



Computer Chess © 1985-2015, D Barnette

# Board Representation



8 x 8 Matrix

• Problem: doubles subscript computation

#### Array (64)

- Subscript range checking: every move generated must be checked
- Illegal moves easily generated: example King on square 8, adding 1 generates move to square 9

Computer Chess © 1985-2015, D Barne

# Board Representation: Array 120



#### Illegal Move Detection Assume Machine plays White: CASE destination square value

- 7 : Off Board Move
- 1...6 : Square occupied by Machine's piece
  - 0 : open square
  - < 0 : capturing move

#### **Move Generation**

- Routine required for each type of piece
- Piece location passed as parameter

#### 111 112 113 114 115 116 117 118 119 120



Computer Chess © 1985-2015, D Barnette

## Board Representation: Array 120 Move Generation









Kings, Knights, pawns move number applied only once.

Queens, Bishops, Rooks move number applied repeatedly until off board square or occupied square generated.



## Bit-Board Representation

- 64 Bits, 1 per square, used to store pieces' locations
- Bit value:
  - 1 = represents a piece
  - 0 = empty square
- 12 x 64 bits to store a position
  - 6 white pieces
  - 6 black pieces
  - 64 bits for each different type of piece: King, Queen, Bishop, Knight, Rook, Pawn
- 2 extra words are used to store all piece locations for both sides for efficient move generation



## Bit-Board Move Generation

- For each type of piece all possible moves for every square must be stored: 6 x 64 x 64 = 3K bytes
- Move bit-maps are searched for the piece to be moved & the square piece is on to locate piece move bit-map
- Logical operations are performed to generate legal bit-map moves:

## & (piece move bit-map,~(bit-map ∀ pieces of moving side))

• Generates bit-map for all legal piece moves in a given position



# Bit-Map Move Generation







## Evaluation Function

Numerical value assigned to a given position Programs next move decision is based *entirely* on Fn Tradeoff:

Complex Fn	more computing time less search time
Simple Fn	less computing time more search time

### Material Evaluation

Values:	Pawn	Knight	Bishop	Rook	Queen	King
	100	300	300	500	900	$\infty$

```
//all computation performed with integer operations
FUNCTION Chess ( P , N , B , R , Q ) : INTEGER ;
Chess = P * 100 + N * 300 + B * 300 + R * 500 + Q * 900
where:
P = # of m/c's pawns - # of opponent's pawns
N = # of m/c's knights - # of opponent's knights
```



## Evaluation Function Positional Factors

# (game heuristics)

Mobility

bonus for each square attacked by a piece

Pawn Structure

bonus for passed pawns, central pawns penalty for doubled pawns, isolated pawns

Attacks

bonus for attacks against opponent's pieces

Castling

bonus for completion

Development

penalty for unmoved pieces

**Doubled pieces** 

bonus for pieces attacking/defending same square

King Safety

penalty for enemy attacks

bonus for pieces near King (piece tropism)



Position Score = 0

Computer Chess © 1985-2015, D Ba

Symmetrical Positions = 0









Can a computer play the perfect game of chess?

Can the entire game tree be formed & searched?

Chess 50 move Rule:

If 50 moves occur without a capture or pawn move the game is a draw.

[*IJ Good 1969*] (Va Tech Emeritus University Distinguished Professor)

- Maximum length of a game is 3150 moves.
- Maximum number of possible moves in a game = 320
- Theoretical maximum number of possible chess games =

 $320^{6300} \cong 10^{15790}$ 





Average middle game position:

- Approximately 30 legal moves
- Tree of depth 4 (2 moves)
  - 10<sup>6</sup> leaves (positions)
  - 10<sup>4</sup> positions for move

## Assume

- Scoring = 10  $\mu$ secs & move generation = 1 $\mu$ sec
- A move for tree of depth 4 would take 40 secs
- A tree of 6 ply (3 moves) would take 11 hours





Tree Size Control Methods: Width

- *Width Control*: number of moves considered per position
- Fanout Parameters
  - Limits the number of moves considered at each level.
  - Decreased with depth 0 [25,30] . . . depth 10 [5,10]
- Plausibility Ordering (iterative deepening) *plausibility* score: heuristic value that determines the worth of a move being searched.
- assigns a *plausibility* score to each generated move
- moves ordered by the *plausibility* score for application of the depth fanout parameters





Tree Size Control Methods: Depth

- Depth Control: number of moves a line (tree path) is searched
- Fixed Depth  $\rightarrow$  Horizon Effect Problem
  - Key move is missed at the next immediate depth.
  - Example: previous position & 2-ply tree program would choose the highest score = S<sub>41</sub> (Qc1) missing the 2 move mate.
- Quiescence
  - Position must be static (quiescent: no captures, checks, =) before analysis can be stopped & score function applied.
- Depth Cutoff Lists
  - Lists containing positional features to determine quiescence.
  - Number of features decreases with depth.
  - When applied to a position if no feature is found position is considered static, otherwise it is explored further.



# Game Theory: Min Max Method

Technique of selecting the best move in a tree

Assumes best play by opponent

- High score is best possible move (position) for program.
- Low score is best possible move for opponent.

Evaluation Fn is applied to static (quiescent) positions

- Minimum or Maximum (depending upon player's/opponent's move) of scores for lower level nodes are selected assigned to the higher level nodes.
- Process repeats until root (current position) is reached.
- The move that is selected is the one that leads to the backed-up value for the root.



## Game Theory: Min Max Method



Thus  $M_2$  is the selected move.

Computer Chess © 1985-



and searching of positions

Alpha-cutoff

Computer Chess © 1985-

Based upon Move Refutation strategy

After a move has been searched & scored (tree path), when other moves are searched if they lead to positions that allow the opponent to force a lower score then the program need not explore the moves.

#### E **Alpha-Beta Algorithm** Tree pruning method Po M<sub>1</sub> $M_2$ P2 M<sub>13</sub> M11 M12 M21 M22 M23 P22 P21 P23 P<sub>11</sub> S22 S21 S23 Alpha-cutoff $S_1 = Min(S_{11}, S_{12}, S_{13}, \equiv) = 27$ $S_{21} = 22$ Since $S_{21} < S_1$ positions $P_{22}$ , $P_{23} \equiv$ need not be evaluated.

#### Reasoning:

IF S<sub>22</sub>, S<sub>23</sub>,  $\equiv$  > S<sub>21</sub> THEN opponent would not choose them (*min*) IF S<sub>22</sub>, S<sub>23</sub>,  $\equiv$  < S<sub>21</sub> THEN S<sub>2</sub> < S<sub>1</sub> & m/c will not choose M<sub>2</sub> (*max*)

Computer Chess © 1985-2015, D Barnette

# Alpha-Beta Algorithm Tree pruning method

<u>Alpha</u> highest score currently achieved for the machine

<u>Beta</u> highest score (negative) currently achieved for the opponent

Alpha Cutoffs

 Occur at odd depths when opponent's move could reduce program's score

Beta Cutoffs

- Occur at even depths when program could reduce the "best" score achieved for the opponent
- $\alpha$ – $\beta$  pruning reduces the search time by 80%

 $\rightarrow$  15 minute move is reduced to a 3 minute move.

 $\alpha-\beta$  pruning combined with iterative deepening (plausibility scoring) achieves a total reduction of 90%

 $\rightarrow$  30 minute move is reduced to a 3 minute move.





## YouTube: Computer Chess

- <u>The History of Chess Computers</u>
- Alan Turing The first ever Chess program
- <u>Chess Computer Science Artifical Intelligence Paper: How</u> <u>do Chess Engines work?</u>
- Nova: Kasparov versus Deep Thought
  - <u>Part 1</u>
  - <u>Part 2</u>
  - <u>Part 3</u>
  - <u>Part 4</u>
- Deep Blue beat G. Kasparov in 1997
- Chess Programming on the Web
  - <u>chessprogramming.wikispaces.com</u>

