging and Recovery**  
**(due April 24th, 2014, 3:30pm, in class—hard-copy please)**

**Reminders:**

- a. Out of 100 points. Contains 4 pages.
- b. Rough time-estimates: 2~4 hours.
- c. Please type your answers. Illegible handwriting may get no points, at the discretion of the grader. Only drawings may be hand-drawn, as long as they are neat and legible.
- d. There could be more than one correct answer. We shall accept them all.
- e. Whenever you are making an assumption, please state it clearly.
- f. Lead TA for this HW: Pranav Nakate.

**Q1. Serializability [30 points]**

Consider the following schedules. The actions are listed in the order they are scheduled, and prefixed with the transaction name.

S1: T1:R(X), T2:R(Y), T1:W(X), T2:W(Y), T1:W(Y), T2:R(Z)

S2: T1:R(X), T2:W(X), T2:W(Y), T3:W(Y), T1:W(Y), T1:R(X), T3:R(Y)

Q1.1. (15 points) Consider the schedule S1.

- A. (5 points) Draw the precedence graph for S1.
- B. (5 points) Is S1 a conflict serializable schedule? If yes, what is the equivalent serial schedule? If no, explain in 1-2 lines.
- C. (5 points) From the following options, choose which of the anomalies is present in S1 (if any). Explain in 1-2 lines which actions cause that anomaly.
  - a. Dirty read (WR Conflict)
  - b. Unrepeatable read (RW Conflict)
  - c. Lost update (WW Conflict)

Q1.2. (15 points) Consider the schedule S2.

- A. (5 points) Draw the precedence graph for S2.
- B. (5 points) Is S2 a conflict serializable schedule? If yes, what is the equivalent serial schedule? If no, explain in 1-2 lines.
- C. (5 points) From the following options, choose which of the anomalies is present in S2 (if any). Explain in 1-2 lines which actions cause that anomaly.
  - a. Dirty read (WR Conflict)

- b. Unrepeatable read (RW Conflict)
- c. Lost update (WW Conflict)

*Note:* The anomalies due to interleaved execution of transactions are explained on slide 38-44 of the lecture 17 or page 526-529 of the textbook. The concept of conflict serializable schedules is explained on slide 46 of lecture 17 or page 550-551. The precedence graphs are explained on slide 51-52 of the lecture 17 or page 550-551 of the textbook.

## Q2. Locking Protocols [20 points]

Consider the following schedules. The actions are listed in the order they are scheduled, and prefixed with the transaction name.

S1: T1:R(Z), T2:W(X), T2:R(Y), T1:W(X), T1:Commit, T2:Commit

S2: T1:R(Z), T2:W(Y), T2:W(X), T2:Commit, T3:R(X), T3:W(Y), T1:W(Z), T1:Commit, T3: Commit.

S3: T2:W(A), T3:W(B), T1:W(C), T3:R(A), T2:R(B), T2:W(D), T1:R(D), T1:Commit, T2:Commit, T3:Commit.

- Q2.1. (10 points) Consider schedule S1.
- A. (5 points) Write the schedule S1 in a table format and draw the precedence graph.
  - B. (5 points) Will the actions of S1 be allowed by the following concurrency control protocols?
    - a. 2PL
    - b. Strict 2PL
 If YES, show in the table form where the lock requests can happen; If NO, explain briefly in 1-2 lines.
- Q2.2. (10 points) Consider schedule S2.
- A. (5 points) Write the schedule S2 in a table format and draw the precedence graph.
  - B. (5 points) Will the actions of S2 be allowed by the following concurrency control protocols?
    - a. 2PL
    - b. Strict 2PL
 If YES, show in the table form where the lock requests can happen; If NO, explain briefly in 1-2 lines.
- Q2.3. (10 points) Consider schedule S3.
- A. (5 points) Write the schedule S3 in a table format and draw the precedence graph.
  - B. (5 points) Will the actions of S3 be allowed by the following concurrency control protocols?

- a. 2PL
  - b. Strict 2PL
- If YES, show in the table form where the lock requests can happen; If NO, explain briefly in 1-2 lines.

*Note:* “2PL” and “Strict 2PL” protocols are explained on slide 18-33 of the lecture 18 or textbook page 551-552. “Lock Management” which explains how the lock requests are granted is explained on slide 36-41 of the lecture 18 or on page 553-554 of the textbook.

### Q3. Deadlock Management [20 points]

Consider the following sequence of actions, listed in the order it is submitted to the DBMS (S is a shared lock, X is an exclusive lock):

S1: T1:S(A), T2:X(C), T3:S(A), T1:X(C), T3:X(B), T2: S(A), T3: X(A)  
 S2: T1:S(A), T2:X(B), T3:S(C), T2:S(C), T1:X(B), T3:S(A),T2:X(A), T3:X(A)

For S1 and S2 as given above, answer the following questions:

- Q3.1. (5 points) For S1, write whether lock requests of its actions will be granted or blocked by the lock manager.
- Q3.2. (5 points) Draw the waits-for graphs for S1 and write if the schedule will result in a deadlock condition. If there is no deadlock condition, write the order of completion of the schedule. Explain in 1-2 lines.
- Q3.3. (5 points) For S2, write whether the lock requests of its actions will be granted or blocked by the lock manager.
- Q3.4. (5 points) Draw the waits-for graphs for S2 and write if the schedule will result in a deadlock condition. Explain in 1-2 lines.

*Note:* The implementation of lock and unlock requests is given in detail on slide 5-10 of lecture 18 or on page 554 of the textbook. The waits-for graphs are explained on slide 13 of lecture 18 or on page 556 of the textbook.

### Q4. Logging [23 points]

Consider Exercise 18.4 in your textbook. Here we show a slightly easier to read Log table.

| LSN | Log entry                 | prevLSN | undonextLSN |
|-----|---------------------------|---------|-------------|
| 00  | begin checkpoint          |         |             |
| 10  | end checkpoint            |         |             |
| 20  | update: T1 writes Page P2 |         |             |
| 30  | update: T1 writes Page P1 |         |             |
| 40  | update: T2 writes Page P5 |         |             |
| 50  | update: T3 writes Page P3 |         |             |
| 60  | T3 commit                 |         |             |
| 70  | update: T2 writes Page P5 |         |             |

|    |                           |  |  |
|----|---------------------------|--|--|
| 80 | update: T2 writes Page P3 |  |  |
| 90 | T2 abort                  |  |  |

- Q4.1. (10 points) Write the appropriate values in the prevLSN and undonextLSN columns.
- Q4.2. (5 points) Describe the actions taken to rollback transaction T2.
- Q4.3. (8 points) Show the log after T2 is rolled back, including all the prevLSN and undonextLSN values in the log records.

*Note:* The Log record and CLR are explained in detail on page 582-585 of the textbook and also in Lecture 20.

### Q5. Crashing now [5 points]

Consider the following log. It is the same one as given in Q4 above, but with a CRASH after LSN 90.

| LSN | Log entry                 | prevLSN | undonextLSN |
|-----|---------------------------|---------|-------------|
| 00  | begin checkpoint          |         |             |
| 10  | end checkpoint            |         |             |
| 20  | update: T1 writes Page P2 |         |             |
| 30  | update: T1 writes Page P1 |         |             |
| 40  | update: T2 writes Page P5 |         |             |
| 50  | update: T3 writes Page P3 |         |             |
| 60  | T3 commit                 |         |             |
| 70  | update: T2 writes Page P5 |         |             |
| 80  | update: T2 writes Page P3 |         |             |
| 90  | T2 abort                  |         |             |
|     | CRASH                     |         |             |

The recovery manager now sees this log (of course with the correct prevLSN and undonextLSN values filled in) after the CRASH. Which transactions will be redone and undone?

*Note:* The concept of crash recovery is explained in detail on page 588-592 of the textbook and also in Lecture 19.

### Q6. Character fights [2 points]

Who would win the fight, if James Bond fought with Jason Bourne? Explain your answer. Is there anyone who can defeat both of them? ☺