

ABSTRACT DATA TYPES

- **Based on the fundamental concept of ABSTRACTION:**
 - process abstraction
 - data abstraction
- **Both provide:**
 - information hiding
 - reliability
 - security
 - reuse

ADT PROPERTIES

- **Defines a collection of objects, and**
- **a Set of applicable operations**

- **Representation of objects is hidden**
- **Operations by outsiders is restricted to only those operations that are visible**

COMMON FORMAT OF DESCRIPTION

- **SPECIFICATION**

- Defines type name and parameters
- Names visible operations and results types

- **BODY**

- Describes syntax of type objects
- Describes visible and hidden operations

IMPLEMENTATIONS

- **SIMULA 67 - first to introduce classes, retrospectively recognized to be ADTs**
- **CLU - an experimental language that introduced ADTs**
- **Modula-2 - first generally accessible implementations in a widely used language**
- **Smalltalk - used ADTs as basis for objects**
- **Ada - used ADTs in packages**
- **C++ - added ADTs to C**

ENCAPSULATION and INSTANTIATION

- **ENCAPSULATION**

- the syntax of the specification as a separate module
- builds a “fire-wall” around the type
- provides a reusable, portable object
- the development of the idea of a type as an object

- **INSTANTIATION**

- the creation of an instance of the type
- the operation of importing an ADT into another program unit
- may include initialization actions
- scope may be limited to the lifetime of the user module or may be specified

Abstract Data Types - details

- **Type representation and operations on that type are defined together.**
- **Representation is hidden from user of the type -- objects of type t can only be manipulated by operations defined for t.**
- **Advantages of user-defined ADTs**
 - encapsulation
 - protection
 - extensibility
- **We'll look at three languages:**
 - Simula 67
 - Ada
 - Modula-2

Simula 67: Classes

- **A class consists of:**
 - variable declarations
 - procedure declarations
 - code (for initialization)
- **If C is a class with variables $x_1 \dots x_n$ and procedures $p_1 \dots p_k$, an instance of C is a dynamically created object, say r.**

```
ref (C) r;  
...  
r := new C;  
...  
... r.xi ...  
... r.pj(y1 ... ym) ...
```



Stack Example

```
class stack;  
  begin  
    integer array a(1 .. 100);  
    integer top;  
  
    boolean procedure empty;  
    ...  
  end;  
  
  procedure push (element);  
  ...  
  end;  
  
  procedure pop;  
  ...  
  end;  
  
  procedure look;  
  ...  
  end;  
  
  top := 0; ← initialization code  
end stack;
```


Using the Stack Class

```
ref (stack) s1,s2;  
...  
s1 :- new stack;  
s2 :- new stack;  
s1.pop;      -- error  
s1.push(5);  
...  
s1.look;     -- 5  
...  
s2.look ;    -- error  
...
```

→ **But no protection!**

```
s2.a(4) := 1000;  allowed, but unsafe.  
s1.top := 0; 
```

Inheritance in Simula

- **If x is a subclass of y, then instances of x have all of x's attributes plus all of y's attributes.**

→ *x inherits* the attributes of y.

- **Example: defining a heterogeneous stack**

```
class stack_mem
  begin ref(stack_mem) next_mem
    next_mem :- none
  end stack_mem;
```

Example Continued: Define stack

```
class stack;
  begin
    ref (stack_mem) first;
    ref (stack_mem) procedure top
      top :- first;
  procedure pop;
    if not(empty) then
      first :- first.next_mem;
  boolean procedure empty;
    empty :- (first = = none);
  procedure push(e);
    ref(stack_mem) e;
  begin
    if first /= none then
      e.next_mem :- first;
    first :- e;
  end
  first :- none;
end stack;
```

Example Continued: Stackable Objects

- **Stackable objects must be instances of a *subclass* of `stack_mem`:**

```
stack_mem class complex(. . .) -- declare complex as
                                subclass of stack_mem
...
end complex
```

- **Another example:**

```
class mammal;
mammal class dog;
mammal class cat;
dog class golden_retriever;
```

Packages in Ada

- **Two parts:**
 - **specification: provides interface, defines visibility.**
 - **body: provides implementation**
- **Important:**
 - **Support separate compilation so that if package p1 uses package p2, p1 can be compiled given only the specification part of p2.**

Package Example

```
package stack is                                     -- the specification
    type stacktype;
    function empty (s: in stacktype)
        return boolean;
    procedure push (e: in integer;
        s: in out stacktype);
    procedure pop (s: in out stacktype);
    function top(s: in stacktype)
        return integer;
end stack;
```

```
package body stack is                               -- the body
    type stacktype is . . .
    function empty (. . .) is . . .
    ...
```

Package example (continued)

- **Does our separate compilation rule hold:**
 - **No!**
 - **Definition for stacktype must be in the interface too.**
- **Problem: We didn't want stacktype's definition to be exported.**
 - **Solution: Divide the specification into a *public* part and a *private* part.**

New Specification for stack

```
package stack is                                     -- the visible part
  type stacktype is private;
  function empty(. . .) . . .
  procedure push . . .
  ...
private                                             -- the private part
  type list_type is array (1..100) of int;
  type stacktype is
    record
      list : list_type;
      top : integer range 0..100 := 0
    end record;

end stack;
```


Using Packages

```
with stack;  
procedure p is  
  s : stack.stacktype;  
  begin  
  . ..  
  stack.push(4,s);  
  ...stack.top(s)...;  
  ...  
  end
```

OR...

```
with stack; use stack;  
procedure p is  
  s : stacktype;  
  begin  
  ...  
  push(4,s);  
  ...top(s)...;  
  ...  
  end
```

Ada Generic -- Abstract Package

generic

-- A private generic type means assignment and equality must be
-- defined on that type

type Elem is private ;

package List is

type T is private ;

-- Create operation is implicit. Lists created by declaration

procedure Head (L: T ; V: out Elem ; Err: out ERROR_INDICATOR) ;

-- Length can't fail so no need for error indicator

function Length (L: T) return NATURAL ;

procedure Tail (L: T ; LT: out T; Err: out ERROR_INDICATOR) ;

-- Cons can't fail so no need for error indicator

function Cons (L: T ; V: Elem) return T ;

private

-- an Ada access type corresponds to a Pascal pointer

-- the entity referenced by the pointer is defined in the package body

-- In this case, it would be a record with one field pointing to the next

-- list element

type LISTREC ;

type T is access LISTREC ;

end List ;

Modules in Modula-2

- **Very similar to Ada packages, but only pointer types can be exported.**

```
Definition module stack;                                -- public
    type stacktype;
    procedure empty . . .
    ...
end stack;
```

```
Implementation module stack;                            -- private
    type stacktype = pointer to record
        list : . . .
        topsub: . . .
```

Modula-2 Modules (continued)

- **What are the repercussions of this design decision?**
 - **separate compilation is easy (+)**
 - **module must supply a creation/initialization routine (-)**
 - **extra use of pointers (-)**