CS 4104: Data and Algorithm Analysis

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Countable vs. Uncountably Infinite Sets

Two sets have the **same cardinality** if there is a bijection between them.

Notation: |A| = |B|.

This concept can also be applied to infinite sets.

Example: Let Odd and Even be the sets of odd and even natural numbers, respectively.

Then, |Odd| = |Even| because the function $f : |\text{Odd} \rightarrow \text{Even}|$ defined by f(x) = x - 1 is a bijection.

How about $|\text{Even}| = |\mathbb{N}|$?

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Counting Infinite Sets

A set C is **countable** if it is finite or if $|C| = |\mathbb{N}|$.

If a set is not countable, then it is **uncountable**.

If A is a finite alphabet, then A^* is countably infinite.

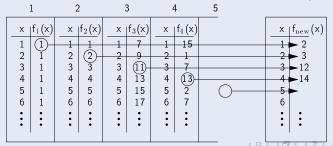
Proof: Arrange the strings in order by length, and within a given length by alphabetical order. This provides a bijection.

As a corollary, the set of all computer programs is countable.

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More Functions than Programs

- Consider set of functions f(x) = y for x, y natural numbers.
- The set of such functions is uncountable.
- Diagonalization argument
- Not all functions on natural numbers are computable.



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Halting Problem for Programs

Does the following terminate?

```
while (n > 1)
  if (ODD(n))
    n = 3 * n + 1;
  else
    n = n / 2;
```

Can a **C**++ program be written to solve the following problem?

Halting Problem:

- Input: A program *P* and input *X*.
- Output: "Halts" if P halts when run with X as input.
 "Does not Halt" otherwise.

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Halting Problem Proof

Theorem: There is no program to solve the Halting Problem.

Proof: (by contradiction).

Assumption: There is a **C**++ program that solves the Halting Problem.

```
bool halt(char* prog, char* input)
{
  Code to solve halting problem
  if (prog does halt on input) then
    return(TRUE);
  else
    return(FALSE);
}
```

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Two More Procedures

```
bool selfhalt(char *proq) {
  // Return TRUE if program halts
  // when given itself as input.
  if (halt(proq, proq))
    return(TRUE);
  else
    return(FALSE);
void contrary(char *prog) {
  if (selfhalt(prog))
    while(TRUE); // Go into an infinite loop
```

The Punchline

- What happens when function contrary is run on itself?
- Case 1: selfhalt returns TRUE.
 - contrary will go into an infinite loop.
 - ▶ This contradicts the result from selfhalt.
- selfhalt returns FALSE.
 - ► contrary will halt.
 - This contradicts the result from selfhalt.
- Either result is impossible.
- The only flaw in this argument is the assumption that halt exists.
- Therefore, halt cannot exist.

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Computability Reduction Proof

Given arbitrary program M, does it halt on the EMPTY input?

This is uncomputable. Proof:

- Suppose that program M₀ determines if M halts on the EMPTY input.
- Given arbitrary program M and string w, we can create a new program M_w that operates as follows on empty input:
 - Write w into a static variable.
 - Simulate the execution of M.
- So, we can take arbitrary program M and string w, create M_w , and invoke M_0 on M_w (with empty input) to solve the original halting problem.
- Thus, M_0 must not exist.

Another Reduction Proof

Does there exist ANY input for which an arbitrary program halts?

Proof that this is uncomputable:

- Suppose that program M₀ could decide if arbitrary program M halts on ANY input.
- We can take an arbitrary program M and string w, and modify it so that it ignores its input before proceeding.
- Thus, arbitrary program *M* is modified to be *M'* that effectively is *M* operating on the empty input.
- Thus, we can take arbitrary program M and string w, modify it to become M' and feed that to M_0 to solve the problem of deciding if M halts on the empty input.
- We already know that is undecidable.
- Thus, M₀ cannot exist.

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Other Noncomputable Functions

Does a program halt on EVERY input?

Do two programs compute the SAME function?

Does a particular line in a program get executed?

Does a program compute a particular function?

Does a program contain a "computer virus"?

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