

# Introduction

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# What is a network?

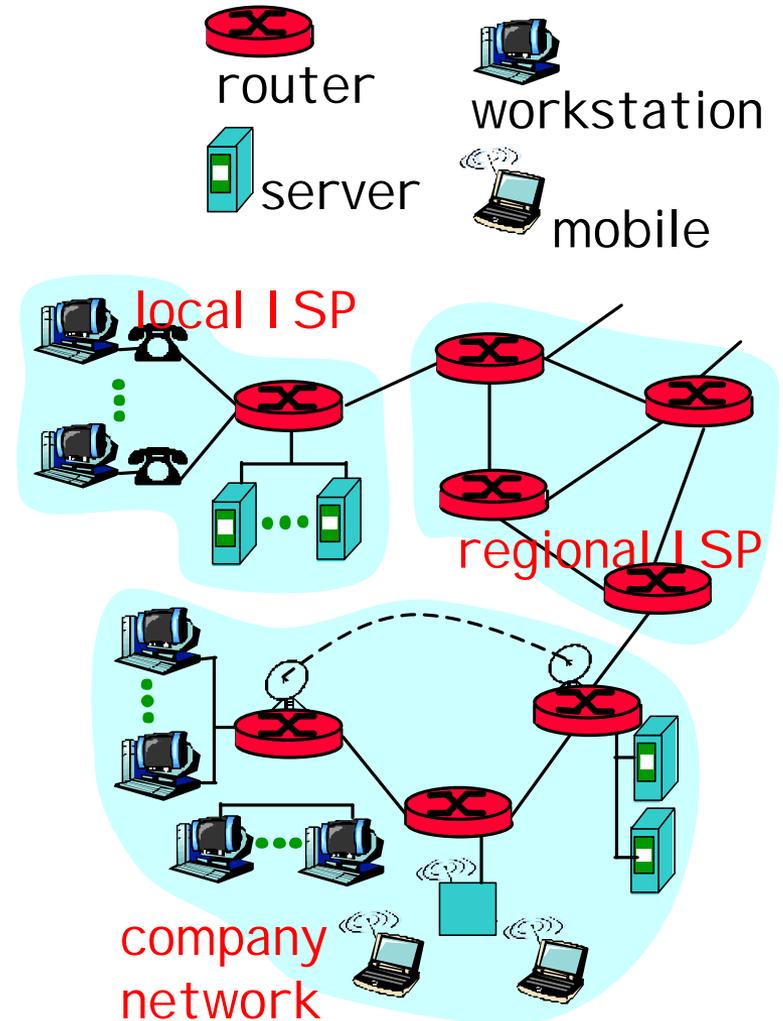
- **Carrier of information between connected entities**
- **What does a network consist of?**
  - End hosts connected to the network
  - Routers/switches that move data through the network
  - Physical links that carry information
    - E.g. Ethernet, FDDI, ATM, Token Ring
  - Applications that communicate with each other to provide services
    - E-Mail, File Transfer, Web Browser

# What is an internetwork?

- A set of connected networks is called an internetwork.
- The *Internet* is a specific example of an internetwork.
  - The Internet is characterized by the use of a common network protocol (IP).
- An internetwork need NOT be connected to the *Internet*

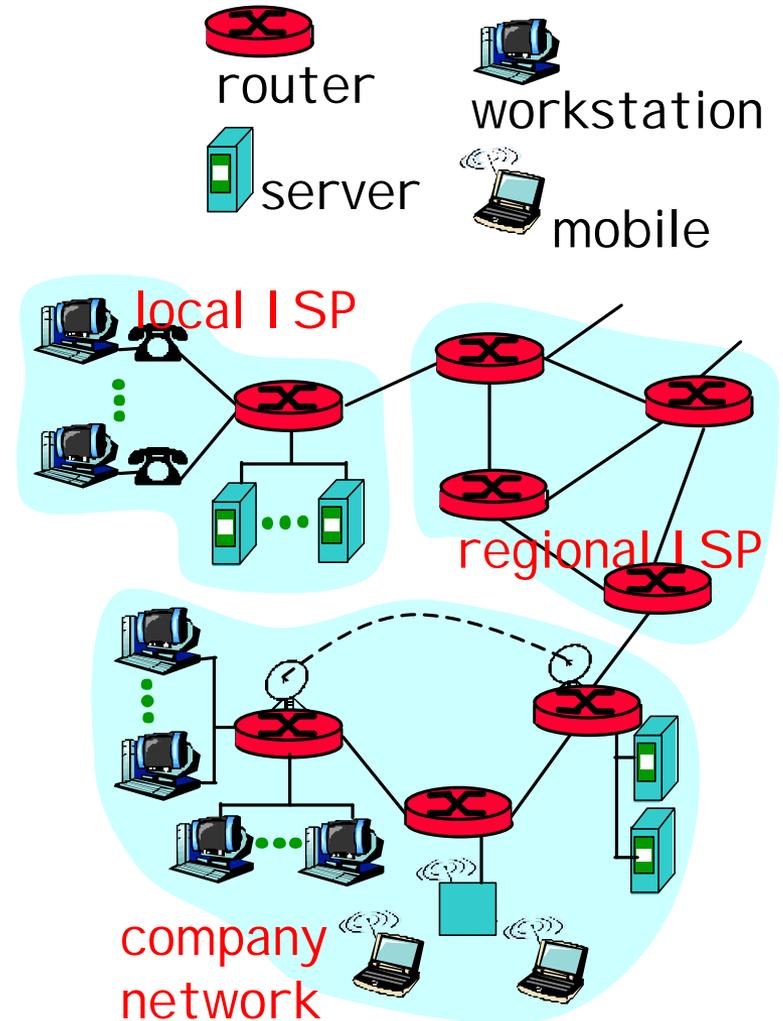
# Internet: Components

- **Network edge:**
  - millions of connected computing devices running **network apps**
  - pc's workstations, servers
  - PDA's phones, toasters
- **Network core:**
  - **routers:** forward packets (chunks) of data thru network
- **Media:**
  - **communication links:** fiber, copper, radio, satellite



# Internet: Components

- **protocols:** control sending, receiving of msgs
  - e.g., TCP, IP, HTTP, FTP, PPP
- **Internet: “network of networks”**
  - loosely hierarchical
  - public Internet versus private intranet
- **Internet standards**
  - RFC: Request for comments
  - IETF: Internet Engineering Task Force



# Internet: Services

- **communication *infrastructure*** enables distributed applications:
  - WWW, email, games, e-commerce, database., voting,
  - more?
- **communication services provided:**
  - connectionless
  - connection-oriented
- **cyberspace [Gibson]:**
  - “a consensual hallucination experienced daily by billions of operators, in every nation, ....”

# What's a protocol?

## human protocols:

- “what’s the time?”
- “I have a question”
- introductions

... specific msgs sent

... specific actions  
taken when msgs  
received, or other  
events

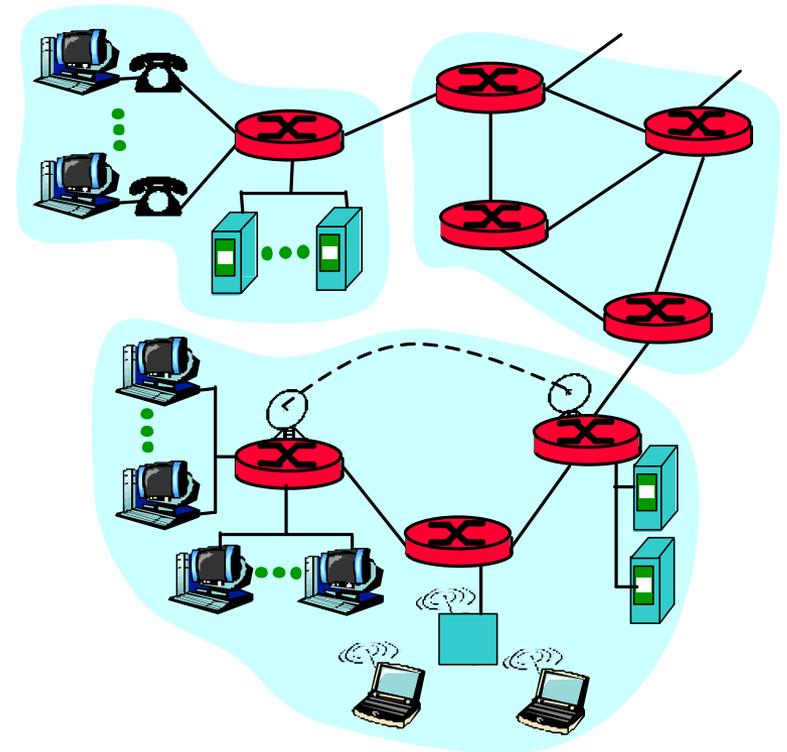
## network protocols:

- machines rather than humans
- all communication activity in Internet governed by protocols

*protocols define format,  
order of msgs sent and  
received among network  
entities, and actions  
taken on msg  
transmission, receipt*

# A closer look at network structure

- **network edge:**  
applications and hosts
- **network core:**
  - routers
  - network of networks
- **access networks, physical media:**  
communication links



# The Network Edge

- **end systems (hosts):**

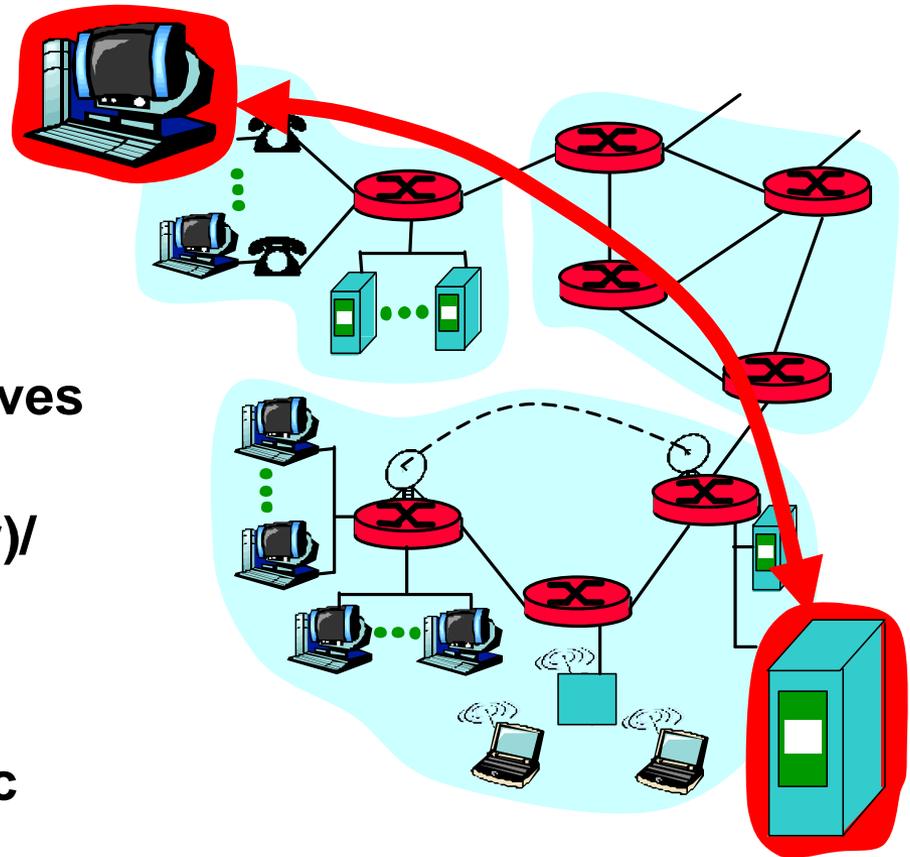
- run application programs
- e.g., WWW, email
- at “edge of network”

- **client/server model**

- client host requests, receives service from server
- e.g., WWW client (browser)/ server; email client/server

- **peer-peer model:**

- host interaction symmetric
- e.g.: teleconferencing



# Network edge: connection-oriented service

**Goal:** data transfer  
between end sys.

- ***handshaking:*** setup (prepare for) data transfer ahead of time
  - ***set up “state”*** in two communicating hosts
- **TCP - Transmission Control Protocol**
  - Internet’s connection-oriented service

**TCP service** [RFC 793]

- ***reliable, in-order*** byte-stream data transfer
  - loss: acknowledgements and retransmissions
- ***flow control:***
  - sender won’t overwhelm receiver
- ***congestion control:***
  - senders “slow down sending rate” when network congested

# Network edge: connectionless service

- Goal:** data transfer between end systems
- same as before!
  - **UDP** - User Datagram Protocol [RFC 768]: Internet's connectionless service
    - unreliable data transfer
    - no flow control
    - no congestion control

## **App's using TCP:**

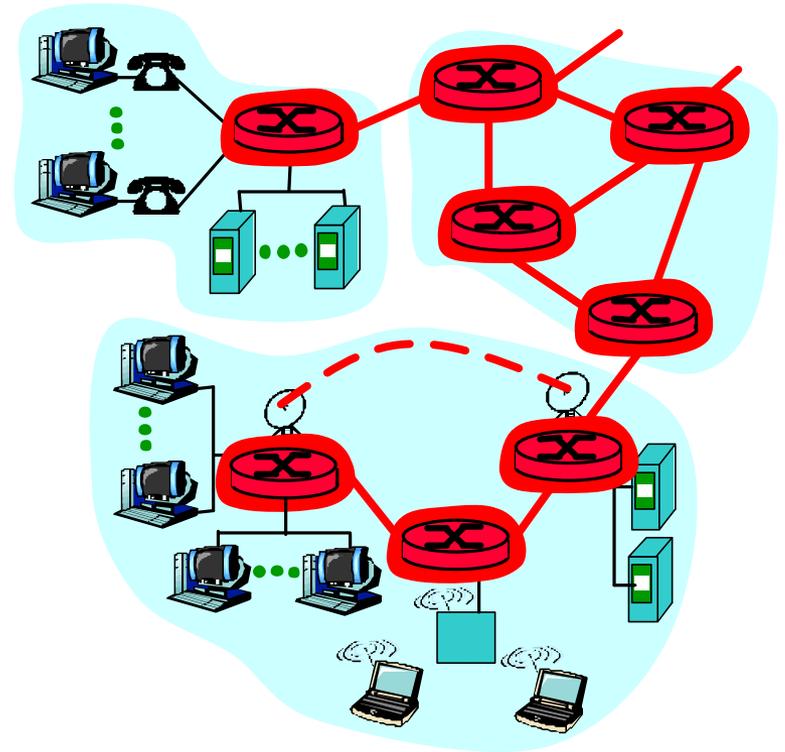
- HTTP (WWW), FTP (file transfer), Telnet (remote login), SMTP (email)

## **App's using UDP:**

- streaming media, teleconferencing, Internet telephony

# The Network Core

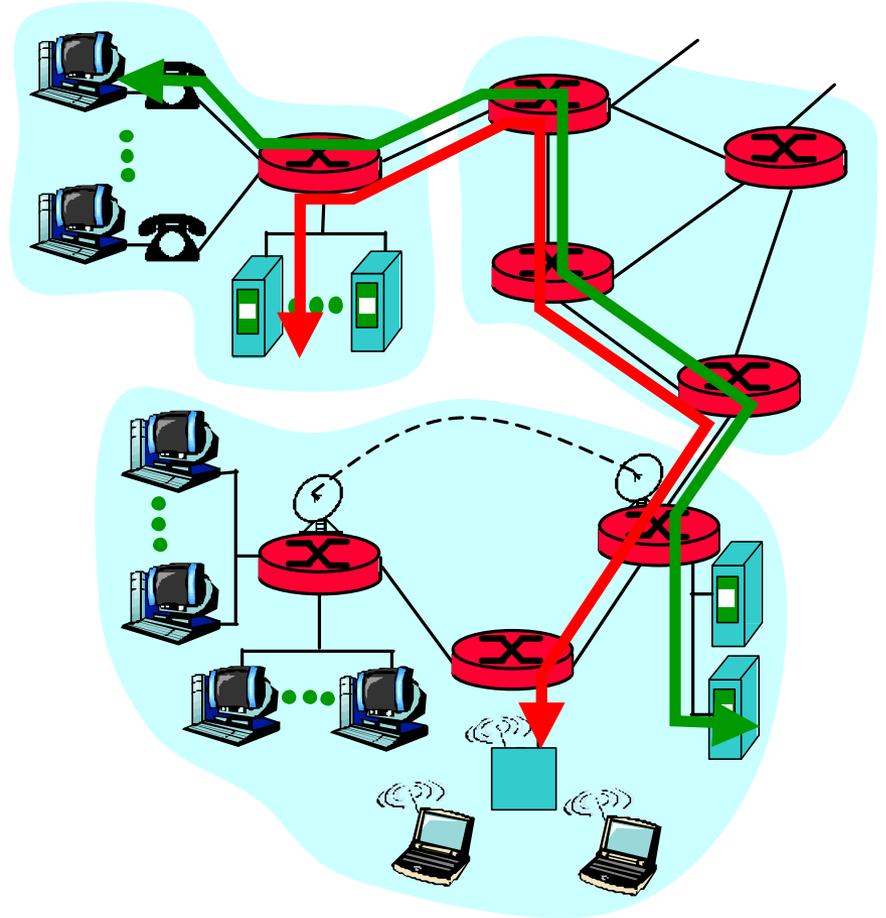
- mesh of interconnected routers
- **the fundamental question:** how is data transferred through net?
  - **circuit switching:** dedicated circuit per call: telephone net
  - **packet-switching:** data sent thru net in discrete “chunks”



# Network Core: Circuit Switching

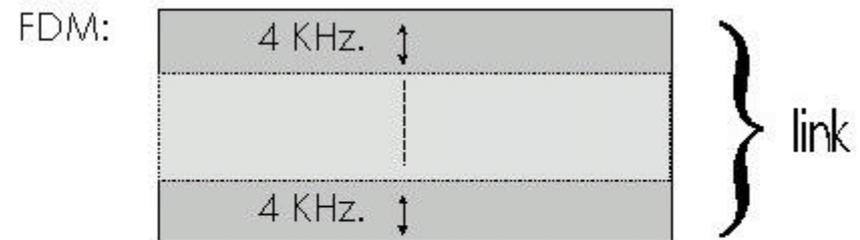
**End-end resources reserved for “call”**

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required

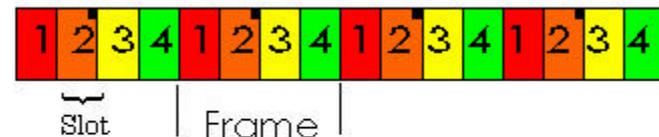


# Network Core: Circuit Switching

- network resources (e.g., bandwidth) **divided into “pieces”**
- pieces allocated to calls
- resource piece *idle* if not used by owning call (*no sharing*)
- dividing link bandwidth into “pieces”
  - frequency division
  - time division



TDM:



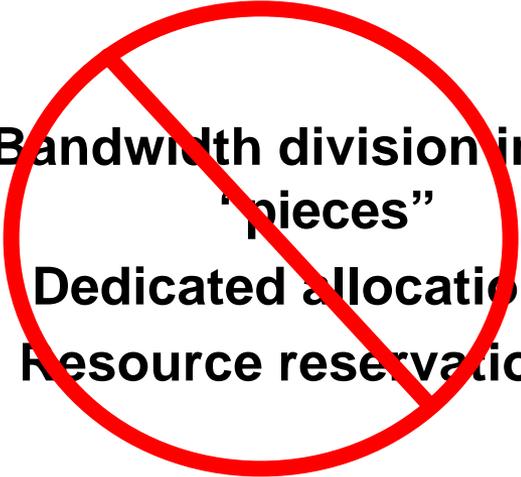
All slots labelled  are dedicated to a specific sender-receiver pair.

# Network Core: Packet Switching

**each end-end data stream  
divided into *packets***

- user A, B packets *share* network resources
- each packet uses full link bandwidth
- resources used as *needed*,

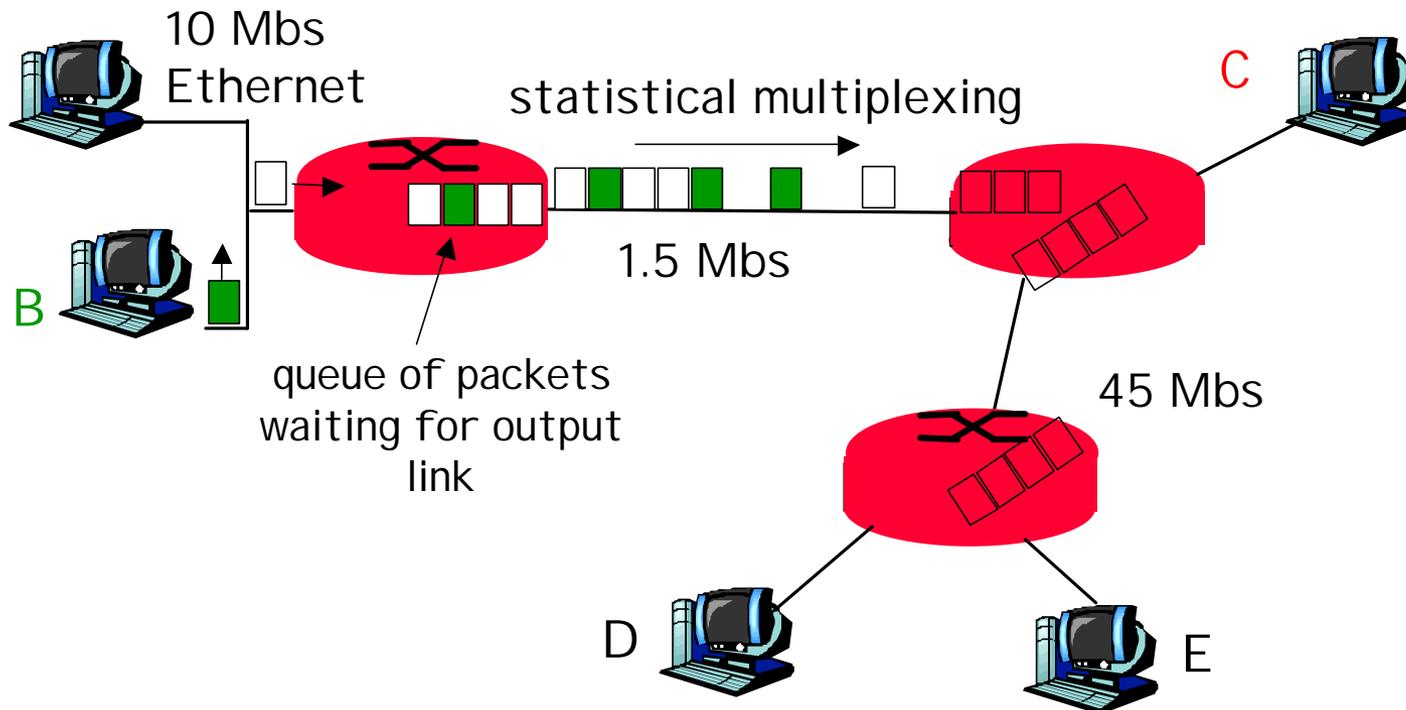
Bandwidth division into  
“pieces”  
Dedicated allocation  
Resource reservation



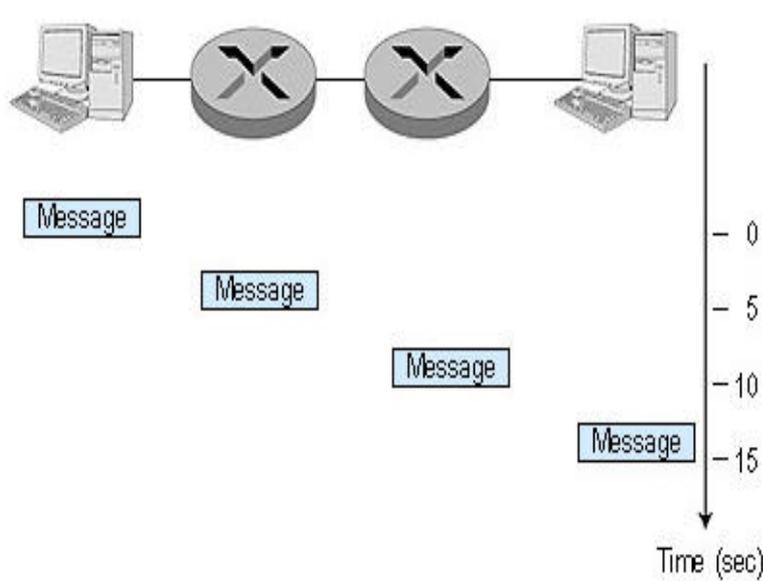
**resource contention:**

- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward: packets move one hop at a time
  - transmit over link
  - wait turn at next link

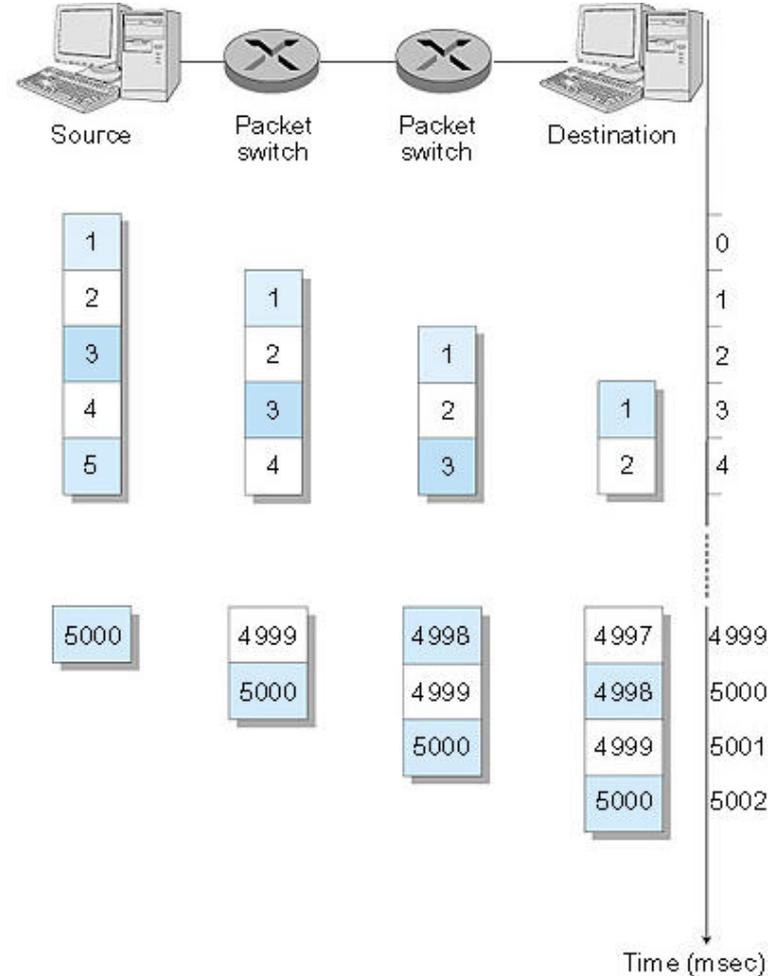
# Network Core: Packet Switching



# Network Core: Packet Switching



**Message Switching**

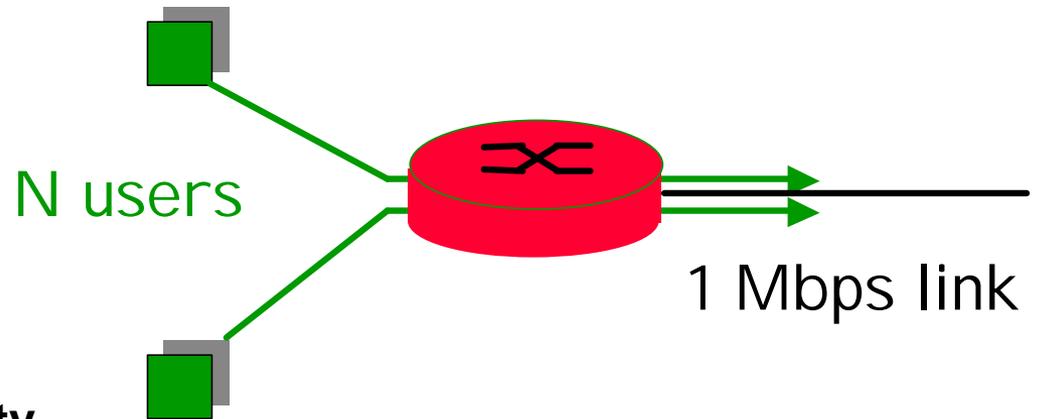


**Segmented Packet Switching**

# Packet switching versus circuit switching

**Packet switching allows more users to use network!**

- **1 Mbit link**
- **each user:**
  - 100Kbps when “active”
  - active 10% of time
- **circuit-switching:**
  - 10 users
- **packet switching:**
  - with 35 users, probability > 10 active less than .004



# Packet switching versus circuit switching

## Is packet switching the “ultimate solution”

- **Great for bursty data**
  - resource sharing
  - no call setup
- **Excessive congestion:** packet delay and loss
  - protocols needed for reliable data transfer, congestion control
- **Q: How to provide circuit-like behavior?**
  - bandwidth guarantees needed for audio/video apps

still an unsolved problem

# Packet-switched networks: routing

- **Goal:** move packets among routers from source to destination
- **datagram network:**
  - *destination address* determines next hop
  - routes may change during session
  - analogy: driving, asking directions
- **virtual circuit network:**
  - each packet carries tag (virtual circuit ID), tag determines next hop
  - fixed path determined at *call setup time*, remains fixed thru call
  - routers maintain per-call state

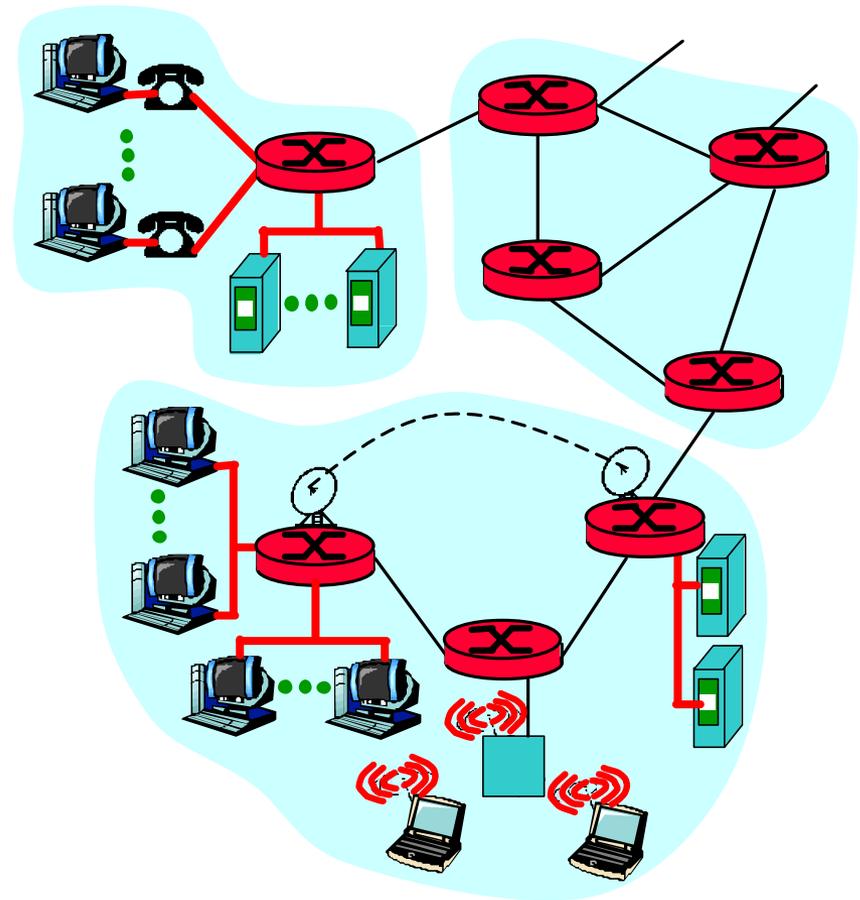
# Access networks and physical media

**Q: How to end systems connect to an edge router?**

- residential access nets
- institutional access networks (school, company)
- mobile access networks

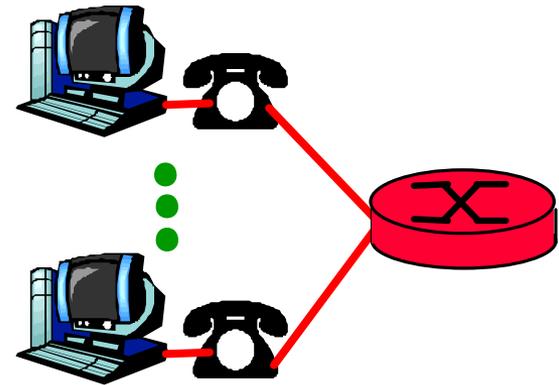
**Keep in mind:**

- bandwidth (bits per second) of access network?
- shared or dedicated?



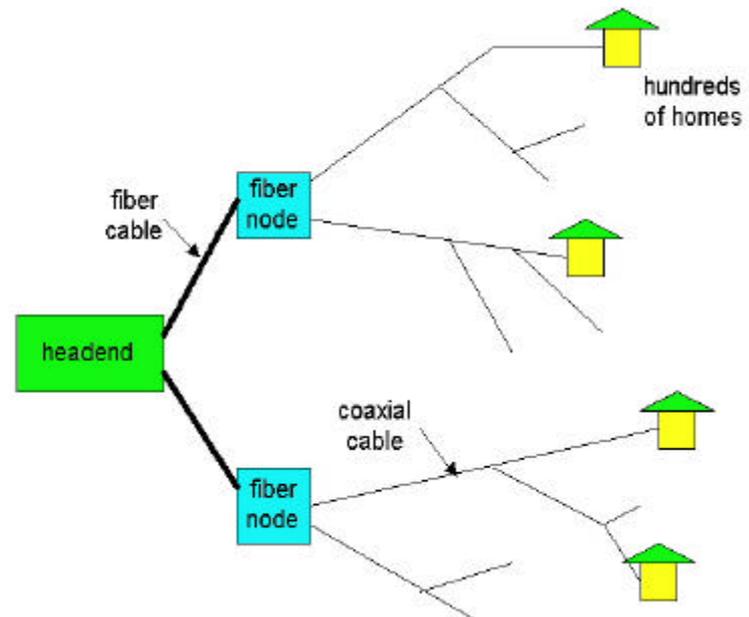
# Residential access: point to point access

- **Dialup via modem**
  - up to 56Kbps direct access to router (conceptually)
- **ISDN**: integrated services digital network: 128Kbps all-digital connect to router
- **ADSL**: asymmetric digital subscriber line
  - up to 1 Mbps home-to-router
  - up to 8 Mbps router-to-home



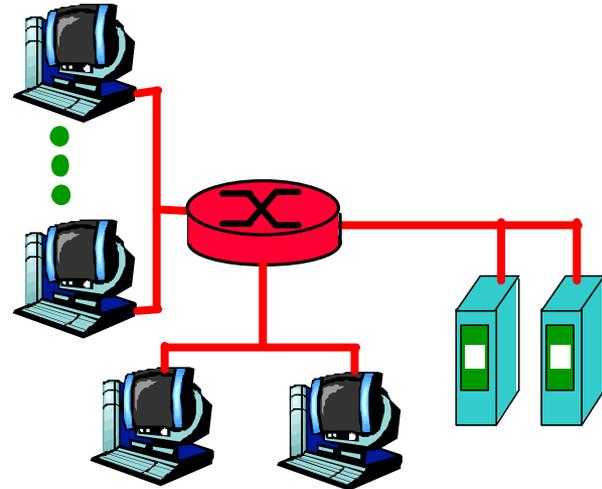
# Residential access: cable modems

- **HFC: hybrid fiber coax**
  - asymmetric: up to 10Mbps upstream, 1 Mbps downstream
- **network of cable and fiber attaches homes to ISP router**
  - shared access to router among home
  - issues: congestion, dimensioning
- **deployment: available via cable companies, e.g., MediaOne**



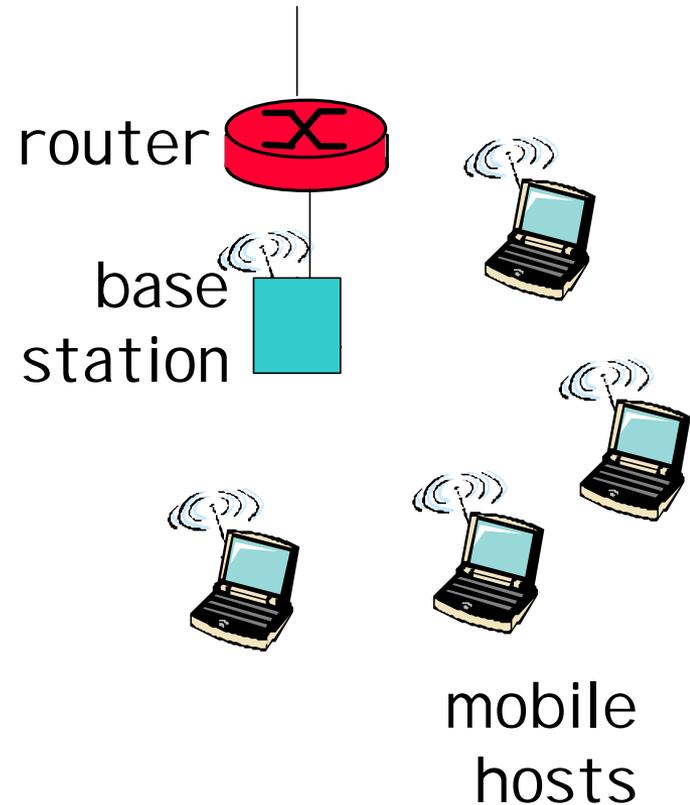
# Institutional access: local area networks

- company/univ **local area network** (LAN) connects end system to edge router
- **Ethernet:**
  - shared or dedicated cable connects end system and router
  - 10 Mbs, 100Mbps, Gigabit Ethernet
- **deployment:** institutions, home LANs soon



# Wireless access networks

- shared *wireless access network* connects end system to router
- **wireless LANs:**
  - radio spectrum replaces wire
  - e.g., Lucent Wavelan 10 Mbps
- **wider-area wireless access**
  - CDPD: wireless access to ISP router via cellular network



# Physical Media

- **physical link:**  
transmitted data bit propagates across link
- **guided media:**
  - signals propagate in solid media: copper, fiber
- **unguided media:**
  - signals propagate freely e.g., radio

## Twisted Pair (TP)

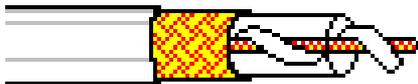
- two insulated copper wires
  - Category 3: traditional phone wires, 10 Mbps ethernet
  - Category 5 TP: 100Mbps ethernet



# Physical Media: coax, fiber

## Coaxial cable:

- wire (signal carrier) within a wire (shield)
  - baseband: single channel on cable
  - broadband: multiple channels on cable
- bidirectional
- common use in 10Mbps Ethernet



## Fiber optic cable:

- glass fiber carrying light pulses
- high-speed operation:
  - 100Mbps Ethernet
  - high-speed point-to-point transmission (e.g., 5 Gbps)
- low error rate



# Physical media: radio

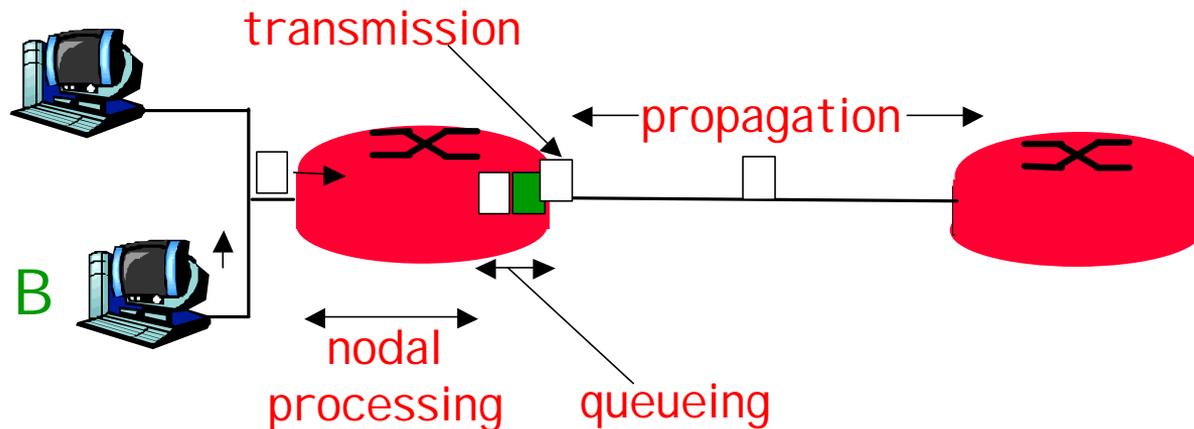
- signal carried in electromagnetic spectrum
- no physical “wire”
- bidirectional
- propagation environment effects:
  - reflection
  - obstruction by objects
  - interference

## Radio link types:

- **microwave**
  - e.g. up to 45 Mbps channels
- **LAN** (e.g., waveLAN)
  - 2Mbps, 11Mbps
- **wide-area** (e.g., cellular)
  - e.g. CDPD, 10’s Kbps
- **satellite**
  - up to 50Mbps channel (or multiple smaller channels)
  - 270 Msec end-end delay
  - geosynchronous versus LEOS

# Delay in packet-switched networks

- packets experience **delay** on end-to-end path
- **four** sources of delay at each hop
- **nodal processing**:
  - check bit errors
  - determine output link
- **queueing**
  - time waiting at output link for transmission
  - depends on congestion level of router



# Delay in packet-switched networks

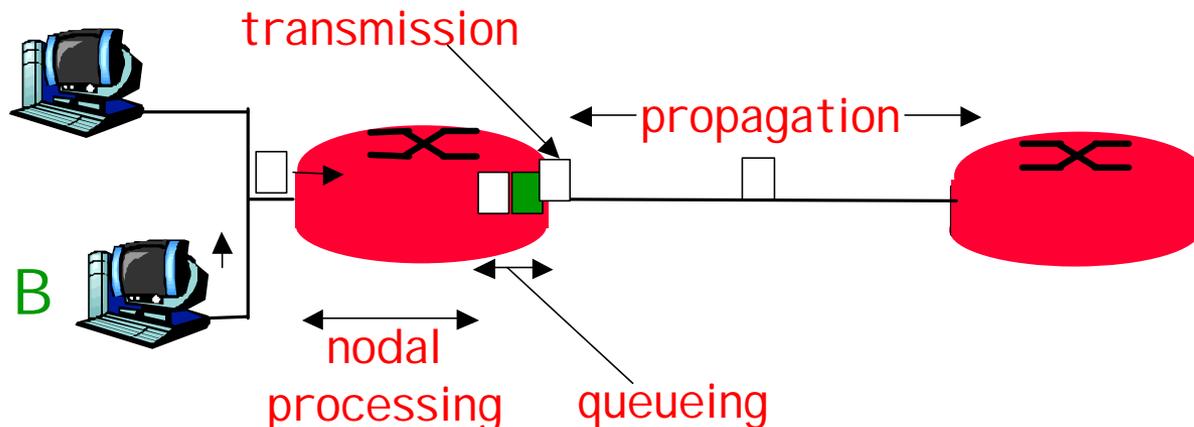
## Transmission delay:

- $R$  = link bandwidth (bps)
- $L$  = packet length (bits)
- time to send bits into link =  $L/R$

## Propagation delay:

- $d$  = length of physical link
- $s$  = propagation speed in medium ( $\sim 2 \times 10^8$  m/sec)
- propagation delay =  $d/s$

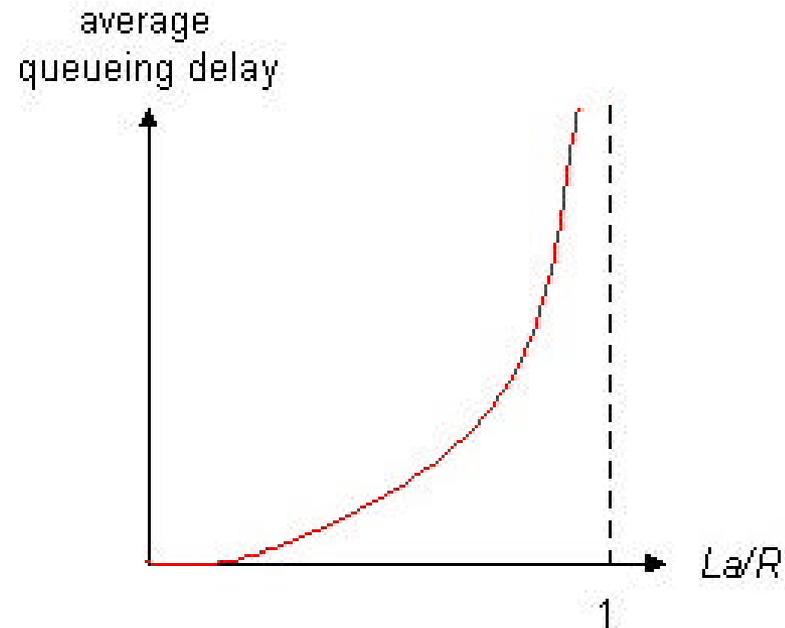
**Note:**  $s$  and  $R$  are very different quantities!



# Queuing delay (revisited)

- $R$ =link bandwidth (bps)
- $L$ =packet length (bits)
- $a$ =average packet arrival rate

**traffic intensity =  $La/R$**



- $La/R \sim 0$ : average queuing delay small
- $La/R \leq 1$ : delays become large
- $La/R > 1$ : more “work” arriving than can be serviced, average delay infinite!

# Protocol “Layers”

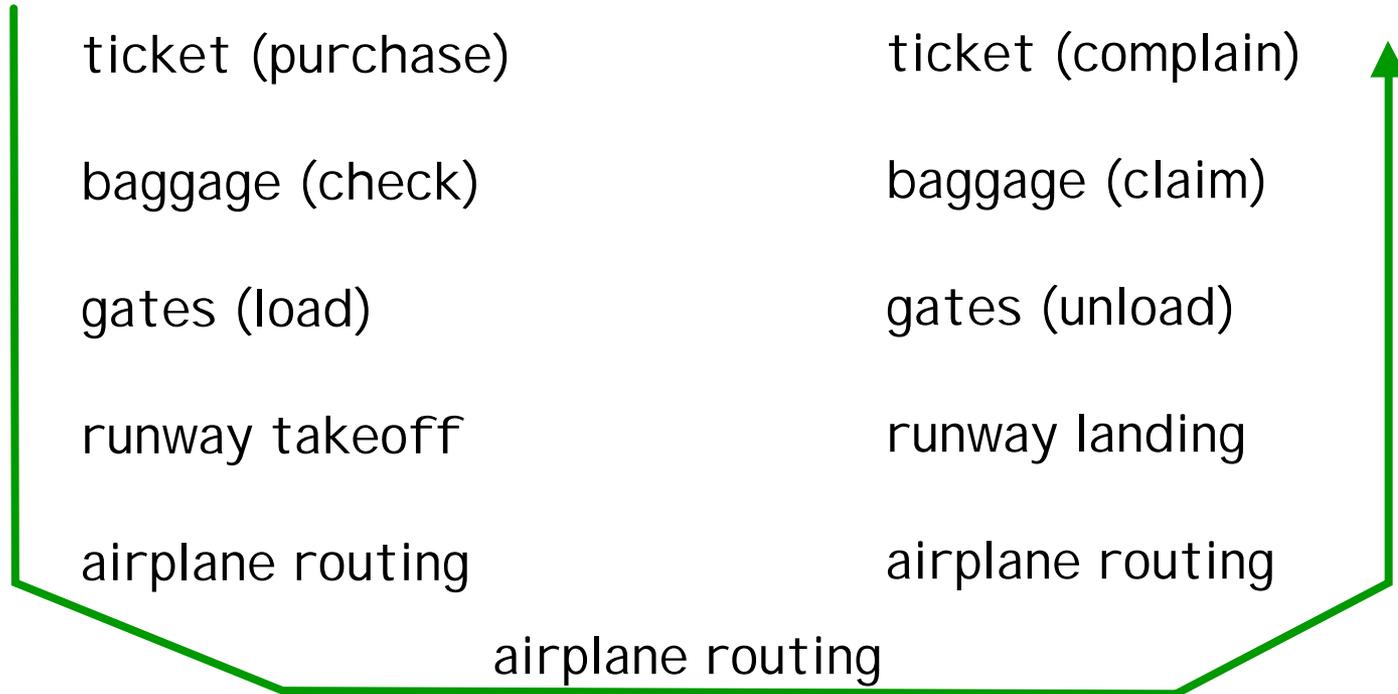
## Networks are complex!

- many “pieces”:
  - hosts
  - routers
  - links of various media
  - applications
  - protocols
  - hardware, software

## Question:

Is there any hope of  
*organizing*  
structure of  
network?

# Organization of air travel



- **a series of steps**

# Organization of air travel: a different view

ticket (purchase)	ticket (complain)
baggage (check)	baggage (claim)
gates (load)	gates (unload)
runway takeoff	runway landing
airplane routing	airplane routing
airplane routing	

**Layers:** each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

# Layered air travel: services

Counter-to-counter delivery of person+bags

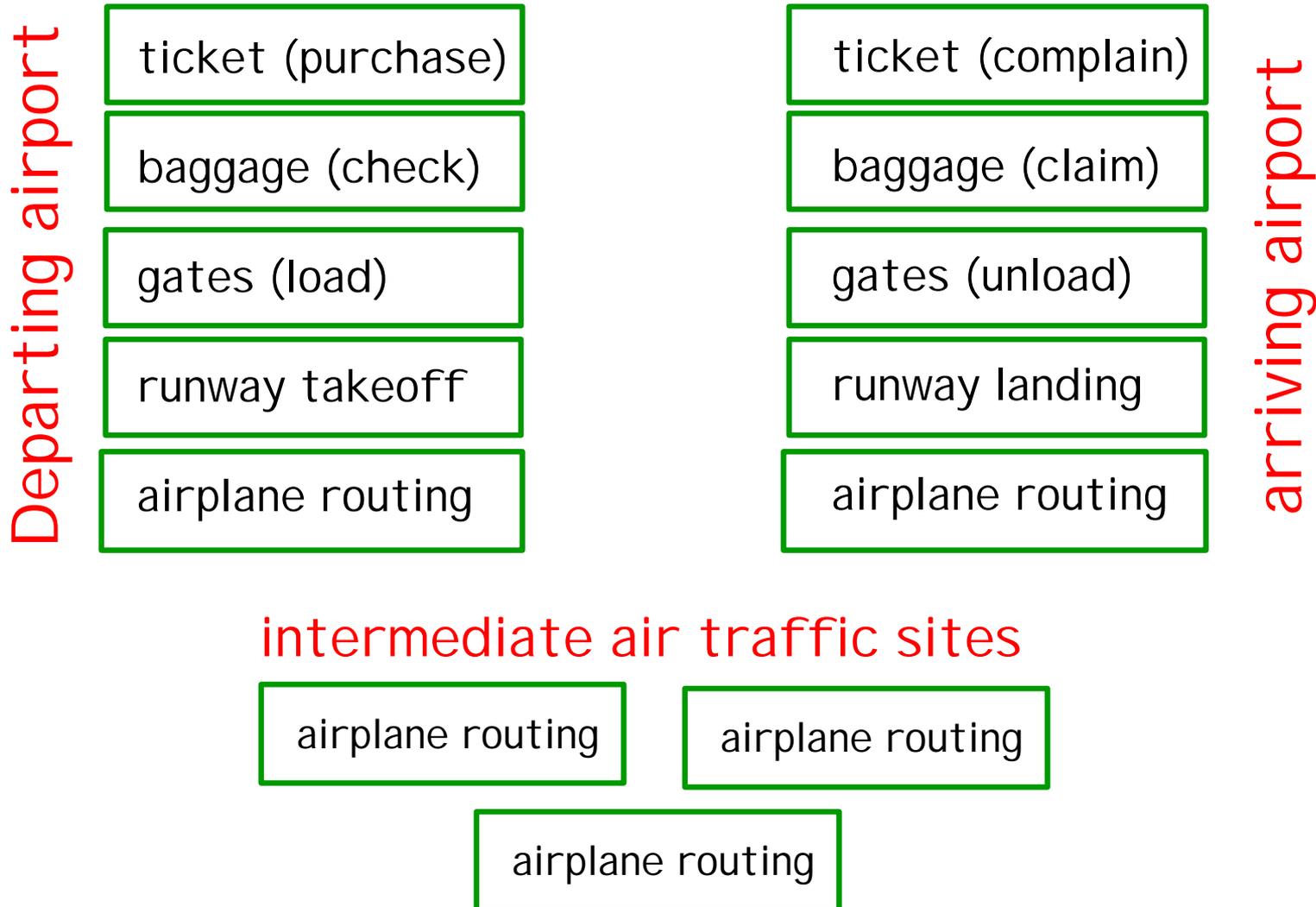
baggage-claim-to-baggage-claim delivery

people transfer: loading gate to arrival gate

runway-to-runway delivery of plane

airplane routing from source to destination

# Distributed implementation of layer functionality



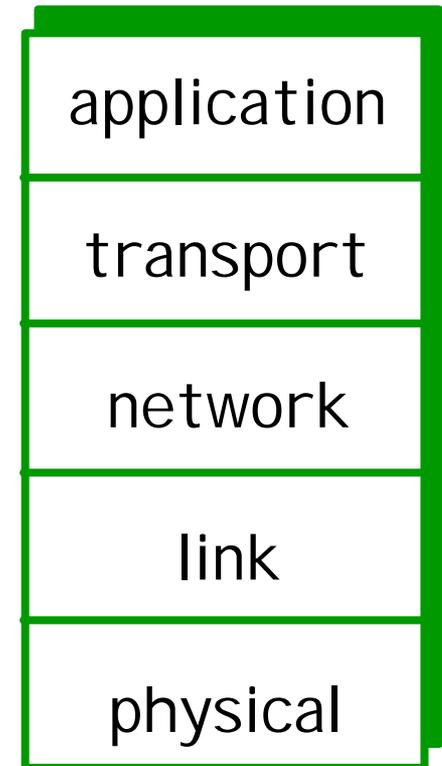
# Why layering?

## Dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
  - layered **reference model** for discussion
- modularization eases maintenance, updating of system
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful?

# Internet protocol stack

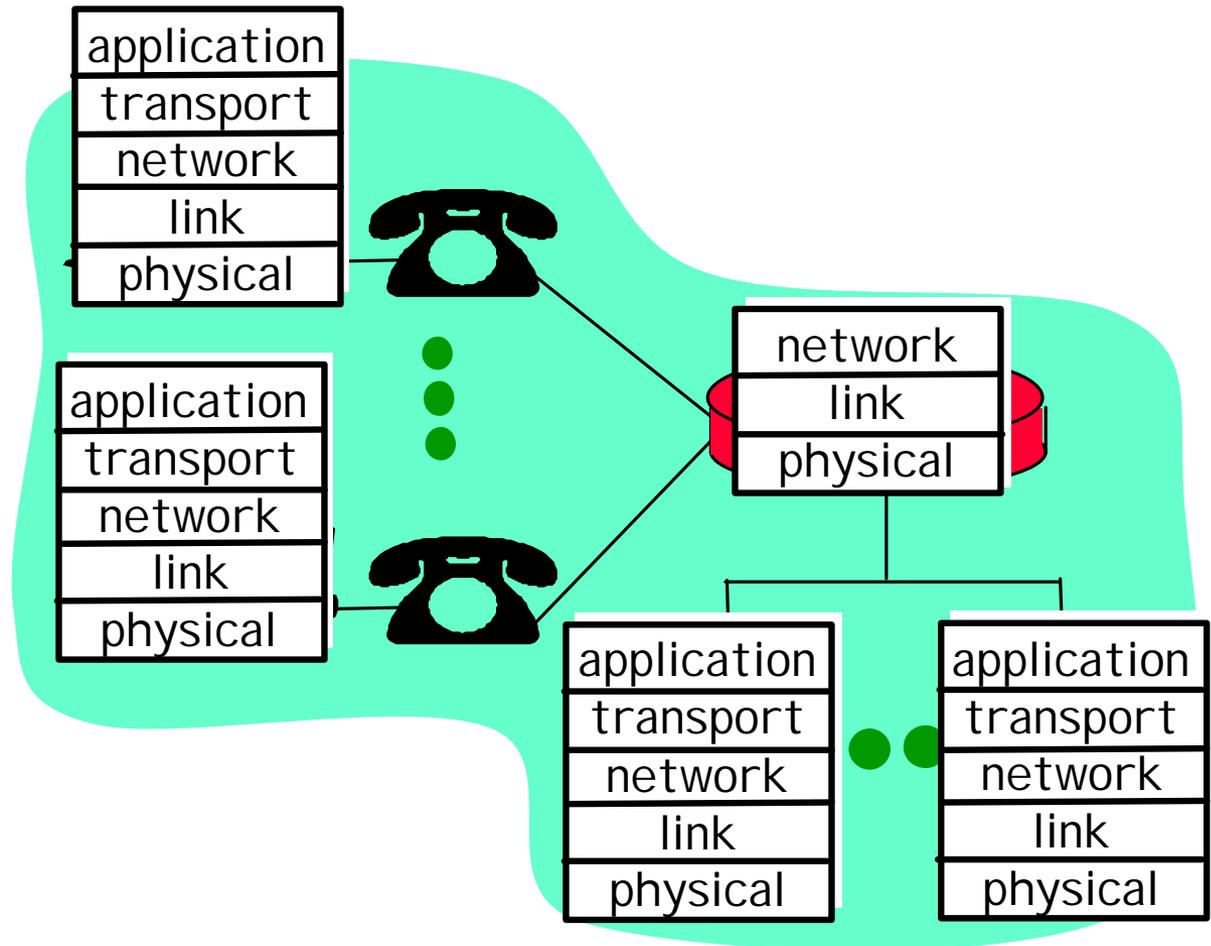
- **application:** supporting network applications
  - ftp, smtp, http
- **transport:** process-process data transfer
  - tcp, udp
- **network:** routing of datagrams from source to destination
  - ip, routing protocols
- **link:** data transfer between neighboring network elements
  - ppp, ethernet
- **physical:** bits “on the wire”



# Layering: logical communication

## Each layer:

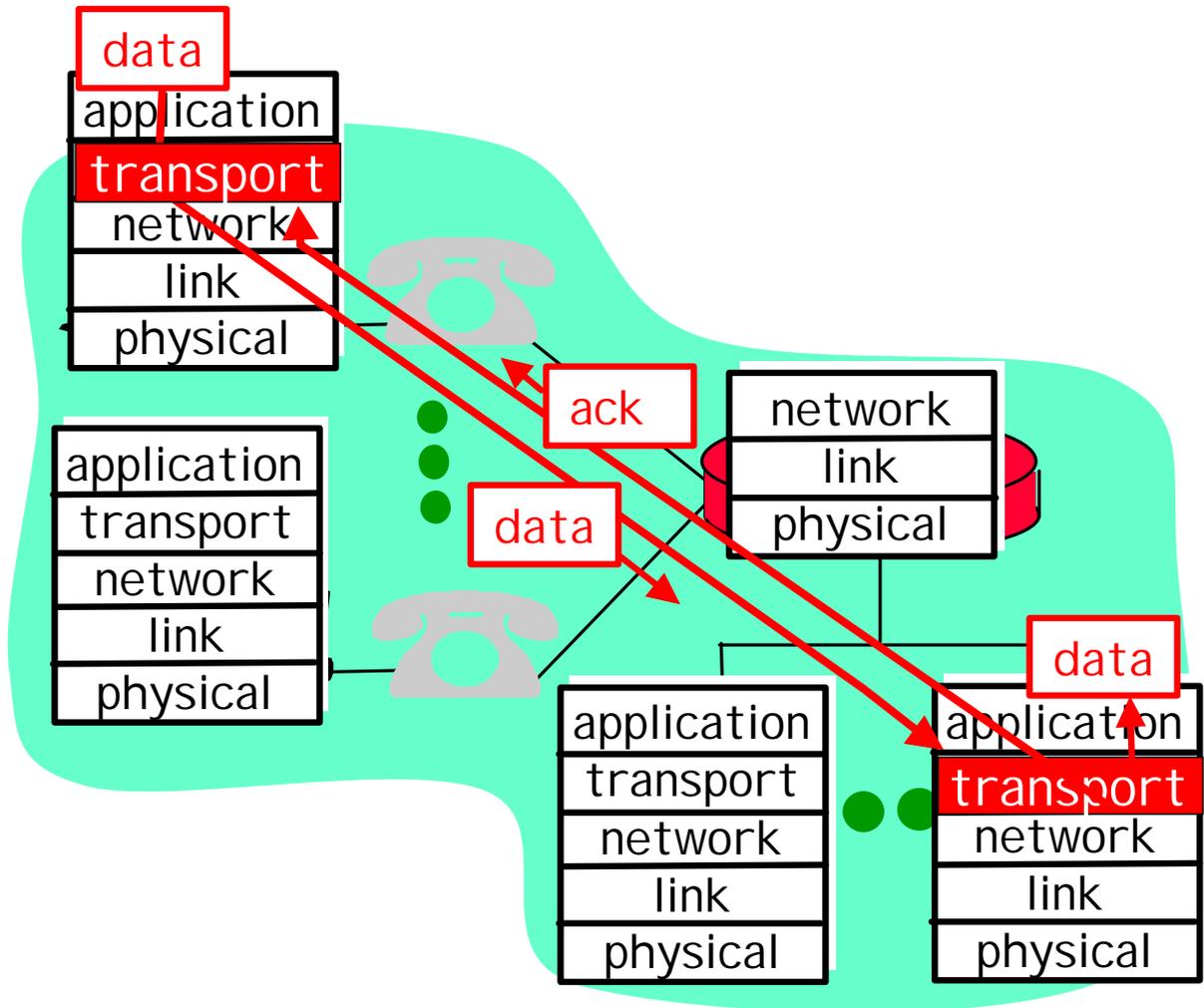
- distributed
- “entities” implement layer functions at each node
- entities perform actions, exchange messages with peers



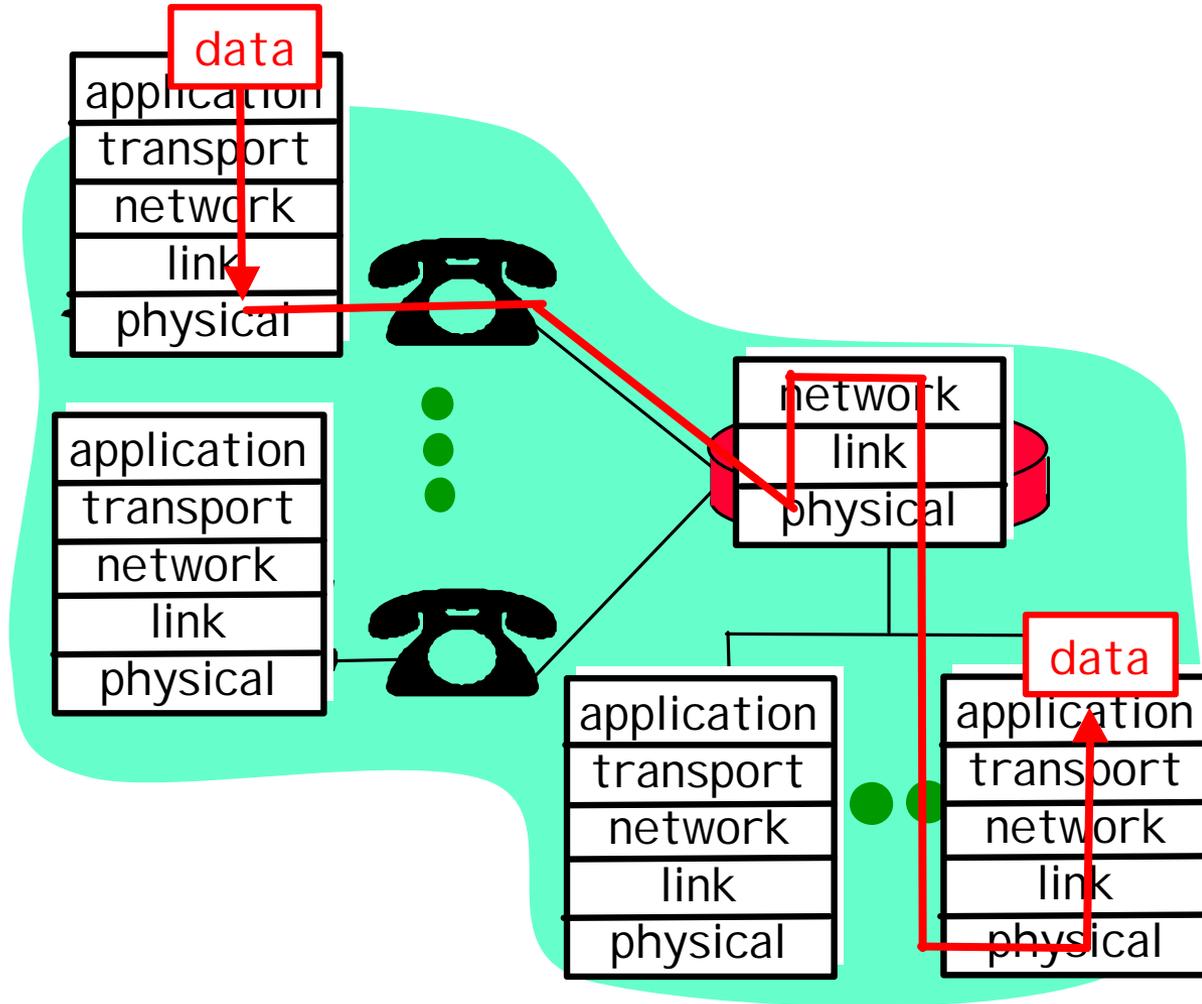
# Layering: *logical* communication

## E.g.: transport

- take data from app
- add addressing, reliability check info to form “datagram”
- send datagram to peer
- wait for peer to ack receipt
- analogy: post office delivery to person



# Layering: physical communication



# Protocol layering and data

**Each layer takes data from above**

- **adds header information to create new data unit**
- **passes new data unit to layer below**

