

Let  $N \geq 0$ , let  $A$ ,  $B$ , and  $C$  be constants, and let  $f$  and  $g$  be any functions. Then:

$$\sum_{k=1}^N C f(k) = C \sum_{k=1}^N f(k)$$

**S1: factor out constant**

$$\sum_{k=1}^N (f(k) \pm g(k)) = \sum_{k=1}^N f(k) \pm \sum_{k=1}^N g(k)$$

**S2: separate summed terms**

$$\sum_{k=1}^N C = NC$$

**S3: sum of constant**

$$\sum_{k=1}^N k = \frac{N(N+1)}{2}$$

**S4: sum of  $k$**

$$\sum_{k=1}^N k^2 = \frac{N(N+1)(2N+1)}{6}$$

**S5: sum of  $k$  squared**

$$\sum_{k=0}^N 2^k = 2^{N+1} - 1$$

**S6: sum of  $2^k$**

$$\sum_{k=1}^N k 2^{k-1} = (N-1)2^N + 1$$

**S7: sum of  $k 2^{(k-1)}$**

Let  $b$  be a real number,  $b > 0$  and  $b \neq 1$ . Then, for any real number  $x > 0$ , the *logarithm of  $x$  to base  $b$*  is the power to which  $b$  must be raised to yield  $x$ . That is:

$$\log_b(x) = y \text{ if and only if } b^y = x$$

For example:

$$\log_2(64) = 6 \text{ because } 2^6 = 64$$

$$\log_2(1/8) = -3 \text{ because } 2^{-3} = 1/8$$

$$\log_2(1) = 0 \text{ because } 2^0 = 1$$

If the base is omitted, the standard convention in mathematics is that log base 10 is intended; in computer science the standard convention is that log base 2 is intended.

Let  $a$  and  $b$  be real numbers, both positive and neither equal to 1. Let  $x > 0$  and  $y > 0$  be real numbers.

$$\text{L1: } \log_b(1) = 0$$

$$\text{L2: } \log_b(b) = 1$$

$$\text{L3: } \log_b(x) < 0 \text{ for all } 0 < x < 1$$

$$\text{L4: } \log_b(x) > 0 \text{ for all } x > 1$$

$$\text{L5: } \log_b(b^y) = y$$

$$\text{L6: } b^{\log_b(x)} = x$$

$$\text{L7: } \log_b(xy) = \log_b(x) + \log_b(y)$$

$$\text{L8: } \log_b\left(\frac{x}{y}\right) = \log_b(x) - \log_b(y)$$

$$\text{L9: } \log_b(x^y) = y \log_b(x)$$

$$\text{L10: } \log_b(x) = \frac{\log_a(x)}{\log_a(b)}$$

Definition:

Let  $f(x)$  be a function with domain  $(a, b)$  and let  $a < c < b$ . The *limit of  $f(x)$  as  $x$  approaches  $c$  is  $L$*  if, for every positive real number  $\varepsilon$ , there is a positive real number  $\delta$  such that whenever  $|x-c| < \delta$  then  $|f(x) - L| < \varepsilon$ .

The definition being cumbersome, the following theorems on limits are useful. We assume  $f(x)$  is a function with domain as described above and that  $K$  is a constant.

**c1:**  $\lim_{x \rightarrow c} K = K$

**c2:**  $\lim_{x \rightarrow c} x = c$

**c3:**  $\lim_{x \rightarrow c} x^r = c^r$  for all  $r > 0$

Here assume  $f(x)$  and  $g(x)$  are functions with domain as described above and that  $K$  is a constant, and that both the following limits exist (and are finite):

$$\lim_{x \rightarrow c} f(x) = A$$

$$\lim_{x \rightarrow c} g(x) = B$$

Then:

$$\text{c4: } \lim_{x \rightarrow c} Kf(x) = K \lim_{x \rightarrow c} f(x)$$

$$\text{c5: } \lim_{x \rightarrow c} (f(x) \pm g(x)) = \lim_{x \rightarrow c} f(x) \pm \lim_{x \rightarrow c} g(x)$$

$$\text{c6: } \lim_{x \rightarrow c} (f(x) * g(x)) = \lim_{x \rightarrow c} f(x) * \lim_{x \rightarrow c} g(x)$$

$$\text{c7: } \lim_{x \rightarrow c} (f(x) / g(x)) = \lim_{x \rightarrow c} f(x) / \lim_{x \rightarrow c} g(x) \text{ provided } B \neq 0$$

Definition:

Let  $f(x)$  be a function with domain  $[0, \infty)$ . The *limit of  $f(x)$  as  $x$  approaches  $\infty$  is  $L$* , where  $L$  is finite, if, for every positive real number  $\varepsilon$ , there is a positive real number  $N$  such that whenever  $x > N$  then  $|f(x) - L| < \varepsilon$ .

The definition being cumbersome, the following theorems on limits are useful. We assume  $f(x)$  is a function with domain  $[0, \infty)$  and that  $K$  is a constant.

**c8:** 
$$\lim_{x \rightarrow \infty} K = K$$

**c9:** 
$$\lim_{x \rightarrow \infty} \frac{1}{x} = 0$$

**c10:** 
$$\lim_{x \rightarrow \infty} \frac{1}{x^r} = 0 \text{ for all } r > 0$$

Given a rational function the last two rules are sufficient if a little algebra is employed:

$$\lim_{x \rightarrow \infty} \frac{7x^2 + 5x + 10}{3x^2 + 2x + 5} = \lim_{x \rightarrow \infty} \frac{7 + \frac{5}{x} + \frac{10}{x^2}}{3 + \frac{2}{x} + \frac{5}{x^2}}$$

**Divide by highest power of  $x$  from the denominator.**

$$= \frac{\lim_{x \rightarrow \infty} 7 + \lim_{x \rightarrow \infty} \frac{5}{x} + \lim_{x \rightarrow \infty} \frac{10}{x^2}}{\lim_{x \rightarrow \infty} 3 + \lim_{x \rightarrow \infty} \frac{2}{x} + \lim_{x \rightarrow \infty} \frac{5}{x^2}}$$

**Take limits term by term.**

$$= \frac{7 + 0 + 0}{3 + 0 + 0}$$

**Apply theorem C3.**

$$= \frac{7 + 0 + 0}{3 + 0 + 0}$$

$$= \frac{7}{3}$$

In some cases, the limit may be infinite. Mathematically, this means that the limit does not exist.

$$\mathbf{c11:} \quad \lim_{x \rightarrow \infty} x^r = \infty \text{ for all } r > 0$$

$$\mathbf{c13:} \quad \lim_{x \rightarrow \infty} (e^x) = \infty$$

$$\mathbf{c12:} \quad \lim_{x \rightarrow \infty} (\log_b x) = \infty$$

Example:

$$\begin{aligned} \lim_{x \rightarrow \infty} \frac{7x^2 + 5x + 10}{2x + 5} &= \lim_{x \rightarrow \infty} \frac{7x + 5 + \frac{10}{x}}{2 + \frac{5}{x}} \\ &= \frac{\lim_{x \rightarrow \infty} 7x + \lim_{x \rightarrow \infty} 5 + \lim_{x \rightarrow \infty} \frac{10}{x}}{\lim_{x \rightarrow \infty} 2 + \lim_{x \rightarrow \infty} \frac{5}{x}} = \infty \end{aligned}$$

In some cases, the reduction trick shown for rational functions does not apply:

$$\lim_{x \rightarrow \infty} \frac{7x + 5 \log(x) + 10}{2x + 5} = ??$$

In such cases, l'Hôpital's Rule is often useful. If  $f(x)$  and  $g(x)$  are differentiable functions such that

$$\lim_{x \rightarrow c} f(x) = \lim_{x \rightarrow c} g(x) = \infty$$

then:

$$\lim_{x \rightarrow c} \frac{f(x)}{g(x)} = \lim_{x \rightarrow c} \frac{f'(x)}{g'(x)}$$

**This also applies if the limit is 0.**

Applying L'Hôpital's Rule:

$$\lim_{x \rightarrow \infty} \frac{7x + 5 \log(x) + 10}{2x + 5} = \lim_{x \rightarrow \infty} \frac{7 + \frac{5}{x}}{2} = \frac{7}{2}$$

Another example:

$$\lim_{x \rightarrow \infty} \frac{x^3 + 10}{e^x} = \lim_{x \rightarrow \infty} \frac{3x^2}{e^x} = \lim_{x \rightarrow \infty} \frac{6x}{e^x} = \lim_{x \rightarrow \infty} \frac{6}{e^x} = 0$$

Recall that:

$$D[e^{f(x)}] = e^{f(x)} D[f(x)]$$

Mathematical induction is a technique for proving that a statement is true for all integers in the range from  $N_0$  to  $\infty$ , where  $N_0$  is typically 0 or 1.

## First (or Weak) Principle of Mathematical Induction

Let  $P(N)$  be a proposition regarding the integer  $N$ , and let  $S$  be the set of all integers  $k$  for which  $P(k)$  is true. If

- 1)  $N_0$  is in  $S$ , and
- 2) whenever  $N$  is in  $S$  then  $N+1$  is also in  $S$ ,

then  $S$  contains all integers in the range  $[N_0, \infty)$ .

To apply the PMI, we must first establish that a specific integer,  $N_0$ , is in  $S$  (establishing the basis) and then we must establish that if a arbitrary integer,  $N \geq N_0$ , is in  $S$  then its successor,  $N+1$ , is also in  $S$ .

Theorem: For all integers  $n \geq 1$ ,  $n^2+n$  is a multiple of 2.

*proof:* Let  $S$  be the set of all integers for which  $n^2+n$  is a multiple of 2.

If  $n = 1$ , then  $n^2+n = 2$ , which is obviously a multiple of 2. This establishes the basis, that 1 is in  $S$ .

Now suppose that some integer  $k \geq 1$  is an element of  $S$ . Then  $k^2+k$  is a multiple of 2. We need to show that  $k+1$  is an element of  $S$ ; in other words, we must show that  $(k+1)^2+(k+1)$  is a multiple of 2. Performing simple algebra:

$$(k+1)^2+(k+1) = (k^2 + 2k + 1) + (k + 1) = k^2 + 3k + 2$$

Now we know  $k^2+k$  is a multiple of 2, and the expression above can be grouped to show:

$$(k+1)^2+(k+1) = (k^2 + k) + (2k + 2) = (k^2 + k) + 2(k + 1)$$

The last expression is the sum of two multiples of 2, so it's also a multiple of 2. Therefore,  $k+1$  is an element of  $S$ .

Therefore, by PMI,  $S$  contains all integers  $[1, \infty)$ .

QED

Theorem: Every integer greater than 3 can be written as a sum of 2's and 5's.

(That is, if  $N > 3$ , then there are nonnegative integers  $x$  and  $y$  such that  $N = 2x + 5y$ .)

This is not (easily) provable using the First Principle of Induction. The problem is that the way to write  $N+1$  in terms of 2's and 5's has little to do with the way  $N$  is written in terms of 2's and 5's. For example, if we know that

$$N = 2x + 5y$$

we can say that

$$N + 1 = 2x + 5y + 1 = 2x + 5(y - 1) + 5 + 1 = 2(x + 3) + 5(y - 1)$$

but we have no reason to believe that  $y - 1$  is nonnegative. (Suppose for example that  $N$  is 9.)

There is a second statement of induction, sometimes called the "strong" form, that is adequate to prove the result on the preceding slide:

## Second (or Strong) Principle of Mathematical Induction

Let  $P(N)$  be a proposition regarding the integer  $N$ , and let  $S$  be the set of all integers  $k$  for which  $P(k)$  is true. If

- 1)  $N_0$  is in  $S$ , and
- 2) whenever  $N_0$  through  $N$  are in  $S$  then  $N+1$  is also in  $S$ ,

then  $S$  contains all integers in the range  $[N_0, \infty)$ .

Interestingly, the "strong" form of induction is logically equivalent to the "weak" form stated earlier; so in principle, anything that can be proved using the "strong" form can also be proved using the "weak" form.

Theorem: Every integer greater than 3 can be written as a sum of 2's and 5's.

*proof:* Let  $S$  be the set of all integers  $n > 3$  for which  $n = 2x + 5y$  for some nonnegative integers  $x$  and  $y$ .

If  $n = 4$ , then  $n = 2*2 + 5*0$ . If  $n = 5$ , then  $n = 2*0 + 5*1$ . This establishes the basis, that 4 and 5 are in  $S$ .

Now suppose that all integers from 4 through  $k$  are elements of  $S$ , where  $k \geq 5$ . We need to show that  $k+1$  is an element of  $S$ ; in other words, we must show that  $k+1 = 2r + 5s$  for some nonnegative integers  $r$  and  $s$ .

Now  $k+1 \geq 6$ , so  $k-1 \geq 4$ . Therefore by our assumption,  $k-1 = 2x + 5y$  for some nonnegative integers  $x$  and  $y$ . Then, simple algebra yields that:

$$k+1 = k-1 + 2 = 2x + 5y + 2 = 2(x+1) + 5y,$$

whence  $k+1$  is an element of  $S$ .

Therefore, by the Second PMI,  $S$  contains all integers  $[4, \infty)$ .

**QED**

Consider the sequences  $\{a_k\}$  and  $\{d_k\}$  :

$$a_k = 2^k - 1 \text{ for } k \geq 0$$

$$\begin{cases} d_0 = 0, d_1 = 1 \\ d_k = 3d_{k-1} - 2d_{k-2} \text{ for } k \geq 2 \end{cases}$$

Then for all  $k \geq 0$ .  $a_k = d_k$ .

*proof:* Let  $S = \{k \geq 0 \mid a_k = d_k\}$ .

Trivial calculations show that  $a_0 = d_0$  and  $a_1 = d_1$ , so 0 and 1 are in S.

Suppose that there is some  $N \geq 1$  such that 0, 1, ..., N are in S. In other words, assume that there is some  $N \geq 1$  such that  $a_i = d_i$  for all i from 0 to N.

Now,  $N + 1 \geq 2$ , so from the definition of  $\{d_k\}$  and the inductive assumption we have:

$$\begin{aligned} d_{N+1} &= 3d_N - 2d_{N-1} \\ &= 3(2^N - 1) - 2(2^{N-1} - 1) \\ &= 3 \cdot 2^N - 3 - 2^N + 2 \\ &= 2 \cdot 2^N - 1 \\ &= 2^{N+1} - 1 \end{aligned}$$

Therefore  $N + 1$  is in S, and so by the principle of induction,  $S = \{k \geq 0\}$ . **QED**