

# CS 3214: Computer Systems

## Lecture 7: Signals

Instructor: Huaicheng Li

Sept 13 2022



VIRGINIA TECH™

# Signals

- ❑ A small message to notify a process of an event
- ❑ Similar to exceptions and interrupts
- ❑ Who generates signals?
  - Self, other-processes, the kernel
- ❑ Signal types (integer ID's, e.g., <32)

<i>ID</i>	<i>Name</i>	<i>Default Action</i>	<i>Corresponding Event</i>
2	SIGINT	Terminate	User typed ctrl-c
9	SIGKILL	Terminate	Kill program (cannot override or ignore)
11	SIGSEGV	Terminate	Segmentation violation
14	SIGALRM	Terminate	Timer signal
17	SIGCHLD	Ignore	Child stopped or terminated

# Synchronous Signals

- ❑ **SIGILL (1)** Illegal Instruction
- SIGABRT (1)** Program called abort()
- SIGFPE (1)** Floating Point Exception (e.g. integer division by zero)
- SIGSEGV (1)** Segmentation Fault - catch all for memory and privilege violations
- SIGPIPE (1)** Broken Pipe - attempt to write to a closed pipe
- SIGTTIN (2)** Terminal input - attempt to read from terminal while in background
- SIGTTOU (2)** Terminal output - attempt to write to terminal while in background

(1) Default action: terminate the process

(2) Default action: stop the process

# Asynchronous Signals

- **SIGINT (1, 3)**      Interrupt: user typed Ctrl-C
- SIGQUIT (1, 3)**    Interrupt: user typed Ctrl-\
- SIGTERM (3)**        User typed kill pid (default)
- SIGKILL (2, 3)**     User typed kill -9 pid (urgent)
- SIGALRM (1, 3)**    An alarm timer went off (alarm(2))
- SIGCHLD (1)**        A child process terminated or was stopped
- SIGTSTP (1)**        Terminal stop: user typed Ctrl-Z
- SIGSTOP (2)**        User typed kill -STOP pid

(1) These are sent by the kernel, e.g., terminal device driver

(2) SIGKILL and SIGSTOP cannot be caught or ignored

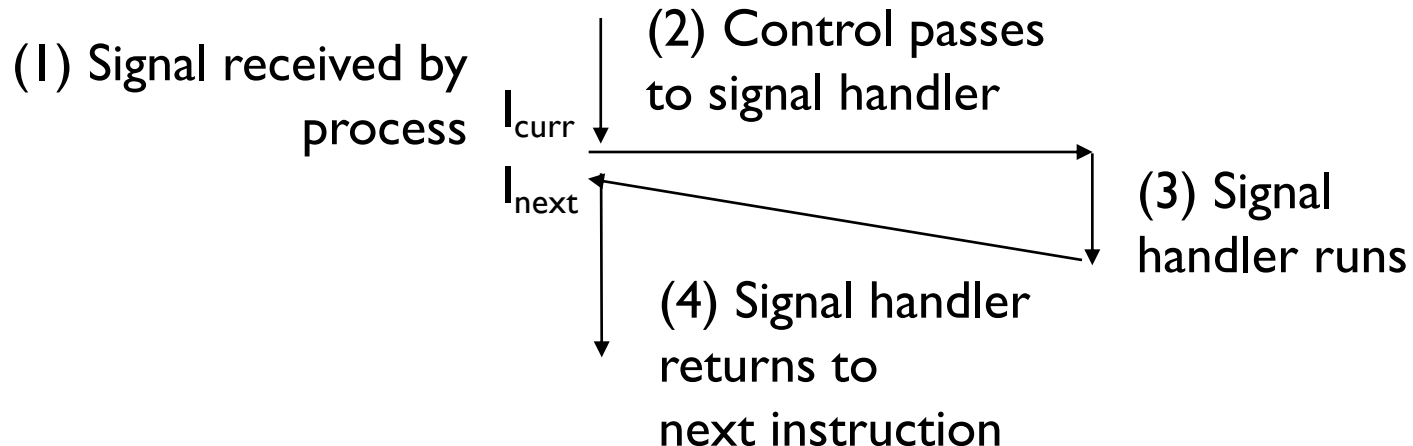
(3) Default action: terminate the process

# Sending a Signal

- ❑ Kernel sends a signal to a destination process by updating some state in the context of the destination process
  - divide-by-zero (SIGFPE)
  - Termination of a child process (SIGCHLD)
- ❑ Another process has invoked *kill()* system call to explicitly request the kernel to send a signal to the destination process
- ❑ *raise()*

# Receiving a Signal

- A destination process receives a signal when it is forced by the kernel to react in some way to the delivery of the signal
- Possible reactions
  - Ignore the signal (do nothing)
  - Terminate the process (e.g., with core dump)
  - Catch the signal by executing a user-level function called signal handler



# Pending and Blocked Signals

- ❑ Pending: sent but not yet received
  - At most one pending signal of any particular type
  - Signals are not queued (On/Off)
- ❑ A process can block the receipt of certain signals
  - Blocked signal can be delivered, but will not be received until the signal is unblocked
- ❑ A pending signal is received at most once
- ❑ Kernel maintains pending and blocked bit vectors in the context of each process
  - Pending: kernel sets/clears certain bits when a signal is delivered/received
  - Blocked: *sigprocmask()*, aka, *signal mask*

# Safe Signal Handling

```
void
list_insert (struct list_elem *before,
             struct list_elem *elem)
{
    elem->prev = before->prev;
    elem->next = before;
    before->prev->next = elem;
    before->prev = elem;
}
```

```
list_insert:
    movq    (%rdi), %rax
    movq    %rdi, 8(%rsi)
    movq    %rax, (%rsi)
    movq    %rsi, 8(%rax)
    movq    %rsi, (%rdi)
    ret
```

If a signal arrives in the middle of `list_insert()`, the manipulated list's list element are in a partially linked state. If the signal handler now takes a path where the same list is being accessed (iterated over, etc.), inconsistent behavior will result. This situation must be avoided.



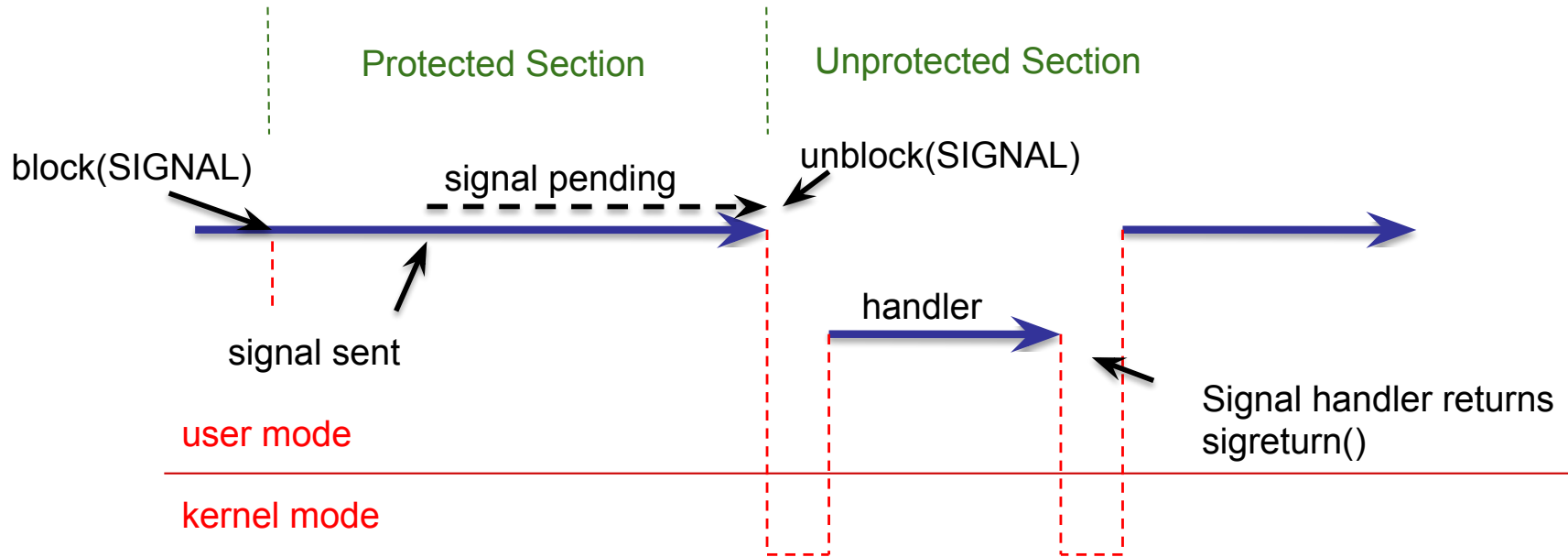
# Safe Signal Handling

- ❑ Concurrent with main program
- ❑ Guidelines to avoid trouble
  - Keep handlers simple
  - Only use async-signal-safe functions (no *printf*)
  - Save and restore *errno* on entry and exit to avoid overwrite
  - Temporarily blocking all signals to protect access to shared data structures
  - Declare global variables as *volatile* to prevent compiler from storing them in a register
  - Declare global flags as *volatile sig\_atomic\_t*

# Async-Signal-Safety

- ❑ Man 7 signal
- ❑ Safe: `_exit()`, `write()`, `wait()`, `waitpid()`, `sleep()`, `kill()`
- ❑ Unsafe: `printf()`, `sprint()`, `malloc()`, `exit()`

# Blocking/Unblocking Signals



*If signals are masked/blocked most of the time in the main program, signal handlers can call most functions, but signal delivery may be delayed. If a signal is not masked most of the time, signal handlers must be very carefully implemented. In practice, coarse-grained solutions are perfectly acceptable unless there is a requirement that bounds the maximum allowed latency in which to react to a signal. Side note: OS face the same trade-off when implementing (hardware) interrupt handlers.*

# Blocking/Unblocking Signals

❑ Explicit blocking/unblocking: *sigprocmask()*

❑ *Others*

- *sigemptyset()* – create empty set
- *sigfillset()* – Add every signal number to set
- *sigaddset()* – Add signal number to set
- *sigdelset()* – Delete signal number from set

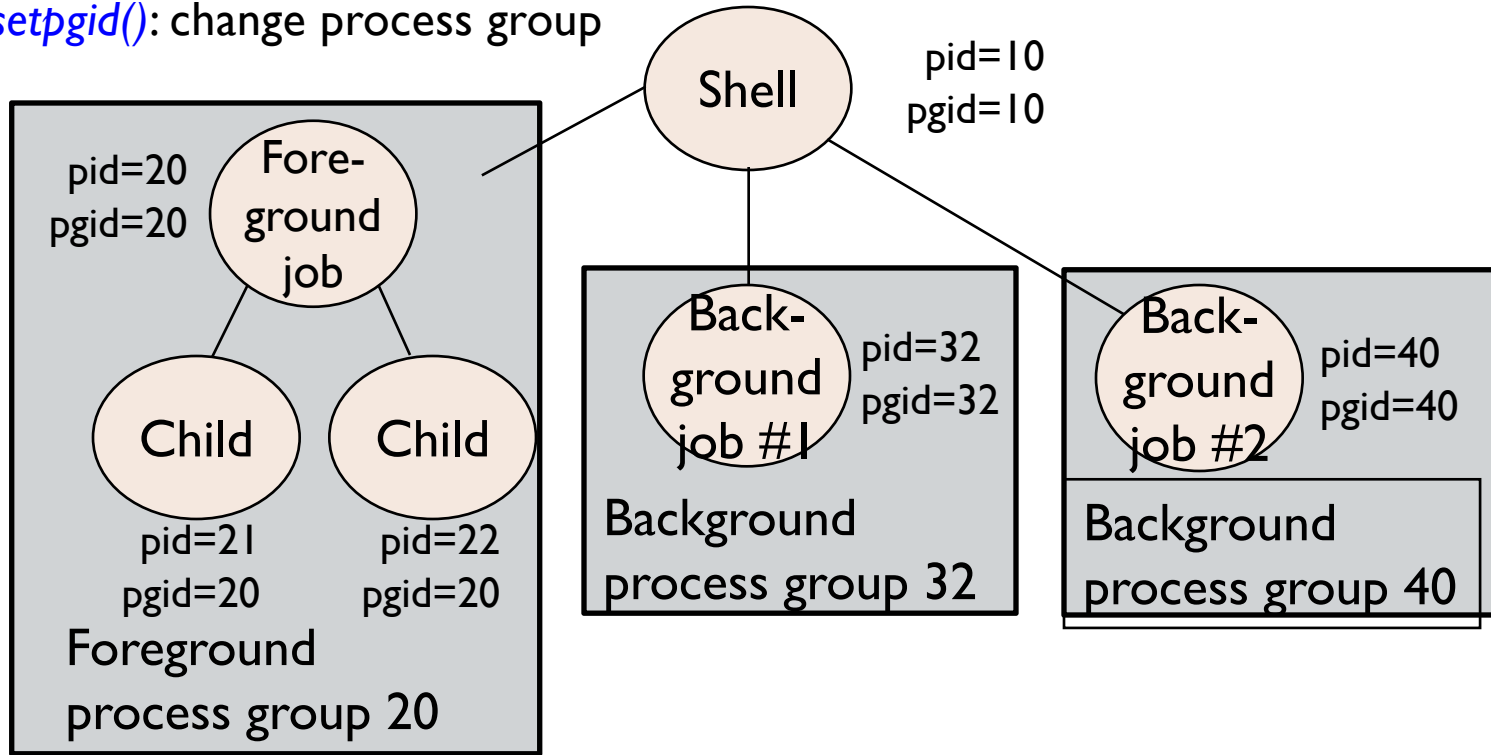
```
sigset_t mask, prev_mask;  
sigemptyset(&mask);  
sigaddset(&mask, SIGINT);
```

```
/* Block SIGINT and save previous blocked set */  
sigprocmask(SIG_BLOCK, &mask, &prev_mask);
```

```
/* Code region that will not be interrupted by SIGINT */  
/* Restore previous blocked set, unblocking SIGINT */  
sigprocmask(SIG_SETMASK, &prev_mask, NULL);
```

# Process Group

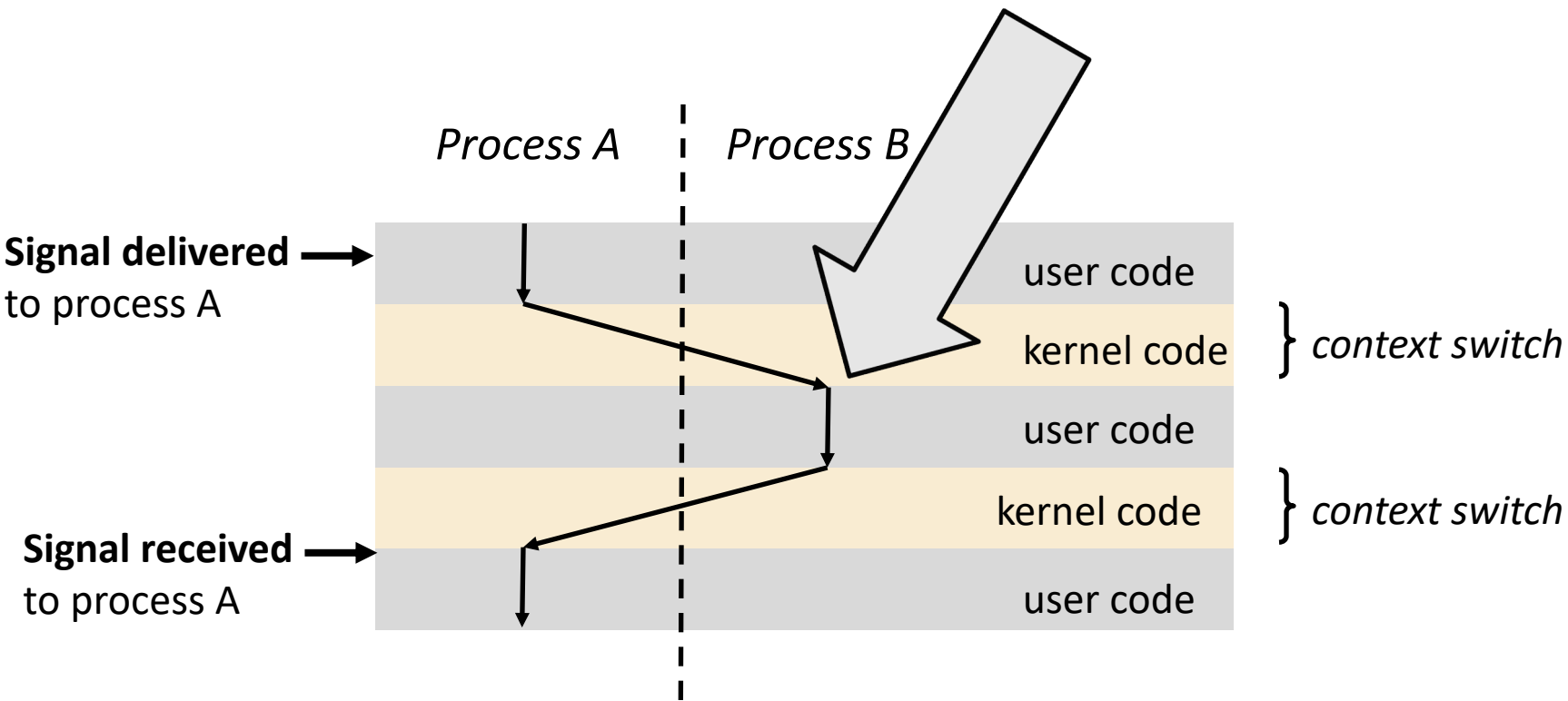
- One process belongs to one process group
  - `getpgrp()`, get process group of current process
  - `setpgid()`: change process group



# Kill: sending signals

- ❑ Kill -9 1000: Send SIGKILL to process 1000
- ❑ Kill -9 -1000: Send SIGKILL to every process in process group 1000
- ❑ Ctrl-C: SIGINT
- ❑ Ctrl-Z: SIGTSTP

# Receiving Signals



# Signal APIs

- Uniform APIs for programs to determine actions to be taken for signals
  - Terminating the process, core dump
  - Ignoring the signal
  - Invoking a user-defined handler
  - Stop the process
  - Continuing the process



# Installing Signal Handlers

□ `Handler_t *signal(int signum, handler_t *handler)`

```
void sigint_handler(int sig) /* SIGINT handler */
{
    printf("So you think you can stop the bomb with ctrl-c, do you?\n");
    sleep(2);
    printf("Well...");
    exit(0);
}

int main()
{
    /* Install the SIGINT handler */
    if (signal(SIGINT, sigint_handler) == SIG_ERR)
        unix_error("signal error");
    /* Wait for the receipt of a signal */
    pause();
    return 0;
}
```