

Implement complex code, such as a PR quadtree, piece-meal.

Write code one small chunk at a time and then test it.

The *code chunks* should implement a logical sub-part of a method/operation.

Carefully consider all of the cases of the operations, e.g. search, insertion and deletion.

Logically order the cases that must handled.

Operation Implementation:

- Implement the first case and test it.
- Implement the second case, test it and then test any interaction with the first case.
- Implement the third case, test it and then test any interaction with the previous cases. etc. etc.
- After all cases are implemented, test the interaction of the operation with other operations.

Search logic in data structures tends to be used in multiple operations, e.g. find, delete, insert). Thus implementing and testing the search logic first will aid in the other operation development.

What are the search cases to be considered?

Tree is empty.

Select branch to search down.

Leaf node comparison.

Item to be found not in tree.

Item to be found not in World.

Item to be found does not exist (null).

Re-order cases for logical implementation:

1. Item to be found does not exist (null parameter).
2. Tree is empty.
3. Item to be found is not in the World.
4. Leaf node comparison: item found or not in tree.
5. Inner node: select branch to search down:
 - a. Compute quadrant division midpoint.
 - b. Determine quadrant in which element to be found is located.

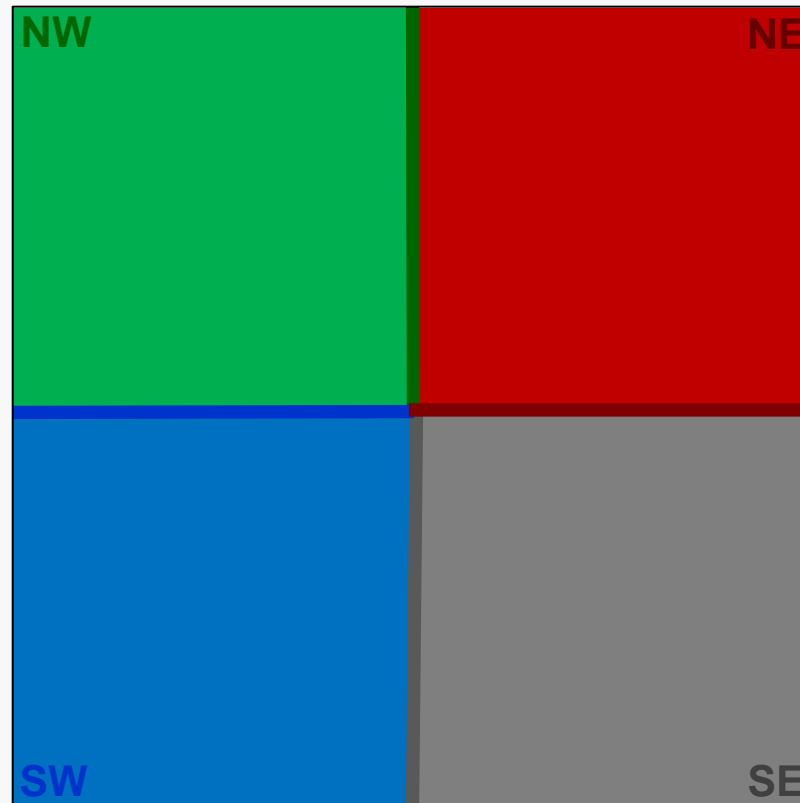
Code one case, then test.

In which quadrant does a point belong when it is on the boundary line of two adjacent quadrants? In which quadrant should the intersection point of four quadrants be placed?

Conventions:

Samet: “...the lower and left boundaries for each block are closed, while the upper and right boundaries for each block are open.”

3114: points on a boundary are considered to belong to its counter-clockwise quadrant and the origin point belongs to the NE quadrant.



Cases:

1. Item to be found does not exist (null).
2. Item to be found not within World.
3. Tree is empty.
4. Current node is a Leaf node.
5. Current node is an Inner node

Case #1: nothing to do.

Case #2: nothing to do.

Case #3: instantiate new leaf node with element to add & attach.

This discussion
assumes one
element per leaf.

Case #4: Current node is a Leaf node.

- a. Coordinates of leaf node element and element to add are equal.
Operation is either an update or non-allowed duplicate.
- b. Coordinates of leaf node element and element to add are NOT equal.
 - i. Determine midpoint of quadrant.
 - ii. Instantiate new inner node, creating sub-quadrant references.
 - iii. Determine in which sub-quadrant existing leaf element falls and assign it to reference the existing leaf node.
 - iv. Determine in which sub-quadrant (SQ) element to add falls
 - 1) SQ is empty (null): instantiate new leaf node with element to add & attach.
 - 2) SQ is NOT empty: descend recursively ...
(existing leaf elem & elem to add falls in the same SQ)



Case #5: Current node is an Inner node.

- a. Determine in which sub-quadrant (SQ) element to add falls.
- b. Descend recursively, update SQ reference upon ascent.

Cases:

1. Item to be found does not exist (null parameter).
2. Item to be found not within World.
3. Tree is empty.
4. Current node is a Leaf node.
5. Current node is an Inner node

Case #1: nothing to do.

Case #2: nothing to do.

Case #3: return unsuccessful deletion.

This discussion
assumes one
element per leaf.

Case #4: Current node is a Leaf node.

- a. Coordinates of leaf node element and element to delete are NOT equal.
Indicate element NOT found in tree, back out.
- b. Coordinates of leaf node element and element to delete are equal.
Remove the element (set the reference in the node to the element to null).
Return null for recursive ascent parent reference assignment.

Case #5: Current node is an Inner node.

- a. Determine in which sub-quadrant (SQ) element to delete falls.
- b. Descend recursively, update SQ reference upon ascent.
- c. Tree contraction?
 - i. if elem was removed lower in tree & SQ now contains one child:
Contraction: set SQ reference to child to null & return reference to child.
 - ii. if elem was NOT removed lower in tree or SQ now contains $>$ one child:
No contraction, return inner node reference for ascent update.

Searching:

- Average: $O\left(\left\lceil \log\left(\frac{\sqrt{2}s}{d}\right) \right\rceil\right)$

where:

s length of side of world
 d minimum distance
 between any two points
 in tree

Assuming a PR Quadtree of size N , built using random insertion of elements.

Insertion:

- Requires search to locate element or its insertion leaf sub-quadrant plus cost to create new inner node and redistribute existing leaf + added element. For random insertion, collision of existing leaf & new element will occur infrequently. This cost of redistribution on average will be a constant.
- Average: $O\left(\left\lceil \log\left(\frac{\sqrt{2}s}{d}\right) \right\rceil\right)$

Deletion

- Requires search to locate element plus cost of possible contraction during ascent. For randomly formed tree, contraction will occur infrequently. This cost of contraction on average will be a constant.
- Average: $O\left(\left\lceil \log\left(\frac{\sqrt{2}s}{d}\right) \right\rceil\right)$