

Introduction

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What is a network?

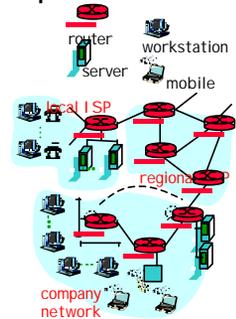
- Carrier of information between connected entities
- What does a network consist of?
 - End hosts connected to the network
 - Routers/switches that move data through the network
 - Physical links that carry information
 - E.g. Ethernet, FDDI, ATM, Token Ring
 - Applications that communicate with each other to provide services
 - E-Mail, File Transfer, Web Browser

What is an internetwork?

- A set of connected networks is called an internetwork.
- The **Internet** is a specific example of an internetwork.
 - The Internet is characterized by the use of a common network protocol (IP).
- An internetwork need **NOT** be connected to the **Internet**

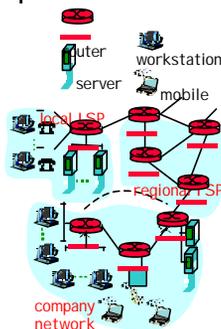
Internet: Components

- **Network edge:**
 - millions of connected computing devices running **network apps**
 - pc's workstations, servers
 - PDA's phones, toasters
- **Network core:**
 - **routers:** forward packets (chunks) of data thru network
- **Media:**
 - **communication links:** fiber, copper, radio, satellite



Internet: Components

- **protocols:** control sending, receiving of msgs
 - e.g., TCP, IP, HTTP, FTP, PPP
- **Internet: "network of networks"**
 - loosely hierarchical
 - public Internet versus private intranet
- **Internet standards**
 - RFC: Request for comments
 - IETF: Internet Engineering Task Force



Internet: Services

- **communication infrastructure** enables distributed applications:
 - WWW, email, games, e-commerce, database., voting,
 - more?
- **communication services provided:**
 - connectionless
 - connection-oriented
- **cyberspace [Gibson]:**
 - "a consensual hallucination experienced daily by billions of operators, in every nation,"

What's a protocol?

human protocols:

- "what's the time?"
- "I have a question"
- introductions

... specific msgs sent

... specific actions taken when msgs received, or other events

network protocols:

- machines rather than humans
- all communication activity in Internet governed by protocols

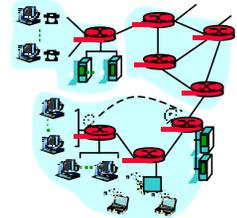
protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt

A closer look at network structure

- **network edge:** applications and hosts

- **network core:**
 - routers
 - network of networks

- **access networks, physical media:** communication links



The Network Edge

• end systems (hosts):

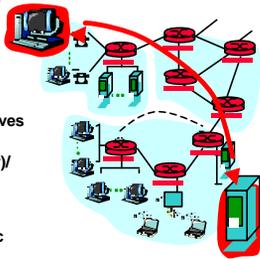
- run application programs
- e.g., WWW, email
- at "edge of network"

• client/server model

- client host requests, receives service from server
- e.g., WWW client (browser)/server; email client/server

• peer-peer model:

- host interaction symmetric
- e.g.: teleconferencing



Network edge: connection-oriented service

Goal: data transfer between end sys.

- **handshaking:** setup (prepare for) data transfer ahead of time
 - set up "state" in two communicating hosts
- **TCP - Transmission Control Protocol**
 - Internet's connection-oriented service

TCP service [RFC 793]

- **reliable, in-order byte-stream data transfer**
 - loss: acknowledgements and retransmissions
- **flow control:**
 - sender won't overwhelm receiver
- **congestion control:**
 - senders "slow down sending rate" when network congested

Network edge: connectionless service

Goal: data transfer between end systems

- same as before!

- **UDP - User Datagram Protocol [RFC 768]:** Internet's connectionless service
 - unreliable data transfer
 - no flow control
 - no congestion control

App's using TCP:

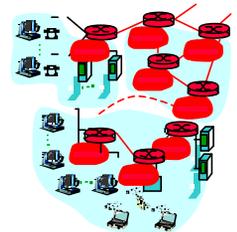
- HTTP (WWW), FTP (file transfer), Telnet (remote login), SMTP (email)

App's using UDP:

- streaming media, teleconferencing, Internet telephony

The Network Core

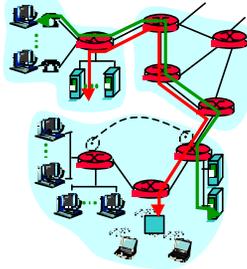
- mesh of interconnected routers
- **the fundamental question:** how is data transferred through net?
 - **circuit switching:** dedicated circuit per call: telephone net
 - **packet-switching:** data sent thru net in discrete "chunks"



Network Core: Circuit Switching

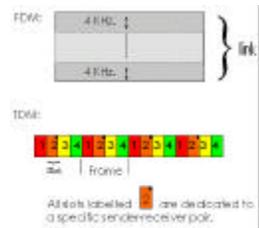
End-end resources reserved for "call"

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required



Network Core: Circuit Switching

- network resources (e.g., bandwidth) divided into "pieces"
- pieces allocated to calls
- resource piece *idle* if not used by owning call (*no sharing*)
- dividing link bandwidth into "pieces"
 - frequency division
 - time division



Network Core: Packet Switching

each end-end data stream divided into *packets*

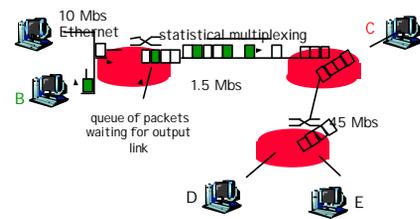
- user A, B packets *share* network resources
- each packet uses full link bandwidth
- resources used as *needed*

~~Bandwidth division into "pieces"~~
~~Dedicated allocation~~
~~Resource reservation~~

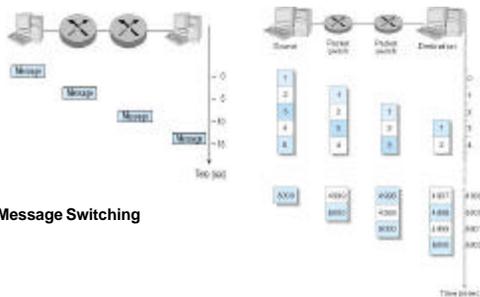
resource contention:

- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward: packets move one hop at a time
 - transmit over link
 - wait turn at next link

Network Core: Packet Switching



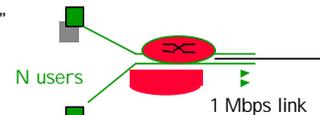
Network Core: Packet Switching



Packet switching versus circuit switching

Packet switching allows more users to use network!

- 1 Mbit link
- each user:
 - 100Kbps when "active"
 - active 10% of time
- circuit-switching:
 - 10 users
- packet switching:
 - with 35 users, probability > 10 active less than .004



Packet switching versus circuit switching

Is packet switching the “ultimate solution”

- Great for bursty data
 - resource sharing
 - no call setup
- **Excessive congestion:** packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- **Q: How to provide circuit-like behavior?**
 - bandwidth guarantees needed for audio/video appsstill an unsolved problem

Packet-switched networks: routing

- **Goal:** move packets among routers from source to destination
- **datagram network:**
 - destination address determines next hop
 - routes may change during session
 - analogy: driving, asking directions
- **virtual circuit network:**
 - each packet carries tag (virtual circuit ID), tag determines next hop
 - fixed path determined at call setup time, remains fixed thru call
 - routers maintain per-call state

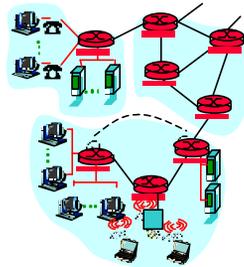
Access networks and physical media

Q: How to end systems connect to an edge router?

- residential access nets
- institutional access networks (school, company)
- mobile access networks

Keep in mind:

- bandwidth (bits per second) of access network?
- shared or dedicated?



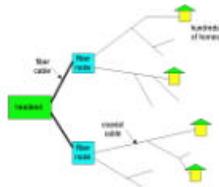
Residential access: point to point access

- **Dialup via modem**
 - up to 56Kbps direct access to router (conceptually)
- **ISDN:** integrated services digital network: 128Kbps all-digital connect to router
- **ADSL:** asymmetric digital subscriber line
 - up to 1 Mbps home-to-router
 - up to 8 Mbps router-to-home



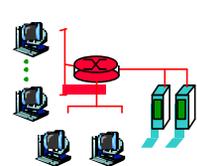
Residential access: cable modems

- **HFC: hybrid fiber coax**
 - asymmetric: up to 10Mbps upstream, 1 Mbps downstream
- **network** of cable and fiber attaches homes to ISP router
 - shared access to router among home
 - issues: congestion, dimensioning
- deployment: available via cable companies, e.g., MediaOne



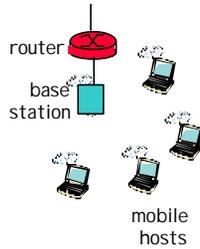
Institutional access: local area networks

- company/univ **local area network (LAN)** connects end system to edge router
- **Ethernet:**
 - shared or dedicated cable connects end system and router
 - 10 Mbs, 100Mbps, Gigabit Ethernet
- **deployment:** institutions, home LANs soon



Wireless access networks

- shared *wireless* access network connects end system to router
- **wireless LANs:**
 - radio spectrum replaces wire
 - e.g., Lucent Wavelan 10 Mbps
- **wider-area wireless access**
 - CDPD: wireless access to ISP router via cellular network



Physical Media

- **physical link:** transmitted data bit propagates across link
- **guided media:**
 - signals propagate in solid media: copper, fiber
- **unguided media:**
 - signals propagate freely e.g., radio

Twisted Pair (TP)

- two insulated copper wires
 - Category 3: traditional phone wires, 10 Mbps ethernet
 - Category 5 TP: 100Mbps ethernet



Physical Media: coax, fiber

Coaxial cable:

- wire (signal carrier) within a wire (shield)
 - baseband: single channel on cable
 - broadband: multiple channels on cable
- bidirectional
- common use in 10Mbps Ethernet



Fiber optic cable:

- glass fiber carrying light pulses
- high-speed operation:
 - 100Mbps Ethernet
 - high-speed point-to-point transmission (e.g., 5 Gbps)
- low error rate



Physical media: radio

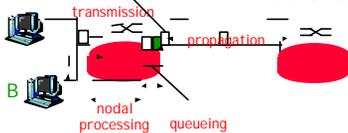
- signal carried in electromagnetic spectrum
- no physical "wire"
- bidirectional
- propagation environment effects:
 - reflection
 - obstruction by objects
 - interference

Radio link types:

- **microwave**
 - e.g. up to 45 Mbps channels
- **LAN** (e.g., waveLAN)
 - 2Mbps, 11Mbps
- **wide-area** (e.g., cellular)
 - e.g. CDPD, 10's Kbps
- **satellite**
 - up to 50Mbps channel (or multiple smaller channels)
 - 270 Msec end-end delay
 - geosynchronous versus LEOs

Delay in packet-switched networks

- packets experience **delay** on end-to-end path
- **four** sources of delay at each hop
 - **nodal processing:**
 - check bit errors
 - determine output link
 - **queueing**
 - time waiting at output link for transmission
 - depends on congestion level of router



Delay in packet-switched networks

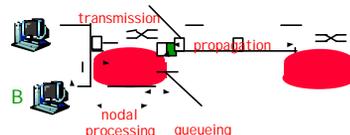
Transmission delay:

- R = link bandwidth (bps)
- L = packet length (bits)
- time to send bits into link = L/R

Propagation delay:

- d = length of physical link
- s = propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- propagation delay = d/s

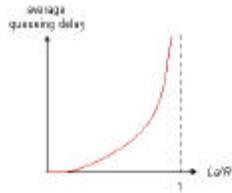
Note: s and R are very different quantities!



Queuing delay (revisited)

- R =link bandwidth (bps)
- L =packet length (bits)
- a =average packet arrival rate

traffic intensity = La/R



- $La/R \sim 0$: average queuing delay small
- $La/R \leq 1$: delays become large
- $La/R > 1$: more “work” arriving than can be serviced, average delay infinite!

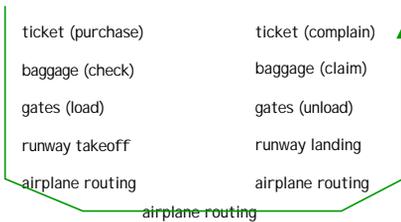
Protocol “Layers”

Networks are complex!

- many “pieces”:
 - hosts
 - routers
 - links of various media
 - applications
 - protocols
 - hardware, software

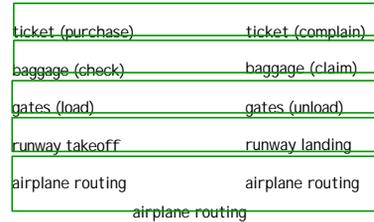
Question:
Is there any hope of
organizing
structure of
network?

Organization of air travel



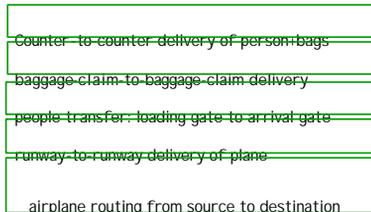
- a series of steps

Organization of air travel: a different view

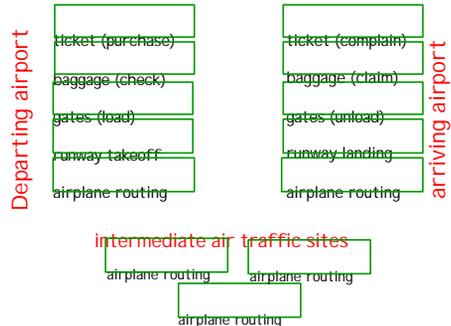


- **Layers:** each layer implements a service
 - via its own internal-layer actions
 - relying on services provided by layer below

Layered air travel: services



Distributed implementation of layer functionality



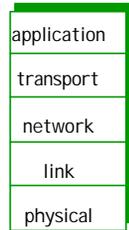
Why layering?

Dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
 - layered **reference model** for discussion
- modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful?

Internet protocol stack

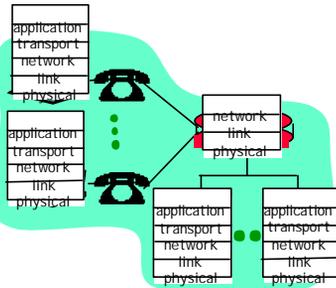
- **application:** supporting network applications
 - ftp, smtp, http
- **transport:** process-process data transfer
 - tcp, udp
- **network:** routing of datagrams from source to destination
 - ip, routing protocols
- **link:** data transfer between neighboring network elements
 - ppp, ethernet
- **physical:** bits "on the wire"



Layering: logical communication

Each layer:

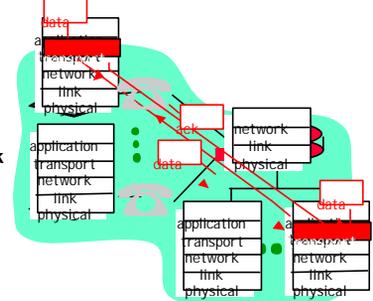
- distributed
- "entities" implement layer functions at each node
- entities perform actions, exchange messages with peers



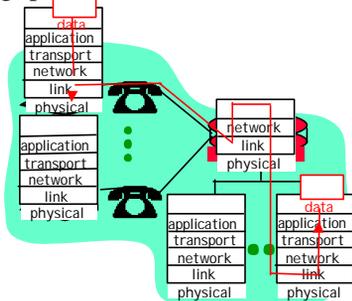
Layering: logical communication

E.g.: transport

- take data from app
- add addressing, reliability check info to form "datagram"
- send datagram to peer
- wait for peer to ack receipt
- analogy: post office delivery to person



Layering: physical communication



Protocol layering and data

Each layer takes data from above

- adds header information to create new data unit
- passes new data unit to layer below

