ABSTRACT DATA TYPES

- Based on the fundamental concept of ABSTRACTION:
  - process abstraction
  - data abstraction
- Both provide:
  - information hiding
  - reliability
  - security
  - reuse

ADT PROPERTIES

- Defines a collection of objects, and
- a Set of applicable operations
- Representation of objects is hidden
- Operations by outsiders is restricted to only those operations that are visible

COMMON FORMAT OF DESCRIPTION

- **SPECIFICATION**
  - Defines type name and parameters
  - Names visible operations and results types
- **BODY**
  - Describes syntax of type objects
  - Describes visible and hidden operations

IMPLEMENTATIONS

- **SIMULA 67 - first to introduce classes, retrospectively recognized to be ADTs**
- **CLU - an experimental language that introduced ADTs**
- **Modula-2 - first generally accessible implementations in a wdiely used language**
- **Smalltalk - used ADTs as basis for objects**
- **Ada - used ADTs in packages**
- **C++ - added ADTs to C**

ENCAPSULATION and INSTANTIATION

- **ENCAPSULATION**
  - **the syntax of the specification as a separate module**
  - **builds a "fire-wall" around the type**
  - **provides a reusable, portable object**
  - **the development of the idea of a type as an object**
- **INSTANTIATION**
  - **the creation of an instance of the type**
  - **the operation of importing an ADT into another program unit**
  - **may include initialization actions**
  - **scope may be limited to the lifetime of the user module or may be specified**

Abstract Data Types - details

- **Type representation and operations on that type are defined together.**
- **Representation is hidden from user of the type -- objects of type t can only be manipulated by operations defined for t.**
- **Advantages of user-defined ADTs**
  - **encapsulation**
  - **protection**
  - **extensibility**
- **We'll look at three languages:**
  - **Simula 67**
  - **Ada**
  - **Modula-2**

Simula 67: Classes

- **A class consists of:**
  - **variable declarations**
  - **procedure declarations**
  - **code (for initialization)**
- **If C is a class with variables $x_1 \ldots x_n$ and procedures $p_1 \ldots p_k$, an instance of C is a dynamically created object, say r.**

```
ref (C) r;
...
r :- new C;
...
. . .r.x_i . . .
. . .r. p_j (y_1 . . . y_m). . .
```


Stack Example

```
class stack;
   begin
      integer array a(1 . . 100);
      integer top;

   boolean procedure empty;
      ...
   end;

   procedure push (element);
   ...
   end;

   procedure pop;
   ...
   end;

   procedure look;
   ...
   end;

                          <-- initialization code
top := 0;
end stack;
```

Using the Stack Class

```
ref (stack) s1,s2;
. . .
s1 :- new stack;
s2 :- new stack;
s1.pop;        -- error
s1.push(5);
. . .
s1.look;       -- 5
. . .
s2.look ;      -- error
. . .
```

➡ **But no protection!**

```
s2.a(4) := 1000;
s1.top := 0;
```

allowed, but unsafe.

Inheritance in Simula

- **If x is a subclass of y, then instances of x have all of x's attributes plus all of y's attributes.**
  - ➡ **x *inherits* the attributes of y.**
- **Example: defining a heterogeneous stack**

```
class stack_mem
    begin ref(stack_mem) next_mem
        next_mem :- none
    end stack_mem;
```

Example Continued: Define stack

```
class stack;
    begin
        ref (stack_mem) first;
        ref (stack_mem) procedure top
            top :- first;
    procedure pop;
        if not(empty) then
            first :- first.next_mem;
    boolean procedure empty;
        empty :- (first = = none);
    procedure push(e);
        ref(stack_mem) e;
    begin
        if first =/= none then
            e.next_mem :- first;
        first :- e;
    end
first :- none;
end stack;
```


Example Continued: Stackable Objects

- **Stackable objects must be instances of a *subclass* of stack_mem:**

```
stack_mem class complex(. . .)   -- declare complex as
                                    subclass of stack_mem
...
end complex
```

- **Another example:**

```
class mammal;
mammal class dog;
mammal class cat;
dog class golden_retriever;
```

Packages in Ada

- **Two parts:**
  - **specification: provides interface, defines visibility.**
  - **body: provides implementation**
- **Important:**
  - **Support separate compilation so that if package p1 uses package p2, p1 can be compiled given only the specification part of p2.**

Package Example

```
package stack is                          -- the specification
   type stacktype;
   function empty (s: in stacktype)
                  return boolean;
   procedure push (e: in integer;
                  s: in out stacktype);
   procedure pop (s: in out stacktype);
   function top(s: in stacktype)
                  return integer;
end stack;

package body stack is                     -- the body
   type stacktype is . . .
   function empty (. . .) is . . .
...
```

Package example (continued)

- **Does our separate compilation rule hold:**
  - **No!**
  - **Definition for stacktype must be in the interface too.**
- **Problem: We didn't want stacktype's definition to be exported.**
  - **Solution: Divide the specification into a *public* part and a *private* part.**

New Specification for stack

```
package stack is                                    -- the visible part
    type stacktype is private;
    function empty(. . .) . . .
    procedure push . . .
    ...
private                                             -- the private part
    type list_type is array (1. .100) of int;
    type stacktype is
        record
            list : list_type;
            top : integer range 0. .100 := 0
        end record;

end stack;
```

Using Packages

```
with stack;
procedure p is
    s : stack.stacktype;
    begin
.   ..
    stack.push(4,s);
    ...stack.top(s)...;
    ...
    end
```

OR...

```
with  stack; use stack;
procedure p is
    s : stacktype;
    begin
    ...
    push(4,s);
    ...top(s)...;
    ...
    end
```


Ada Generic -- Abstract Package

```
generic
        -- A private generic type means assignment and equality must be
        -- defined on that type
        type Elem is private ;
package List is
        type T is private ;
        -- Create operation is implicit. Lists created by declaration
        procedure Head  (L: T ; V: out Elem ; Err: out ERROR_INDICATOR) ;
        -- Length can't fail so no need for error indicator
        function Length (L: T) return NATURAL ;
        procedure Tail (L: T ; LT: out T; Err: out ERROR_INDICATOR ) ;
        -- Cons can't fail so no need for error indicator
        function Cons (L: T ; V: Elem ) return T ;
private
        -- an Ada access type corresponds to a Pascal pointer
        -- the entity referenced by the pointer is defined in the package body
        -- In this case, it would be a record with one field pointing to the next
        -- list element
        type LISTREC ;
        type T is access LISTREC ;
end List ;
```

Modules in Modula-2

- **Very similar to Ada packages, but only pointer types can be exported.**

```
Definition module stack;                -- public
    type stacktype;
    procedure empty . . .
    ...
end stack;

Implementation module stack;            -- private
    type stacktype = pointer to record
                        list : . . .
                        topsub: . . .
```


Modula-2 Modules (continued)

- **What are the repercussions of this design decision?**
  - **separate compilation is easy (+)**
  - **module must supply a creation/initialization routine (-)**
  - **extra use of pointers (-)**