FP Foundations, Scheme (2)

In Text: Chapter 15

Functional programming

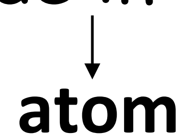
- LISP: John McCarthy 1958 MIT
  - List Processing => Symbolic Manipulation
- First functional programming language
  - Every version after the first has imperative features, but we will discuss the functional subset

LISP Data Types

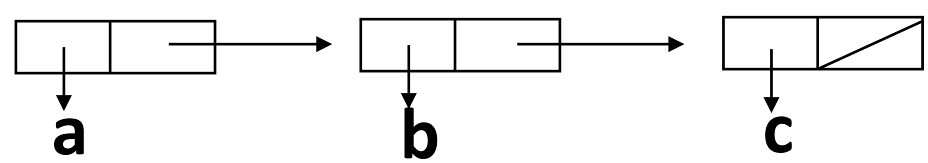
- There are only two types of data objects in the original LISP
  - Atoms: symbols, numbers, strings, …
    - E.g., a, 100, “foo”
  - Lists: specified by delimiting elements with parentheses
    - Simple lists: elements are only atoms
      - E.g., (A B C D)
    - Nested lists: elements can be lists
      - E.g., (A (B C) D (E (F G)))

LISP Data Types

- Internally, lists are stored as **single-linked list** structures
  - Each node has two pointers: one to element, the other to next node in the list
  - Single atom:

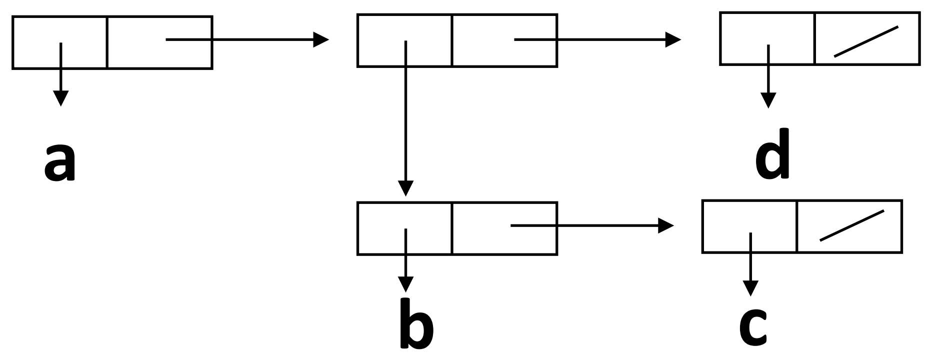


  - List of atoms: (a b c)



LISP Data Types

– List containing list (a (b c) d)



Scheme

- Scheme is a dialect of LISP, emerged from MIT in 1975
- Characteristics
  - simple syntax and semantics
  - small size
  - exclusive use of static scoping
  - treating functions as first-class entities
    - As first-class entities, Scheme functions can be the values of expressions, elements of lists, assigned to variables, and passed as parameters

Interpreter

- Most Scheme implementations employ an interpreter that runs a “read-eval-print” loop
  - The interpreter repeatedly reads an expression from a standard input, evaluates the expression, and prints the resulting value

Primitive Numeric Functions

- Primitive functions for the basic arithmetic operations: +, -, *, /
  - + and * can have zero or more parameters. If * is given no parameter, it returns 1; if + is given no parameter, it returns 0
  - - and / can have two or more parameters
  - Prefix notation

Expression	Value
42	42
(* 3 6)	18
(+ 1 2 3)	6
(sqrt 16)	4

Numeric Predicate Functions

- Predicate functions return Boolean values (#T or #F): =, <>, >, <, >=, <=, EVEN?, ODD?, ZERO?

Expression	Value
(= 16 16)	#T
(even? 29)	#F
(> 10 (* 2 4))	
(zero? (-10(* 2 5)))	

Type Checking

- Dynamic type checking
- Type predicate functions

```
(boolean? x)  ; Is x a Boolean?
(char? x)
(string? x)
(symbol? x)
(number? x)
(pair? x)
(list? x)
```

Lambda Expression

- E.g., lambda(x) (* x x) is a nameless function that returns the square of its given numeric parameter
- Such functions can be applied in the same ways as named functions
  – E.g., ((lambda(x) (* x x)) 7) = 49
- It allows us to pass function definitions as parameters

"define"

- To bind a name to the value of a variable:
  (define symbol expression)
  - E.g., (define pi 3.14159)
  - E.g., (define two_pi (* 2 pi))
- To bind a function name to an expression:
  (define (function_name parameters)
  (expression)
)
  - E.g., (define (square x) (* x x))

“define”

- To bind a function name to a lambda expression

  (define function_name
  (lambda_expression)
)

  – E.g., (define square (lambda (x) (* x x)))

Control Flow

- Simple conditional expressions can be written using if:
  - E.g. (if (< 2 3) 4 5) => 4
  - E.g., (if #f 2 3) => 3

Control Flow (cont'd)

- It is modeled based on the evaluation control used in mathematical functions:

```
(COND
    (predicate_1 expression)
    (predicate_2 expression)
    …
    (predicate_n expression)
    [ELSE expression]
)
```

An Example

$$f(x) = \begin{cases} 1 \; if \;\; x = 0 \\ x * f(x-1) \; if \;\; x > 0 \end{cases}$$

```
( define ( factorial x )
      ( cond
            (( < x 0 ) #f)
            (( = x 0 ) 1)
            ( #t (* x  (factorial (- x 1)))) ; or else (…)
      )
)
```

Bindings & Scopes

- Names can be bound to values by introducing a nested scope
- `let` takes two or more arguments:
  - The first argument is a list of pairs
    - In each pair, the first element is the name, while the second is the value/expression
  - Remaining arguments are evaluated in order
  - The value of the construct as a whole is the value of the final argument
  - E.g. (let ((a 3)) a)

`let` Examples

- E.g., (let ((a 3)
                (b 4)
                (square (lambda (x) (* x x)))
                (plus +))
        (sqrt (plus (square a) (square b))))
- The scope of the bindings produced by `let` is its second and following arguments

`let` Examples

- E.g., (let ((a 3))
```
            (let ((a 4)
                  (b a))
                (+ a b))) => ?
```

- b takes the value of the outer a, because the defined names are visible "all at once" at the end of the declaration list

`let*` Example

- `let*` makes sure that names become available “one at a time”
- E.g., (let*((x 1) (y (+ x 1)))
  (+ x y)) => ?

Functions

- quote: identity function
  - When the function is given a parameter, it simply returns the parameter
  - E.g., (quote A) => A
    (quote (A B C)) => (A B C)
- The common abbreviation of quote is apostrophe (')
  - E.g., (' A) => A
    (' (A B C)) => (A B C)

List Functions

- car: returns the first element of a given list
  - E.g., (car '(A B C)) => A
    (car '((A B) C D)) => (A B)
    (car 'A) => ?
    (car '(A)) => ?
    (car '()) => ?

List Functions

- cdr: returns the remainder of a given list after the first element has been removed
  - E.g., (cdr '(A B C)) => (B C)
    (cdr '((A B) C D)) => (C D)
    (cdr 'A) => ?
    (cdr '(A)) => ?
    (cdr '()) => ?

List Functions

- cons: concatenates an element with a list
- cons builds a list from its two arguments
  - The first can be either an atom or a list
  - The second is usually a list
  - E.g., (cons 'A '()) => (A)
    (cons 'A '(B C)) => (A B C)
    (cons '() '(A B)) => ?
    (cons '(A B) '(C D)) => ?
  - How to compose a list (A B C) from A, B, and C?

List Functions

- Note that cons can take two atoms as parameters, and return a dotted pair
  - E.g., (cons 'A 'B) => (A . B)
  - The dotted pair indicates that this cell contains two atoms, instead of
    an atom + a pointer
    or
    a pointer + a pointer

More Predicate Functions

- The following returns #t if the symbolic atom is of the indicated type, and #f otherwise
  - E.g., (symbol? 'a) => #t
    (symbol? '()) => #f
  - E.g., (number? '55) => #t
    (number? 55) => #t
    (number? '(a)) => #f
  - E.g., (list? '(a)) => #t
  - E.g., (null? '()) => #t

More Predicate Functions

- eq? returns true if two objects are equal through pointer comparison
  - Guaranteed to work on symbols
  - E.g., (eq? 'A 'A) => #T
(eq? 'A '(A B)) => #F
- equal? recursively compares two objects to determine if they are equal
  - The objects can be atoms or lists

How do we implement equal?

```
(define (simple? atm)
  (cond
    ((list? atm) (null? atm))
    (else #T)
  )
)
```

```
(define (equal? lis1 lis2)
  (cond
    ((simple? lis1) (eq? lis1 lis2))
    ((simple? lis2) #F)
    ((equal? (car lis1) (car lis2))
        ((equal? (cdr lis1) (cdr lis2))
    (else #F)
  )
)
```

More Examples

```
(define (member? atm lis)
  (cond
    ((null? lis) #F)
    ((eq? atm (car lis)) #T)
    (else (member? atm (cdr lis)))
  )
)
```

What is returned for the following function?

(member? 'b '(a (b c)))

```
(define (append lis1 lis2)
  (cond
    ((null? lis1) lis2)
    (else (cons (car lis1)
                (append(cdr lis1) lis2)))
  )
)
```

Is lis2 appended to lis1, or lis1 prepended to lis2?

An example: apply-to-all function

```
(define (mapcar fctn lis)
   (cond
      ((null? lis) '())
      (else (cons (fctn (car lis))
                  (mapcar fctn (cdr lis)) ))
)
```

Project 3: A Scheme Parsing Program

- Consider the grammar G = (S, N, T, P) where

S = (Program)
N = (statement_list, statement, expr, symbol, op)
T = (if, bool, then, while, id, const, =, +, -, *, /)

P = (Program → statement_list

statement_list → statement statement_list
→ statement

statement → if bool then statement_list
→ while bool statement_list
→ id = expr

expr → symbol op symbol

symbol → id
→ const

op → + | - | * | /

- Write a Scheme program that correctly parses all valid programs in L(G). *The Scheme program will report (a) the total number of statements in the program, and (b) the maximum nested depth for a program.*
- E.g., given ((id = id - const)), your program will output: (numberofstatements: 1 maximumdepth: 0)

- You can assume that
  - You will be given only valid programs.
  - Each program will be provided as a parenthesized list of statements, each of which is included in its own parentheses.
  - The nesting of parentheses is used to indicate subordinate (or block) statement(s).

Some Hints

- Define two functions to separately count (1) the number of statements and (2) the maximum nested depth

- When counting the number of statements, the function should
  - Check whether the input parameter is an empty list
  - If not, check whether the first element in the list is an empty list
  - If not, obtain the first element in the list, check the statement type and proceed accordingly. The recursive function calls may be involved.

- When calculating the maximum nested depth, the function should do similar checks as what is mentioned above
  - Notice that when an if-statement has both then- and else-branches, you need to compare the depth counts of both branches, obtains the larger number as the maximum nested depth at the current level

On Rlogin …

- You can use command “racket -i” to launch the interactive mode of Racket, and use “(exit)” to exit that mode
- Please name your main function “parser”, which function takes only one argument.
- The program will be tested on Rlogin via “plt-r5rs < filename”
- The source file can be a .txt file