

**CS 3214: Project 1** 

# The Customizable Shell

Help Session: Wednesday Sep 13, 2023 7:30 PM

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## **Topics**

- Shell Concepts
- Project Overview / Logistics
- Version Control (Git & Gitlab)
- Debugging (GDB & Valgrind)
- Advice
- Q & A



# Shell Concepts

#### What is a shell?

- Command Interpreter
  - Reads user input and executes user requests
  - Not to be confused with a "Terminal" (next slide explains distinction)



#### **Terminal vs Shell**

Terminal (the front-end GUI of our shell)



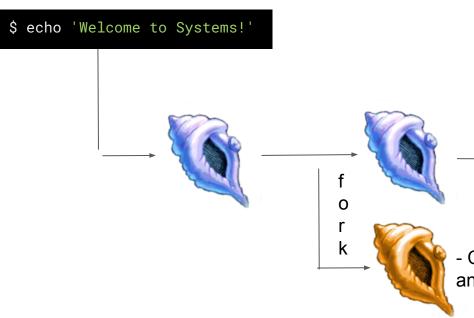
**Examples:** gnome-terminal, terminator, Terminal.app (macOS) etc.

Shell (an executable with no GUI)

```
root@terminal:~# love
-bash: love not found
root@terminal:~# happiness
-bash: happiness not found
root@terminal:~# peace
-bash: peace not found
root@terminal:~# kill
-bash: you need to specify whom
to kill
```

This terminal is running bash, a shell program

#### **Behind the Scenes**



#### FOUR STEPS for *non-built-in*

- 1. Shell waits for user input
- 2. Shell interprets command
- 3. Forks a process
- 4. If the command is a foreground process, the parent waits for the child to finish before taking in new command. If it's a background process, the parent waits for a signal from the child but at the meantime accepts additional command.

Welcome to Systems!

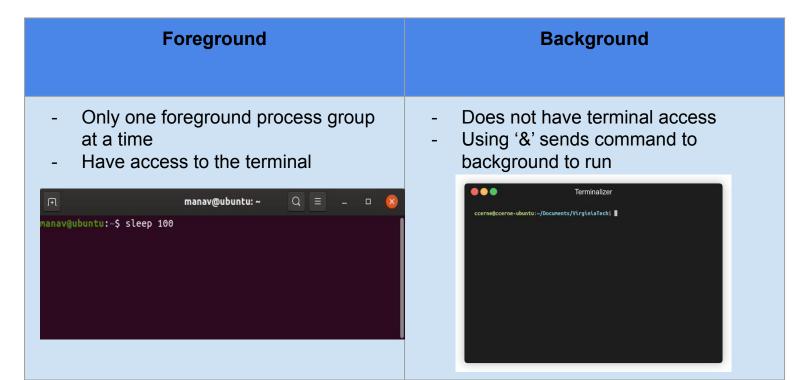
 Child executes the command and writes to stdout

#### Additional Features for the Shell (where you come in)

- Foreground / Background Processes
- Process Groups
- Built-in Commands
- I/O Piping
- I/O Redirection
- Signal Handling

#### Foreground / Background Processes

The shell can fork processes into the foreground or background



#### **Process Groups**

- A Job is essentially a pipelined-command
- Each Job has its own process group
  - Each command within a Job should have the same PGID
  - Two methodologies of creating new processes:
    - o fork() and execvp()
    - posix\_spawn
- Jobs are deleted when they are completed
  - Be careful not to delete a job prematurely
  - See the comment above wait\_for\_job()

#### <justv@cottonwood justv>\$ sleep 20 | sleep 20 | sleep 20 &

```
<justv@cottonwood justv>$ ps xj | head -n 1; ps xj |
                                                    tail -n 6
  PPID
           PID
                  PGID
                           SID TTY
                                          TPGID STAT
                                                       UID
                                                             TIME COMMAND
1357688 1363886 1363886 1357688 pts/0
                                        1365438 S
                                                             0:00 /home/courses/cs3214/bin/cush-gback
                                                     24908
1363886 1365308 1365308 1357688 pts/0
                                        1365438 S
                                                   24908
                                                             0:00 sleep 20
1363886 1365309 1365308 1357688 pts/0
                                                     24908
                                                             0:00 sleep 20
                                        1365438 5
1363886 1365310 1365308 1357688 pts/0
                                                             0:00 sleep 20
                                        1365438 S
                                                     24908
```

Notice the PID and PGID!

#### **POSIX Spawn**

- Replaces fork() + exec() entirely
- Code is "linear" rather than handling multiple processes in if-else statements
- posix\_spawnattr\_t and posix\_spawn\_file\_actions\_t are structs that store information about process groups and I/O redirection/piping respectively. These structs don't do anything until posix\_spawnp is used.
- Example: posix spawn(3) Linux manual page (man7.org)

Note: You need to include "spawn.h" in your cush.c to use these functions. The file is located in the posix\_spawn directory. Also be sure to use the "make" command to compile posix\_spawn.

# fork() + exec()

```
if (fork() == 0) {
    //child stuff

    execvp(/* program arguments */);
}
else {
    //parentstuff
}
```

## posix\_spawn()

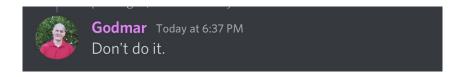
```
posix_spawn_file_actions_t child_file_attr;
posix_spawnattr_t child_spawn_attr;

posix_spawnattr_int(&child_file_attr);
posix_spawn_file_actions_init(&child_file_attr);

// setup for attributes

posix_spawnp(/*pid*/, /*program*/, &child_file_attr,
&child_spawn_attr, /*program arguments*/, environ)
```

You can use fork() + exec() for this project, but our recommendation is:



## **POSIX Spawn Attributes**

- Process Groups posix\_spawnattr\_setpgroup()
- Terminal Control posix\_spawnattr\_tcsetpgrp\_np()
- Piping posix\_spawn\_file\_actions\_adddup2()
- I/O Redirection posix\_spawn\_file\_actions\_addopen()

More listed on both the spec and <spawn.h>.

#### **Built-in Commands**

- Commands that are defined within the program by you
  - No need to fork off and execute an external program
- Required Built-In Commands for your shell:
  - o kill kills a process
  - jobs displays a list of jobs
  - stop stops a process
  - fg sends a process to foreground
  - bg sends a process to background
  - exit exits the shell
- Built-in Commands are not considered Jobs
- Two additional built-ins / functionality extenders also required (examples in later slide)
  - One low-effort
  - One high-effort

#### **Built-ins Behind the Scenes**

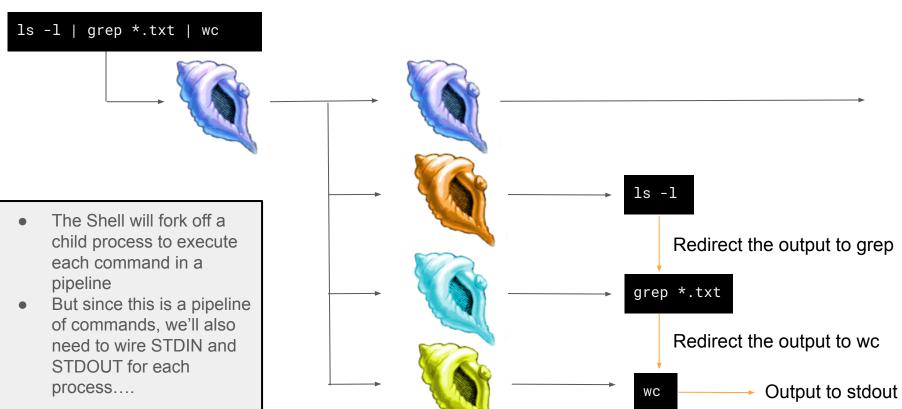
# \$ jobs

#### FOUR STEPS for built-in

- 1. Shell waits for user input
- 2. Shell realizes this is a built in command
- 3. Shell executes built-in (no forking)
- 4. After execution, shell repeats

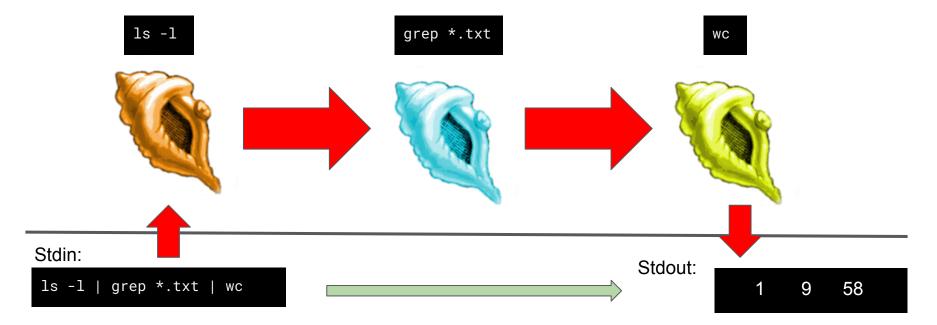
[1]+ Stopped vim
[2]- Running sleep 20 &

#### I/O Piping



#### I/O Piping

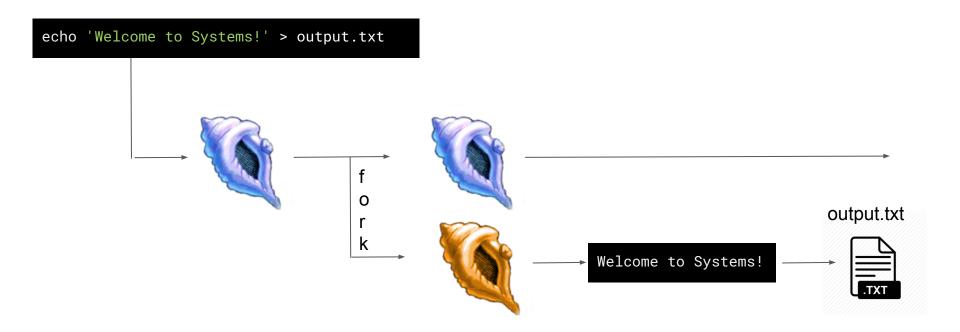
- Processes will wait on previous process, final process outputs to terminal
- STDIN and STDOUT for processes are joined to create the pipeline



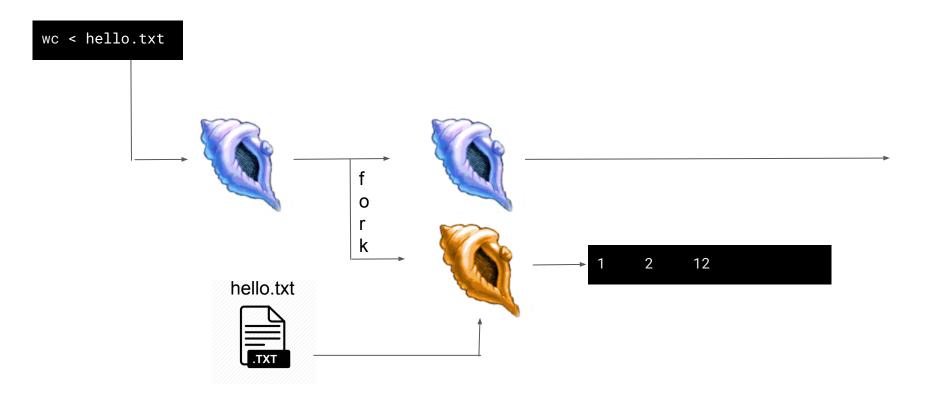
#### I/O Redirection

- overwrites original file contents before writing the new output
- >> appends new content to the end of the original file
- read input from a file rather than STDIN

#### I/O Redirection (Output)



## I/O Redirection (Input)



#### I/O Redirection (Stderr)

Contents written to STDERR can also be piped into other processes using |& and outputted to files using >&.

```
int main() {
    fprintf(stderr, "Write to stderr.\n");
    fprintf(stdout, "Write to stdout\n");
}
```

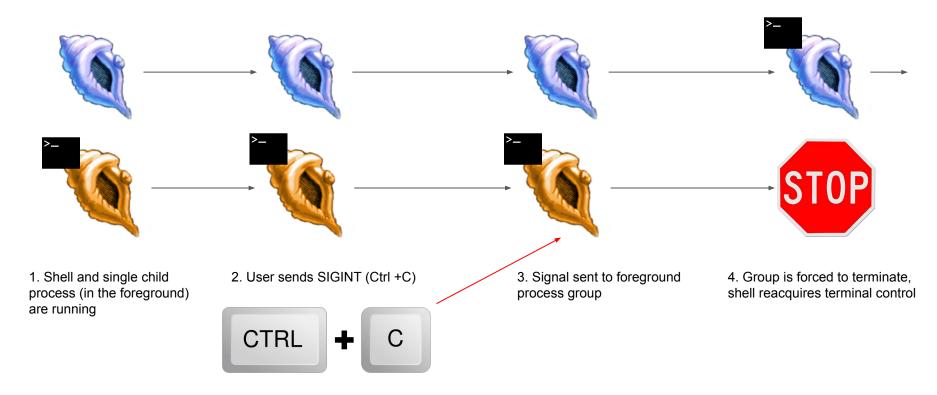
Notice how the message "Write to stderr." was not outputted.

#### **Signal Handling**

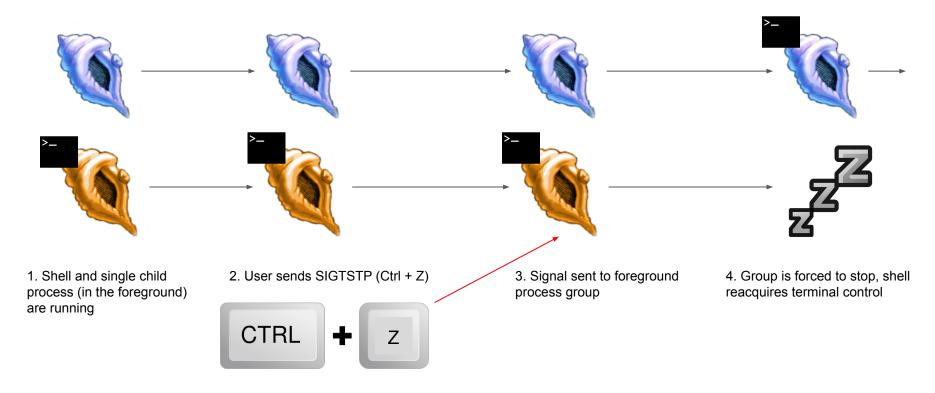
- Shells can handle signals sent to them
  - SIGINT (Ctrl + C)
  - SIGTSTP (Ctrl + Z)
  - SIGCHLD (when a child process terminates)

 Most of the functionality of this will be done in handle\_child\_status(pid\_t pid, int status)

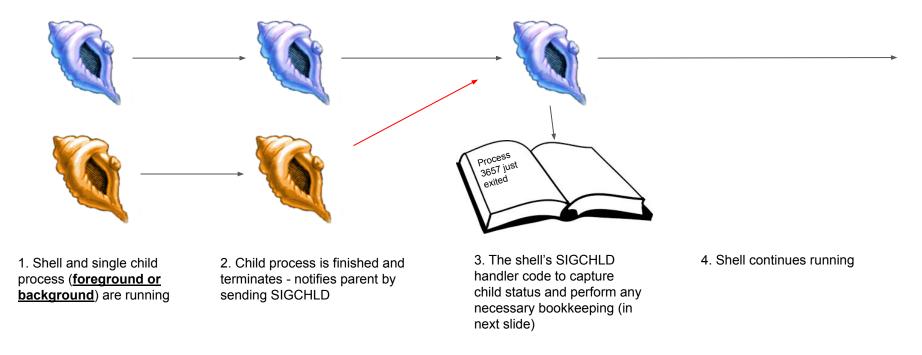
# **Handling SIGINT (Ctrl + C)**



# **Handling SIGTSTP (Ctrl + Z)**



# **Handling SIGCHLD**



## Handling SIGCHLD: WIF\* Macros

- When wait\* is called it will return a pid and a status for a child process that changes state. Using macros, we can decode this status to discover what state a process changed to and how it happened:
  - WIFEXITED(status) did child process exit normally?
  - WIFSIGNALED(status) was child process signaled to terminate?
  - WIFSTOPPED(status) was child process signaled to stop?

process exits via exit()	WIFEXITED	WEXITSTATUS has return code	no	yes
user terminates pro- cess with Ctrl-C	WIFSIGNALED	WTERMSIG equals SIGINT	no	yes
user terminates pro- cess with kill	WIFSIGNALED	WTERMSIG equals SIGTERM	no	yes
user terminates pro- cess with kill -9	WIFSIGNALED	WTERMSIG equals SIGKILL	no	yes
process has been terminated (general case)	WIFSIGNALED	WTERMSIG equals signal number	no	yes
	c/manual/html_n	GNU C library manual, a ode/index.html. Real.		20 C C C C C C C C C C C C C C C C C C C

Additional info

SIGTSTP

SIGSTOP

TIN

WSTOPSIG

WSTOPSIG equals

WSTOPSIG equals

SIGTTOU or SIGT-

equals

Process

yes

yes

yes

stopped?

Process

dead?

no

no

no

How to check for it

WIFSTOPPED

WIFSTOPPED

WIFSTOPPED

Event

User stops fg pro-

User stops process

wants

cess with Ctrl-Z

with kill -STOP

non-foreground

terminal access

process



# **Project Overview**

# Requirements and Grading

- 1. Basic Functionality 50 pts
  - a. Start foreground and background jobs
  - b. Built-in commands: 'jobs', 'fg', 'bg', 'kill', 'stop'
  - c. Signal Handling (SIGINT, SIGTSTP, SIGCHLD)
- Advanced Functionality 50 pts
  - a. I/O Piping
  - b. I/O Redirection
  - c. Running programs requiring exclusive terminal access (ex: vim)
- 3. Two Extra Built-ins 20 pts
  - a. One low effort
  - b. One high effort
- Version Control (git) 10 pts
  - a. At least 3 commits per partner
- 5. Documentation 10 pts

#### Total: 140 points

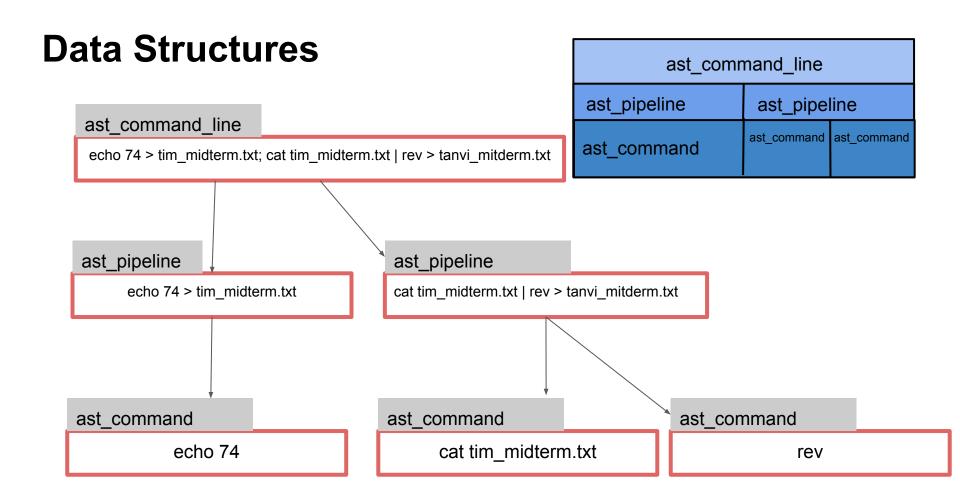
## **Before You Start Coding ....**

- Take time to read over, understand the spec and the starter code
- Read the provided lecture material
- Understand Exercise 1
  - fork() / exec() model
  - Piping : pipe(), dup2(), close()
- Check out Dr. Back's example shell
  - Located at ~cs3214/bin/cush-gback in rlogin
  - Can be useful for comparing outputs with your shell

Super Helpful!

#### **Base Code**

- Already includes a parser!
- Parser spits out hierarchical data structures



#### **List Data Structure**

- You're also provided with a doubly linked list data structure
  - Check out list.h and list.c
- You'll be using this list throughout the semester
- Read through list.h before using it

#### "Data contains node" vs "Node points to data"

#### **Our Linked List**

```
struct list_elem {
    struct list_elem * prev;
    struct list_elem * next;
}
```

#### A Regular Linked List

```
class listnode<T> {
        T data;
        listnode <T> prev;
        listnode<T> next;
}
```

#### "Data contains node" vs "Node points to data"

#### **Our Linked List** struct list elem { struct list elem \* prev; struct list elem \* next; Struct 1 Struct 2 Other Other fields fields list elem list elem Sentinel Sentinel

```
A Regular Linked List
class listnode<T> {
      T data;
      listnode <T> prev;
      listnode<T> next;
Sentinel
               Node
                             Node
                                           Sentinel
               Data
                              Data
```

# So how do I get my data?

struct ast\_command \* cmd = list\_entry(e, struct ast\_command, elem);

Retrieve data from a struct list elem by using the list entry macro:



#### Colin McGee (UTA) 09/11/2023 10:01 PM

A struct list elem is an element of the struct. When you do list operations, it works on this list elem. However, you probably want to get a reference to the struct that the elem is contained in, rather than just the element, right? list entry(elem pointer, struct type, name of list elem in struct) is a way to convert between a struct list elem\* to a pointer to the struct that contains it.

Beautiful explanation by one of our UTA's:)





#### **List Pitfalls**

#### • Don't:

- Use the same list\_elem for multiple lists
- Edit an element while iterating
  - Naive loop to remove elements in a list will fail!
- Forget to list init()

#### BAD IDEA:(

```
// valid example: deallocates a pipeline struct and any commands stored in it while iterating
void ast_pipeline_free(struct ast_pipeline *pipe)
{
    for (struct list_elem * e = list_begin(&pipe->commands); e != list_end(&pipe->commands); ) {
        struct ast_command *cmd = list_entry(e, struct ast_command, elem);
        e = list_remove(e); //Acts as the iterator; stores next element into e
        ast_command_free(cmd);
    }
    free(pipe);
} // make sure to remove an ast_pipeline from a list before adding it to another!
// bottom line with lists? ALWAYS TEST
```

## **Utility Functions (Strongly Recommended)**

- Signal Support (signal\_support.c / .h)
  - signal\_block()
  - o signal\_unblock()
  - singal\_set\_handler()
- Terminal State Management (termstate\_management.c / .h)
  - o termstate init()
  - o termstate\_give\_terminal\_to()
  - termstate\_give\_terminal\_back\_to\_shell()
  - termstate\_get\_current\_terminal\_owner()
  - o termstate\_save()
  - termstate\_restore()

#### Additional Built-ins and extensions

- Your shell must contain two extra built-ins / functionality extensions
  - One high effort and one low effort (bolded is low-effort)
- Ideas include:
  - Customizable Prompt
  - Setting/unsetting env vars
  - Implementing the 'cd' built-in
  - Glob expansion (e.g., \*.c)
  - Timing commands (ex. time)
  - Alias support

- Shell Variables
- Directory Stack
- Command-line history
- Backquote substitution
- Smart command-line completion
- Embedded Apps
- If you have an idea not shown on the list or have any doubts please ask us

## **Testing / Submission**

- Test the driver before submitting, don't just run tests individually
- When grading, tests will be ran 3-5 times. If you crash a single time, it's considered failing
- Make sure you don't have undefined behavior by checking the system call return code and using valgrind to address memory related issues

#### **Test Driver**

- The driver reads from .tst file that describes a test suite (ex. basic.tst)
  - Ex: basic.tst contains a series of test scripts that it will run from the folder /tests/basic

#### cd src/

../tests/stdriver.py [options]

\*- stdriver.py also available at ~cs3214/bin/stdriver.py

#### Options:

- -b : basic tests (processes, built-ins, signals)
- -a: advanced tests (I/O Piping, I/O Redirection, exclusive terminal access)
- -h : list all the options

#### **Additional Tests**

- You are required to write tests for your two extra built-ins
  - Create a .tst file in 'tests' and create a directory that will store your test scripts
- Inside <custom>.tst file:

```
= <custom> Tests
pts <custom>/<test_name>.py
pts <custom>/<test_name>.py
...
```

```
= Milestone Tests
1 basic/foreground.py
1 basic/cmdfail_and_exit_test.py
```

 The driver checks number of total points (pts) to use for a test. Since this is just your own custom tests you can put an arbitrary points here

## **Additional Tests (Part 2)**

Make sure your custom.tst file is of type "ASCII text"



- If it includes Windows line terminators (CR, CRLF, etc) it will fail
- We want \n, not \r\n

## **Design Document**

- When you submit you must include a README.txt describing your implementation of P1
- Explain the custom built-ins created and approach taken to develop them.
- TAs will assign credit only for the functionality for which test cases and documentation exist

Submission. You must submit a design document, README.txt, as an ASCII document using the following format to describe your implementation: Student Information <Student 1 Information> <Student 2 Information> How to execute the shell <describe how to execute from the command line> Important Notes <Anv important notes about your system> Description of Base Functionality <describe your IMPLEMENTATION of the following commands:</pre> jobs, fg, bg, kill, stop, \^C, \^Z > Description of Extended Functionality <describe your IMPLEMENTATION of the following functionality:</pre> I/O, Pipes, Exclusive Access >

List of Additional Builtins Implemented

(Written by Your Team)

<br/>
<br/>
<br/>
<description>

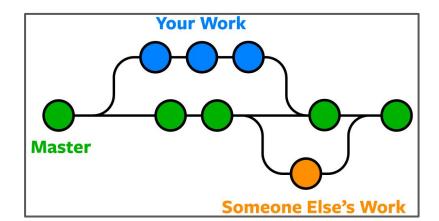


## **Version Control**

#### **Version Control**

- You will be using Git and Gitlab for managing your source code
- Why?
  - Organizes your code
  - Keeps track of features
  - Allows collaborators to work freely without messing up other existing code
  - Back-ups whenever something goes wrong





#### **Basic Git Commands**

• Stage file for commit:

\$ git add <file\_name>

Commit files:

\$ git commit -m 'Add a description here'

Push changes to remote (note: always pull before push!)

\$ git push [origin <branch\_name>]

### **Basic Git Commands**

Fetch changes from remote:

#### \$ git pull

Check status:

#### \$ git status

Revert to the previous commit:

```
$ git reset [--hard]
```

#### **Basic Git Commands**

Create a new branch from the current branch:

\$ git checkout -b <new\_branch\_name>

Switch to another branch:

\$ git checkout <branch\_name>

Merge a branch into the current branch

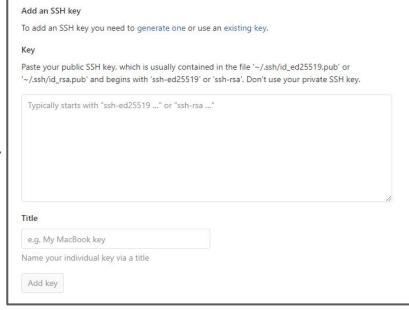
\$ git merge <branch\_name>

## **Setup Git Access**

- You'll need an SSH Key to get access to projects at git.cs.vt.edu
- If you don't already have a key...
  - Create a new key:

```
$ ssh-keygen -t rsa -b 4096 -C "email@vt.edu" \
-f ~/.ssh/id_rsa
```

- Add Key to <a href="https://git.cs.vt.edu/profile/keys">https://git.cs.vt.edu/profile/keys</a>
  - You will paste public key here ----->



## **Verify Git Access**

- Verify you have access
- The first time you connect you will be asked to verify the host, just answer 'Yes' to continue

11 spencetk@linden ~>ssh git@git.cs.vt.edu

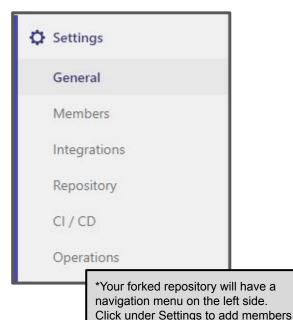
PTY allocation request failed on channel 0
Welcome to GitLab, @spencetk! ← Your pid should be displayed here
Connection to git.cs.vt.edu closed.

- You can get in-depth explanations here:
  - Generate a key
  - Use an existing key

## GitLab Project Setup

- One member will fork the base repository:
  - O https://git.cs.vt.edu/cs3214-staff/cs3214-cush
- 2. Invite partner to collaborate
  - O Go to Settings > Members to add them
  - O Check partner role permissions too
- 3. Both members will clone the forked repository on their machines:





and set repo to private

**IMPORTANT:** Set forked repository to **private** 

Go to Settings > General > Visibility, project features, permissions



# The GNU Project Debugger (GDB)



## **Starting GDB**

1. Invoke GDB with a program:

\$ gdb program

2. Run with command

(gdb) run arg1 arg2

Must be compiled with debug symbols, -g

## **Breakpoints**

Set a breakpoint

```
(gdb) b <func_name> OR (gdb) b line_number>
```

Set a conditional breakpoint:

```
(gdb) b <func_name> if <condition>
```

Ignore breakpoint #1 100 times

```
(gdb) ignore 1 100
```

Show # of times breakpoint was hit

(gdb) info b

### **Backtrace and Frames**

Show backtrace:

#### (gdb) backtrace

- Show frame:
  - After selecting frame, you can print all variables declared in that function call

(gdb) frame <num>

#### Follow-Fork-Mode

Which process to follow after a fork (parent / child):

#### (gdb) set follow-fork-mode <mode>

- o 'parent' = ignore child process and continue debugging the parent
- 'child' = begin debugging the child process when fork() is called
- Retaining debugger control after fork:
  - After a fork, specify whether to freeze the child or allow it to run (this may make it difficult to find race conditions)

(gdb) set detach-on-fork <mode>



# Valgrind

#### 1. Start valgrind

\$valgrind --log-file="output.log" --leak-check=full ./cush

#### 2. Run commands in your shell

cush@cedar in src> jobs cush@cedar in src> ls

#### 3. Check the log file

```
=50520== Memcheck, a memory error detector
=50520== Copyright (C) 2002-2017, and GNU GPL'd, by Julian Seward et al.
=50520== Using Valgrind-3.15.0 and LibVEX; rerun with -h for copyright info
=50520== Command: ./bru.exe
=50520==
50520== HEAP SUMMARY:
            in use at exit: 22,592 bytes in 164 blocks
         total heap usage: 185 allocs, 21 frees, 31,040 bytes allocated
=50520===
=50520== LEAK SUMMARY:
           definitely lost: 3,584 bytes in 56 blocks
=50520===
           indirectly lost: 0 bytes in 0 blocks
             possibly lost: 72 bytes in 3 blocks
50520==
           still reachable: 200 bytes in 6 blocks
=50520==
                suppressed: 18,736 bytes in 99 blocks
=50520== Rerun with --leak-check=full to see details of leaked memory
=50520== For lists of detected and suppressed errors, rerun with: -s
=50520== ERROR SUMMARY: 0 errors from 0 contexts (suppressed: 4 from 4)
```



## Advice

### **Advice**

- START EARLY
- Create a roadmap before starting projects
- Utilize TAs
  - Come with questions prepared, try to figure out what the problem is first
  - **Be organized and have clean code** the cleaner it is, the faster we can help!
  - Run valgrind and try debugging with GDB before consulting us
  - Discord, Zoom, Class Forum
- Understand the Exercises
- Use valgrind! This can isolate many bugs
- Become an expert at the debugger
- Find what works best for communicating with your partner
  - o In-Person Meetings, Discord, Zoom, etc.

#### Sources

- Referred to previous help session slides created by previous UTA's Kent McDonough, Connor Shugg, Joe D'Anna, Chris Cerne, Justin Vita, Sam Lightfoot, and Alex Kyer, Timothy Wu, Tanvi Allada since the Spring 2021 Semester
- Spencer Keefer created the revised slides



# Thanks for attending! Questions?