The first step in design is to understand the problem.

What does your assembler have to do?

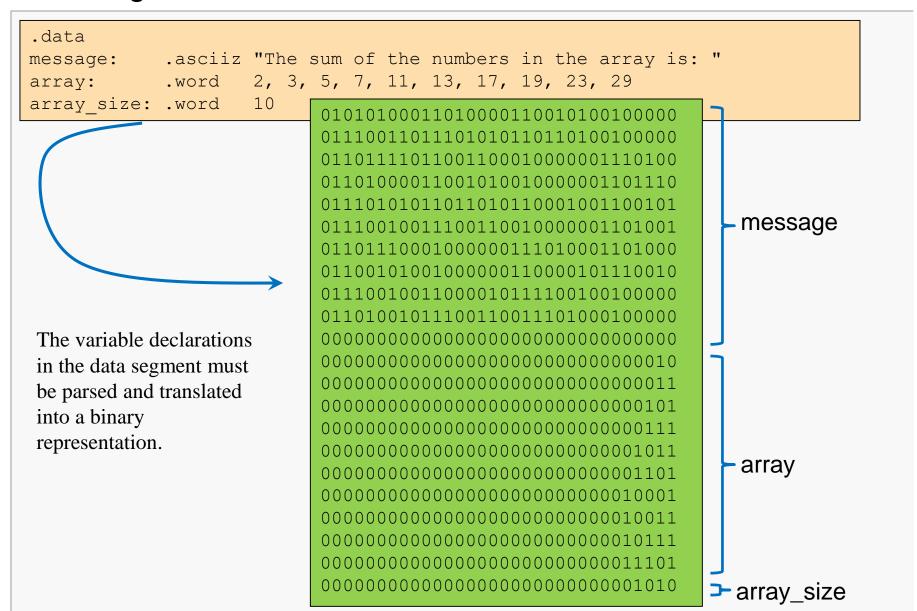
- for the data segment...?
- for the text segment...?

What information does your assembler have to possess?

- for the data segment...?
- for the text segment...?

How are you going to organize that information?

Data Segment



```
.text
main:
          $a0, array
       la
       la $a1, array size
       lw $a1, 0($a1)
loop:
       sll $t1, $t0, 2
       add $t2, $a0, $t1
       sw $t0, 0($t2)
       addi $t0, $t0, 1
       add $t4, $t4, $t0
       slt $t3, $t0, $a1
       bne
            $t3, $zero, loop
       li $v0, 4
       la $a0, message
       syscall
       li $v0, 1
       or $a0, $t4, $zero
       syscall
       li $v0, 10
       syscall
```

The assembly instructions in the text segment must be parsed and translated into a binary representation.

Consider an Example

Keep it simple.

How would YOU translate a particular MIPS assembly instruction to machine code?

Consider: add \$t0, \$s5, \$s3

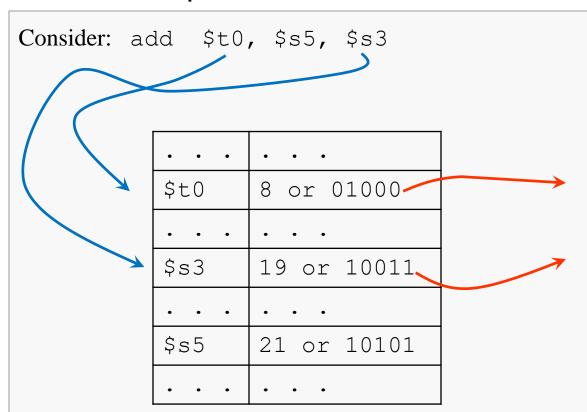
What's the machine code format? (R-type, I-type, J-type, special?)

How do you know that?

What are the correct values for the various fields in the machine instruction?

How do you know that?

More to the point... how will your program "know" those things?



Think of the table as defining a mapping from some sort key of value (e.g., symbolic register name) to another sort of value (e.g., register number, binary text string).

What are the key values?

What are the values we want to map the keys to?

Implementing a Table

Define a struct type that associates a particular key value with other values; for instance:

```
struct _RegMapping {      // register name to number
  char* regName; // symbolic name as C-string
  char* regNumber; // string for binary representation
};
typedef struct RegMapping RegMapping;
```

Define an array of those, and initialize appropriately; for instance:

```
static RegMapping Table[...] = {
   {"$zero", "00000"},
   {"$at", "00001"},
  {"$t0", "01000"},
   {"$ra", "11111"}
};
```

Define a function to manage the lookup you need and you're in business...

Static Tables

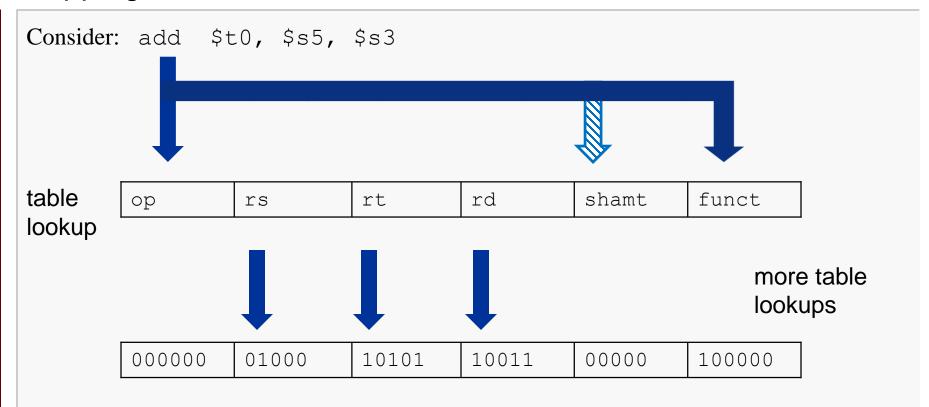
Put the initialization of your table at file scope in an appropriate .c file:

```
// Register lookup module .c file
struct RegMapping {
};
typedef struct RegMapping RegMapping;
static RegMapping Table[...] = {
};
// definitions of lookup functions
```

Now:

- the table is created automatically when your program starts
- the table exists the whole time your program is running
- the table can only be accessed by calling the functions you "publish" in a .h file

Mapping Fields to Bits



If we have the right tables and we break the assembly instruction into its parts, it's easy to generate the machine instruction...

Representing the Machine Instruction

One basic design decision is how to represent various things in the solution.

For the machine instruction, we have (at least) two options:

```
char MI[...]; // array of chars '0' and '1'
uint32 t MI;
           // sequence of actual bits
```

Either will work.

Each has advantages and disadvantages.

But the option you choose will affect things all throughout the design... so decide early!

Either way, you have to decide how to put the right bits at the right place in your representation of the machine instruction.



Alas, C does not provide any standard format specifiers (or some other feature) for displaying the bits of a value. But, we can always roll our own:

It would be fairly trivial to modify this to print the bits of "wider" C types.

It would also be easy to modify this to put the characters into an array...

But, execution of the assembler starts with an assembly program file, like:

```
.data
                                             data segment
Str01: .asciiz "To be or not to be..."
        .text
main:
        la $t0, Str01
        li $s0, 4096
        add $s0, $s1, $s2
                                             text segment
bgloop:
        lw $t1, ($t0)
        beg $t0, $t7, bgloop
        li $v0, 10
        syscall
```

The logic of parsing is different for the data segment and the text segment.

So is the logic of translation to text-binary form.

How are you going to handle the high-level tasks of identifying instructions/variables?

Doing this by hand, you'd probably think of grabbing a line at a time and processing it.

```
safely read a
          .text
                                                       line of text
main:
         la $t0, Str01
         li $s0, 4096
         add $s0, $s1, $s2
                                                                read
bgloop:
                                                                formatted
         lw $t1, ($t0)
                                                                values from
                                                                a C-string
         li $v0, 10
         syscall
                                                                break a C-
C provides a number of useful library functions:
                                                                string into
```

break a Cstring into delimited pieces, destructively

Parsing an Assembly Instruction

Consider: add \$t0, \$s5, \$s3

The specification says some things about the formatting of assembly instructions.

Those things will largely determine how you split an instruction into its parts.

And, don't forget that different instructions take different numbers and kinds of parameters:

How would we figure and out the number ameters kinds of parameters at runtime?

```
la $t0, Str01
li $s0, 4096
...
add $s0, $s1, $s2
bgloop:
lw $t1, ($t0)
...
syscall
```

Parsing an Assembly Instruction

```
Consider: add $t0, $s5, $s3
```

C provides a number of useful functions here.

Do not ignore strtok() ... it's flexible and powerful.

But also, don't ignore the fact that C supports reading formatted I/O:

```
char* array = malloc(MAXLINELENGTH);
. . .
fgets(array, MAXLINELENGTH, source);
. . .
// determine that you read an instruction taking 3 reg's
. . .
sscanf(array, " %s %s, %s, %s", . . .);
```

But... I'm ignoring the return value from sscanf()... that may not be best practice.

Broader Perspective: Labels

We may find labels in the data segment and/or the text segment. .data data segment Str01: .asciiz "To be or not to be..." .text main: la \$t0, Str01 li \$s0, 4096 add \$s0, \$s1, \$s2 bgloop:) lw \$t1, (\$t0) text segment beq \$t0, \$t7, bgloop bne \$t7, \$s4, done done: li \$v0, 10 syscall

```
What's the deal?

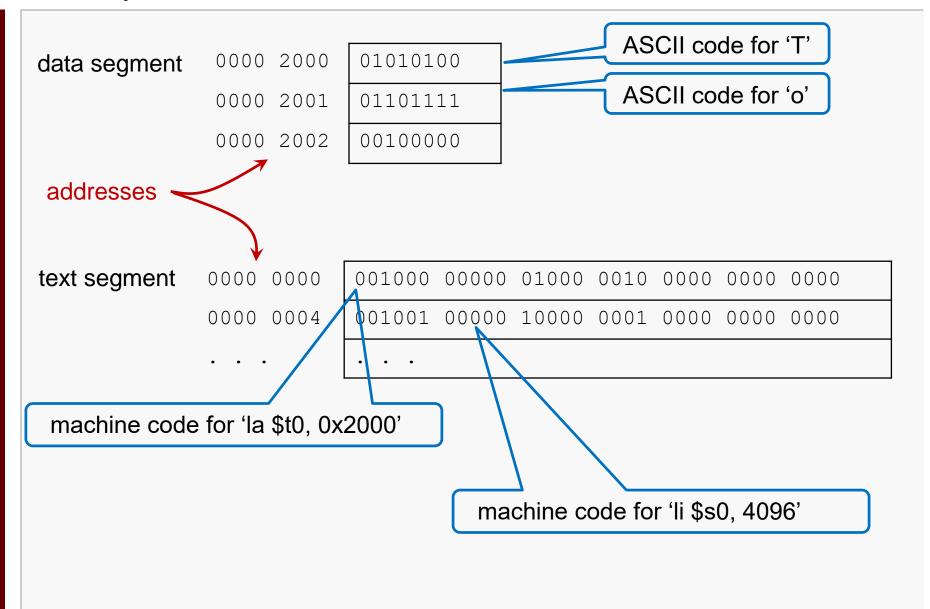
.data
Str01: .asciiz "To be or not to be..."

.text
main:
la actually translates to an addi instruction

addi $t0, $zero, Str01
```

```
What's the deal?
         .data
Str01: .asciiz "To be or not to be..."
                                                     la actually translates
         .text
                                                     to an addi instruction
main:
         la $t0, Str01
                                             addi
                                                     $t0, $zero, Str01
         addi
                $t0, $zero, <address>
                                                  Labels translate to 16-bit
                                                  addresses... how?
```

```
.data
Str01: .asciiz "To be or not to be..."
        .text
main:
        la $t0, Str01
        li $s0, 4096
data segment 0000 2000
                        "To be or ..."
text segment 0000 0000
                        la
                            $t0, 0x2000
             0000 0004
                        li $s0, 4096
```



Symbol Table

The assembler needs to build a *symbol table*, a table that maps symbolic names (labels) to memory addresses:

0000	0000	main
0000	001C	bgloop
0000	2000	Str01

Building the symbol table is a bit tricky:

- need to know where data/text segment starts in memory
- may see a label in an instruction before we actually see the label "defined"

One reason most assemblers/compilers make more than one pass.

We want the symbol table to be built before we start translating assembly instructions to machine code — or else we must do some fancy bookkeeping.

Incremental Development

Plan your development so that you add features one by one.

This requires thinking about (at least) two things:

- How can I decompose the system into a sequence of "features" that make sense.

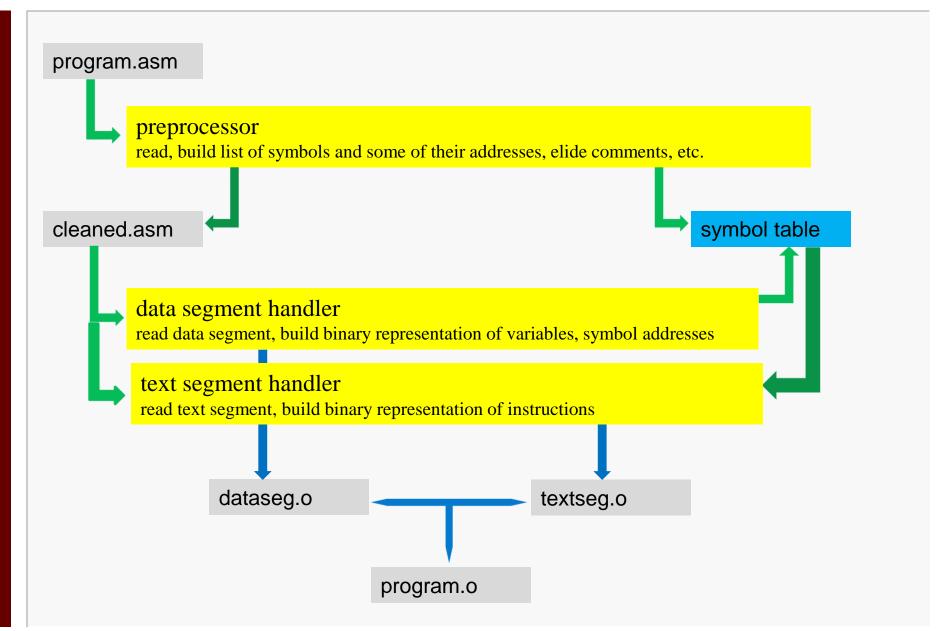
I often start by asking what's the minimum functionality I need to be able to actually do anything? Frequently, that's a matter of data acquisition... so I start with planning how to read my input file.

Now, what can I do next, once I know I can acquire data?

- In what order can I add those "features"?

This is usually not too difficult, provided I've given enough thought to the specification of the system and thought about how I might handle each process that needs to be carried out. But, if I get this wrong, I may have to perform a painful retrofit of something to my existing code.

One View



More Questions

But this analysis leads to more questions, including (in no particular order):

When/where do we deal with pseudo-instructions?

Some map to one basic instruction, some to two... is that an issue?

What "internal" objects and structures might be valuable in the design/implementation?

Instructions (assembly and machine)?

Build a list of instructions in memory at some point?

How should this be broken up into modules?

What focused, smaller parts might make up a text segment handler?

Testing

I will test features as I add them to the system.

NEVER implement a long list of changes and then begin testing. It's much harder to narrow down a logic error.

This may require creating some special test data, often partial versions of the full test data.

For example, I might hardwire a string holding a specific assembly instruction and pass it to my instruction parser/translator module.

Or, I might edit an assembly program so it only contains R-type instructions so I can focus on testing my handling of those.

Take advantage of the diagnostic tools:

gdb

can show you what is really happening at runtime which may not be what you believe is happening at runtime breakpoints, watchpoints, viewing values of variables

valgrind

can show you where you have memory-usage bugs finds memory leaks finds memory overruns where you exceed the bounds of a <u>dynamic</u> array

Pragmatics

Use the right development environment: CentOS 7 with gcc 4.8.x.

Do that from the beginning; don't wait until the last few days to "port" your code.

Read Maximizing Your Results in the project specification.

Small things within your control can make a huge difference in your final score.

There are many of you and few of us, so you cannot expect special treatment.

Use the supplied test harness (shell scripts and test files).

If your packaged submission doesn't run properly with these for you, it won't do that for us either.

There are many of you and few of us, so you will not receive special treatment.

Use the resources...

The course staff is here to help.

Most of us have prior experience with this assignment.