## C Programming

Pointer Accesses to Memory and Bitwise Manipulation
This assignment consists of implementing a function that can be executed in two modes, controlled by a switch specified by a parameter to the function:

```
enum _DataFormat {CLEAR, ENCRYPTED};
typedef enum _DataFormat DataFormat;
struct _WordRecord {
    uint16_t offset; // offset at which word record was found in memory
    char* word; // dynamically alloc'd C-string containing the "word"
};
typedef struct _WordRecord WordRecord;
/**
    * Untangle() parses a chain of records stored in the memory region pointed
    * to by pBuffer, and stores WordRecord objects representing the given data
    * into the array supplied by the caller.
    * Pre: Fmt == CLEAR or ENCRYPTED
    * pBuffer points to a region of memory formatted as specified
    * wordList points to an empty array large enough to hold all the
        WordRecord objects you'll need to create
    * Post: wordList[0:nWords-1] hold WordRecord objects, where nWords is
                            is the value returned by Untangle()
    * Returns: the number of "words" found in the supplied quotation.
    */
uint8_t Untangle(DataFormat Fmt, const uint8_t* pBuffer, WordRecord* const wordlist);
```

The function will access a scrambled quotation, stored in a memory region pointed to by pBuffer. The organization of the memory region is described in detail below. The function will analyze that memory region, and reconstruct the quotation by created a sequence of WordRecord objects and storing them in an array provided by the caller.

You will also implement a function that will deallocate all the dynamic content of such an array of WordRecord objects:

```
/**
    * Deallocates an array of WordRecord objects.
    * Pre: wordList points to a dynamically-allocated array holding nWords
    * WordRecord objects
    * Post: all dynamic memory related to the array has been freed
    */
void clearWordRecords(WordRecord* const wordList, uint8_t nWords);
```

Part of your score on the assignment will depend on the correctness of this function, and your ability to deallocate any other allocations your solution may perform. This will be determined by running your solution on Valgrind; your goal is to achieve a Valgrind report showing no memory leaks or other memory-related errors:

```
==10833== Memcheck, a memory error detector
==10833== Copyright (C) 2002-2013, and GNU GPL'd, by Julian Seward et al.
==10833== Using Valgrind-3.10.0 and LibVEX; rerun with -h for copyright info
==10833== Command: ./driver -clear
==10833== Parent PID: 10832
==10833==
==10833== HEAP SUMMARY:
==10833== in use at exit: 0 bytes in 0 blocks
==10833== total heap usage: 230 allocs, 230 frees, 13,642 bytes allocated
==10833==
==10833== All heap blocks were freed -- no leaks are possible
==10833==
==10833== ERROR SUMMARY: 0 errors from 0 contexts (suppressed: 1 from 1)
```


## Case 1 [80\%]

## Untangling Clear Data Records in Memory

Here is a hexdump of a memory region containing a scrambled quotation:

|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | F |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00000000 | 34 | 00 | $0 f$ | 3 a | 00 | 69 | 6e | 64 | 69 | 66 | 66 | 65 | 72 | 65 | 6 C | 63 | \|4..:.indifferenc| |
| 00000010 | 65 | 05 | 3f | 00 | 62 | 65 | 05 | 47 | 00 | 62 | 79 | 06 | 02 | 00 | 66 | $6 f$ | \|e.?.be.G.by...fo| |
| 00000020 | 72 | 0a | 68 | 00 | 70 | 65 | 6 e | 61 | 6c | 74 | 79 | 09 | 74 | 00 | 70 | 75 | \|r.h.penalty.t.pu| |
| 00000030 | 62 | 6c | 69 | 63 | 06 | 21 | 00 | 54 | 68 | 65 | 05 | 2b | 00 | 74 | $6 f$ | 08 | \|blic.!.The.+.to. |
| 00000040 | 16 | 00 | 72 | 75 | 6c | 65 | 64 | 07 | 4 e | 00 | 65 | 76 | 69 | 6c | 07 | 55 | ..ruled.N.evil.U\| |
| 00000050 | 00 | 6d | 65 | 6 e | 2e | 05 | 5a | 00 | 2d | 2d | 08 | 00 | 00 | 50 | 6c | 61 | .men..z.--...Pla\| |
| 00000060 | 74 | 67 | 06 | 83 | 00 | 6d | 65 | 6 e | 07 | 62 | 00 | 67 | 6 f | 6f | 64 | 05 | \|to...men.b.good. |
| 00000070 | 11 | 00 | 74 | 67 | 0a | 7 e | 00 | 61 | 66 | 66 | 61 | 69 | 72 | 73 | 05 | $6 f$ | \|..to.~.affairs.o| |
| 00000080 | 00 | 69 | 73 | 06 | 1b | 00 | 70 | 61 | 79 |  |  |  |  |  |  |  | \|.is...pay |

$$
\begin{array}{lllllllllllllllll}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & A & B & C & D & E & F
\end{array}
$$

The first two bytes of the memory region contain the offset at which you will begin processing records: $0 \times 0034$.
This offset of the first record is followed by a sequence of word records, each consisting of a positive integer value, another positive integer value, and a sequence of characters:

| Length of record | Offset of next record | Characters in word |
| :---: | :---: | :---: |
| uint8_t | uint16_t | chars |

The first value in each record specifies the total number of bytes in the record. Since words are relatively short, the value will be stored as a uint8_t value, which has a range of $0-255$. The record length is followed immediately by a uint16_t value specifying the offset of the next word record in the list. This is followed by a sequence of ASCII codes for the characters that make up the word. (The term "word" is used a bit loosely here.) There is no terminator after the final character of the string, so be careful about that.

Note that the length of the record depends upon the number of characters in the word, and so these records vary in length. That's one reason we must store the offset for each record.

Since I'm using x86 hardware, integer values are stored in memory in little-endian order; that is, the low-order byte is stored first (at the smallest address) and the high-order byte is stored last (at the largest address). So the bytes of a multibyte integer value appear to be reversed. For example, if we have in an int32_t variable the base-10 value 85147, the corresponding base-16 representation world be 0x14C9B, and the in-memory representation would look like this:

| 9B | 4C | 01 | 00 |
| :---: | :---: | :---: | :---: |
| low | high |  |  |

or, represented in pure binary:

| 10011011 | 01001100 | 00000001 | 00000000 |
| :---: | :---: | :---: | :---: |
| low |  |  |  |

The least-significant byte (corresponding to the lowest powers of 2) is stored at the lowest address, and the most-significant byte (corresponding to the highest powers of 2 ) is stored at the highest address.

As a programmer, you usually do not need to take the byte-ordering into account since the compiler will generate machine language compatible with your hardware, and that will make use of the bytes in the appropriate manner. But, when you're reading memory displays, you must take the byte-ordering into account.

So, looking at the first two bytes of the memory block, we see that the word record we will process first occurs at relative offset $0 \times 0034$ from the beginning of the memory block.

Let's consider how to interpret the hexdump shown earlier:


$$
\begin{array}{llllllllllllllll}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & A & B & C & D & E & F
\end{array}
$$

The first word record consists of the bytes:

$$
\begin{array}{llllll}
06 & 21 & 00 & 54 & 68 & 65
\end{array}
$$

The length of the first record is $0 \times 06$ or 6 in base- 10 , which means that the string is 3 characters long, since the length field occupies 1 byte and the offset of the next record occupies 2 bytes. The ASCII codes are 546865 , which represent the characters "The". The offset of the next record is $0 \times 0021$.

The second word record consists of the bytes:
$0 a 680070656 e 616 c 7479$
The length is $0 \times 0 \mathrm{a}$ (10 in base-10), so the string is 7 characters long (the ASCII codes represent "penalty"), and the next word record is at the offset $0 \times 0068$. And so forth... The complete quotation, with word record offsets, is:

| 0x0037: | The |
| :--- | :--- |
| $0 \times 0024:$ | penalty |
| $0 \times 006 \mathrm{~B}:$ | good |
| $0 \times 0065:$ | men |
| $0 \times 0086:$ | pay |
| $0 \times 001 \mathrm{E}:$ | for |
| $0 \times 0005:$ | indifference |
| $0 \times 003 \mathrm{D}:$ | to |
| $0 \times 002 \mathrm{E}:$ | public |
| $0 \times 0077:$ | affairs |
| $0 \times 0081:$ | is |
| $0 \times 0072:$ | to |
| $0 \times 0014:$ | be |
| $0 \times 0042:$ | ruled |
| $0 \times 0019:$ | by |
| $0 \times 004 \mathrm{~A}:$ | evil |
| $0 \times 0051:$ | men. |
| $0 \times 0058:$ | -- |
| $0 x 005 D:$ | Plato |

To indicate the end of the sequence of word records, the final word record specifies that its successor is at an offset of 0 , which is invalid (since that's the offset of the pointer to the first word record).

For case 1, your function will be passed the value CLEAR for the parameter Fmt.

You will use the WordRecord data type to represent a parsed word record:

```
struct _WordRecord {
    uint16_t offset; // offset at which word record was found in memory
    char* word; // dynamically alloc'd C-string containing the "word"
};
typedef struct _WordRecord WordRecord;
```

You will create one of these struct variables whenever you parse a word record, and place that struct variable into an array supplied by the caller of your function.

## Case 2 [20\%]

## Untangling Mildly Encrypted Data Records in Memory

Read the posted notes on bitwise operations in C, and the related sections in your C reference text.
For this case, your function will be passed the value ENCRYPTED for the parameter Fmt.
The memory region pointed to by pBuffer will be formatted in exactly the same way as for case 1 , except that the bytes that represent the offset of the next record and the characters in the word will have been "masked":

| Length of record | Masked offset of next record | Masked characters in word |
| :---: | :---: | :---: |
| uint8_t | uint16_t | chars |

Each of the ASCII codes in the word field has been XORed with a mask formed by taking the number of characters (bytes) in the word, and reversing the nybbles of that value (remember, the length of the character sequence is stored as a one-byte value). Each byte of the offset field has been XORed with the unmasked first byte in the word. You must "unmask" the masked bytes in order to properly display the quotation.

Part of the assignment is for you to determine what operation(s) you can use to perform this unmasking. We will not answer any questions about how to do that, except to say that you should consider the properties of the various bitwise operations available in C. This is a good opportunity for you to discover the value of the Boolean algebra rules covered in Discrete Mathematics.

Aside from the issue of unmasking the encrypted bytes, the logic of this part is identical to the handling without encryption, so we will not repeat the detailed description given there. However, we will give you an example illustrating what must be done:


$$
\begin{array}{llllllllllllllll}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & A & B & C & D & E & F
\end{array}
$$

The first word record begins at offset $0 \times 0025$, and contains the bytes: 0576496954
The length of the word is $0 \times 02$, so the mask is $0 \times 20$. XORing that with each byte of the string yields 4974 , which represents the character string "It".

And, XORing the first byte of the word with the bytes of the offset yields 3 F 00 , which is the offset of the next record. In this case, the encrypted quotation decodes to:

| 0x0028: | It |
| :--- | :--- |
| 0x0042: | is |
| 0x0069: | the |
| 0x0056: | mark |
| 0x0079: | of |
| 0x0047: | an |
| 0x0083: | educated |
| 0x008E: | mind |
| 0x002D: | to |
| 0x001D: | be |
| 0x0032: | able |
| 0x007E: | to |
| 0x0011: | entertain |
| $0 \times 003 E:$ | a |
| $0 \times 004 C:$ | thought |
| $0 \times 006 F:$ | without |
| $0 \times 0005:$ | accepting |
| $0 \times 0022:$ | it. |
| $0 \times 0039:$ | -- |
| $0 \times 005 D:$ | Aristotle |

## Testing and Grading

A tar file is posted for the assignment containing testing/grading code:

| gradec05.sh | Bash script to perform automated testing; see the header comment for instructions. <br> Note that you cannot use the shell script for debugging purposes; you must execute your <br> program directly instead. |
| :---: | :--- |
| C05Grader.tar: |  |
| driver.c | Test driver |
| Untangle.h | Declarations for specified function <br> checkAnswer.h |
| Declarations for answer-checking and grading function |  |
| checkAnswer.o | 64-bit Linux binary for checking/grading code |
| Generator.h | Declarations for test data generator <br> Generator.o |
| 64-bit Linux binary for test data generator |  |

Create Untangle. c and implement your version of it, then compile it with the files above. Read the header comments in driver.c for instructions on using it.

The test driver produces two output files:

| Log.txt | shows your output, correct output, and score information |
| :--- | :--- |
| Data.bin | binary file containing same bytes as the memory region used in the test |

The second file cannot be viewed in most text editors. However, you can use the hexdump utility to display the contents in a form that's almost identical to the examples shown earlier in this specification: hexdump -C Data.bin

The grading script automates the entire process of running the tests, including Valgrind. If Valgrind detects any errors at all when your solution is tested, we will assess a penalty of $10 \%$ of the project score. We will also assess the penalty if your solution does not allocate the char arrays for the words dynamically, or if your solution allocates arrays that are larger than necessary.

## What to Submit

You will submit your file Unt angle.c to Canvas. That file must include any helper functions you have written and called from your version of Untangle( ); any such functions must be declared (as static) is the file you submit. You must not include any extraneous code (such as implementations of main( ) in that file).

Your submission will be graded by running the supplied test/grading code on it.

If you work with a partner, make sure that the submitted file contains a properly-completed copy of the partners form posted on the assignments page. Failure to do that will result in at least one of you not receiving credit for the assignment.

## Pledge:

Each of your program submissions must be pledged to conform to the Honor Code requirements for this course.
Specifically, you must include the following pledge statement in the submitted file:

```
// On my honor:
//
// - I have not discussed the C language code in my program with
// anyone other than my instructor or the teaching assistants
// assigned to this course.
//
// - I have not used C language code obtained from another student,
// the Internet, or any other unauthorized source, either modified
// or unmodified.
//
// - If any C language code or documentation used in my program
// was obtained from an authorized source, such as a text book or
// course notes, that has been clearly noted with a proper citation
// in the comments of my program.
//
// - I have not designed this program in such a way as to defeat or
// interfere with the normal operation of the grading code.
//
// <Student Name>
// <Student's VT email PID>
```


## We reserve the option of assigning a score of zero to any submission that is undocumented or does not contain this statement.

## Change Log

Any changes or corrections to the specification will be documented here.

| Version | Posted | Pg | Change <br> Base document. |
| :--- | :--- | ---: | :--- |
| 1.00 | June 13 |  |  |

