VirginiaTech

CS 4604: Introduction to Database Management Systems

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Lecture #5: Data Definition, Modification, and Constraints

Data Types in SQL

- Character strings:
  - CHAR(n): fixed-length string of n characters.
  - VARCHAR(n): string of length of up to n characters.
- Bit strings:
  - BIT(n): bit string of length n.
  - BIT VARYING(n): bit string of length upto n.
- BOOLEAN: possible values are TRUE, FALSE, and UNKNOWN (read Chapter 6.1.6 and 6.1.7).
- integers: INTEGER (INT), SHORTINT.
- floats: FLOAT (or REAL), DOUBLE PRECISION.
- fixed point numbers: DECIMAL(n, d): a number with n digits, with the decimal point d positions from the right.
- dates and times: DATE and TIME (read Chapter 6.1.5).

Creating and Deleting Tables

- A table is a relation that is physically stored in a database.
- A table is persistent; it exists indefinitely unless dropped or altered in some way.
- Creating a table:
  - CREATE TABLE Students (PID VARCHAR(8), Name CHAR(20), Address VARCHAR(255));
- Deleting a table:
  - DROP TABLE followed by the name of the table.

Modifying Table Schemas

- ALTER TABLE followed by the name of the relation followed by:
  - ADD followed by a column name and its data type.
  - Add date of birth (Dob) to Students:

    ALTER TABLE Students ADD Dob DATE;
  - DROP followed by a column name.

Null and Default Values

- SQL allows NULL for unknown attribute values. (Read Chapter 6.1.6, especially for how SQL treats comparisons using NULL).
- NULL not allowed in certain cases.
- We can specify a default value for an attribute using the DEFAULT keyword.
  - ALTER TABLE Students ADD Gender char(1) DEFAULT '?';

Inserting Data into a Table

- INSERT INTO R(A1,A2, . . . An) VALUES (v1, v2, . . . , vn).
  - (A1,A2, . . . ,An) can be a subset of R's schema.
  - Remaining attributes get NULL values.
  - Can omit names of attributes if we provide values for all attributes and list values in standard order.
- Insertion: Instead of VALUES, can use a SELECT statement.
  - Insert into the Professors table all professors who are mentioned in Teach but are not in Professors.

```
INSERT INTO Professors(PID)
SELECT ProfessorPID
FROM Teach
WHERE ProfessorPID NOT IN
      (SELECT PID FROM Professors);
```

Deleting Data from a Table

- DELETE FROM R WHERE C.
- Every tuple satisfying the condition C is deleted from R.

Updating Data in a Table

- An update in SQL is a change to one of the tuples existing in the database.
- Example: change the name of a student so that every male student has 'Mr.' added to the name and every female student has 'Ms.' added to the name.
  - UPDATE Students
    SET Name = 'Ms. ' || Name
    WHERE Gender = 'F';
  - UPDATE Students
    SET Name = 'Mr. ' || Name
    WHERE Gender = 'M';
- Can set multiple attributes in the SET clause, separated by commas.
- The WHERE clause can involve a subquery.

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Loading Data: BULK

- Different RDBMs have different syntax.
- PostgreSQL: Use the \copy ' filename' INTO TABLE tablename; at the psql prompt
- File format:
  - Tab-delimited with columns in the same order as the attributes.
  - Use \N to indicate null values.
- SQLite: see homework 1

- Do not make assumptions about how the RDBMS will behave!
- Check to make sure your data is not corrupted.
- Do not delete the original files that contain the raw data.

Saving Data

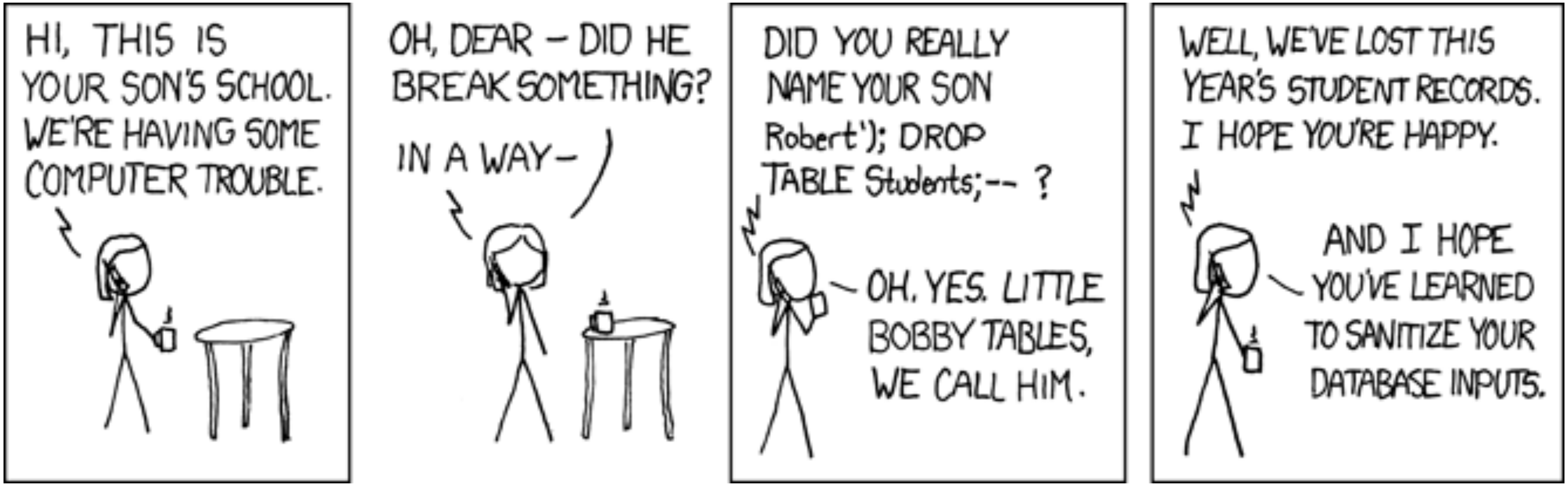
- Use the pg_dump program:
  - pg_dump -t table database

- Use man pg_dump for more information.

Constraints in Relational Algebra and SQL

Maintaining Integrity of Data

- Data is dirty.
- How does an application ensure that a database modification does not corrupt the tables?



Maintaining Integrity of Data

- Data is dirty.
- How does an application ensure that a database modification does not corrupt the tables?

- Two approaches:
  - Application programs check that database modifications are consistent.
  - Use the features provided by SQL.

Integrity Checking in SQL

- PRIMARY KEY and UNIQUE constraints.
- FOREIGN KEY constraints.
- Constraints on attributes and tuples.
- Triggers (schema-level constraints).

- How do we express these constraints?
- How do we check these constraints?
- What do we do when a constraint is violated?

Keys in SQL

- A set of attributes S is a key for a relation R if every pair of tuples in R disagree on at least one attribute in S.

- Select one key to be the PRIMARY KEY; declare other keys using UNIQUE.

Primary Keys in SQL

- Modify the schema of Students to declare PID to be the key.
  - CREATE TABLE Students(
    PID VARCHAR(8) PRIMARY KEY,
    Name CHAR(20), Address VARCHAR(255));

- What about Courses, which has two attributes in its key?
  - CREATE TABLE Courses(Number integer, DeptName:
    VARCHAR(8), CourseName VARCHAR(255), Classroom
    VARCHAR(30), Enrollment integer,
    PRIMARY KEY (Number, DeptName)
    );

Effect of Declaring PRIMARY KEYs

- Two tuples in a relation cannot agree on all the attributes in the key. DBMS will reject any action that inserts or updates a tuple in violation of this rule.
- A tuple cannot have a NULL value in a key attribute.

Other Keys in SQL

- If a relation has other keys, declare them using the UNIQUE keyword.
- Use UNIQUE in exactly the same places as PRIMARY KEY.

- There are two differences between PRIMARY KEY and UNIQUE:
  - A table may have only one PRIMARY KEY but more than one set of attributes declared UNIQUE.
  - A tuple may have NULL values in UNIQUE attributes.

Enforcing Key Constraints

- Upon which actions should an RDBMS enforce a key constraint?
- Only tuple update and insertion.
- RDMBS searches the tuples in the table to find if any tuple exists that agrees with the new tuple on all attributes in the primary key.
- To speed this process, an RDBMS automatically creates an efficient search index on the primary key.
- User can instruct the RDBMS to create an index on one or more attributes (If interested see Chapter 8.3).

Foreign Key Constraints

- Referential integrity constraint: in the relation Teach (that "connects" Courses and Professors), if Teach relates a course to a professor, then a tuple corresponding to the professor must exist in Professors.
- How do we express such constraints in Relational Algebra?
- Consider the Teach(ProfessorPID, Number, DeptName) relation.

  We want to require that every non-NULL value of ProfessorPID inTeach must be a valid ProfessorPID in Professors.
- RA $\pi_{\text{ProfessorPID}}(\text{Teach}) \subseteq \pi_{\text{PID}}(\text{Professors})$.

Foreign Key Constraints in SQL

- We want to require that every non-NULL value of ProfessorPID inTeach must be a valid ProfessorPID in Professors.
- In Teach, declare ProfessorPID to be a foreign key.
- CREATE TABLE Teach(ProfessorPID VARCHAR(8) REFERENCES Professor(PID), Name VARCHAR(30) ...);
- CREATE TABLE Teach(ProfessorPID VARCHAR(8), Name VARCHAR(30) ..., FOREIGN KEY ProfessorPID REFERENCES Professor(PID));
- If the foreign key has multiple attributes, use the second type of declaration.

Requirements for FOREIGN KEYs

- If a relation R declares that some of its attributes refer to foreign keys in another relation S, then these attributes must be declared UNIQUE or PRIMARY KEY in S.
- Values of the foreign key in R must appear in the referenced attributes of some tuple in S.

Enforcing Referential Integrity

- Three policies for maintaining referential integrity.
- Default policy: reject violating modifications.
- Cascade policy: mimic changes to the referenced attributes at the foreign key.
- Set-NULL policy: set appropriate attributes to NULL.

Default Policy for Enforcing Referential Integrity

- Reject violating modifications. There are four cases.

- Insert a new Teach tuple whose ProfessorPID is not NULL and is not the PID of any tuple in Professors.
- Update the ProfessorPID attribute in a tuple in Teach to a value that is not the PID value of any tuple in Professors.
- Delete a tuple in Professors whose PID value is the ProfessorPID value for one or more tuples in Teach.
- Update the PID value of a tuple in Professors when the old PID value is the value of ProfessorPID in one or more tuples of Teach.

Cascade Policy for Enforcing Referential Integrity

- Only applies to deletions of or updates to tuples in the referenced relation (e.g., Professors).

- If we delete a tuple in Professors, delete all tuples in Teach that refer to that tuple.
- If we update the PID value of a tuple in Professors from p1 to p2, update all value of ProfessorPID in Teach that are p1 to p2.

Set-NULL Policy for Enforcing Referential Integrity

- Also applies only to deletions of or updates to tuples in the referenced relation (e.g., Professors).

- If we delete a tuple in Professors, set the ProfessorPID attributes of all tuples in Teach that refer to the deleted tuple to NULL.
- If we update the PID value of a tuple in Professors from p1 to p2, set all values of ProfessorPID in Teach that are p1 to NULL

Specifying Referential Integrity Policies in SQL

- SQL allows the database designer to specify the policy for deletes and updates independently.
- Optionally follow the declaration of the foreign key with ON DELETE and/or ON UPDATE followed by the policy: SET NULL or CASCADE.
- Constraints can be circular, e.g., if there is a one-one mapping between two relations.
- In this case, SQL allows us to defer the checking of constraints. (Read Chapter 7.1.3).
- For your project, you do not have to consider deferring constraints.

Constraining Attributes and Tuples

- SQL also allows us to specify constraints on attributes in a relation and on tuples in a relation.
  - Disallow courses with a maximum enrollment greater than 100.
  - A chairperson of a department must teach at most one course every semester.
- How do we express such constraints in SQL?
- How can we change our minds about constraints?
- A simple constraint: NOT NULL
  - Declare an attribute to be NOT NULL after its type in a CREATE TABLE statement.
  - Effect is to disallow tuples in which this attribute is NULL.

Attribute-Based CHECK Constraints

- Disallow courses with a maximum enrollment greater than 100.
- This constraint only affects the value of a single attribute in each tuple.
- Follow the declaration of the Enrollment attribute with the CHECK keyword and a condition.
- CREATE TABLE Courses(...

  Enrollment INT CHECK (Enrollment <= 100) ...);
- The condition can be any condition that can appear in a WHERE clause.
- CHECK statement may use a subquery to mention other attributes of the same or other relations.
- An attribute-based CHECK constraint is checked only when the value of that attribute changes.

Tuple-Based CHECK Constraints

- Tuple-based CHECK constraints are checked whenever a tuple is inserted into or updated in a relation.
- Designer may add these constraints after the list of attributes in a CREATE TABLE statement.
- A chairperson of a department teaches at most one course in any semester.

```
CREATE TABLE Teach(...
    CHECK ProfessorPID NOT IN
        ((SELECT ProfessorPID FROM Teach)
        INTERSECT
        (SELECT ChairmanPID FROM Departments)
        )
        );
```

Modifying Constraints

- SQL allows constraints to be named.
- Use CONSTRAINT followed by the name of the constraint in front of PRIMARY KEY, UNIQUE, or CHECK.
- Can use constraint names in ALTER TABLE statements to delete constraints: say DROP CONSTRAINT followed by the name of the constraint.
- Can add constraints in an ALTER TABLE statement using ADD CONSTRAINT followed by an optional name followed by the (required) CHECK statement.

Assertions

- These are database-schema elements, like relations
- Defined by:
  - CREATE ASSERTION <name>

    CHECK (<condition>);
- Condition may refer to any relation or attribute in the database schema.

Assertions: Example

- Can’t have more courses than students (‘Pigeonhole Principle’)

```
CREATE ASSERTION FewStudents CHECK (
  (SELECT COUNT(*) FROM Students) <=
  (SELECT COUNT(*) FROM Courses)
);
```

Triggers: Motivation

- Assertions are powerful, but the DBMS often can't tell when they need to be checked.
- Attribute- and tuple-based checks are checked at known times→but not powerful.
- Triggers let the user decide when to check for any condition.

Triggers

- Trigger: procedure that starts automatically if specified changes occur to the DBMS
- A trigger has three parts:
  - Event (activates the trigger)
  - Condition (tests whether the triggers should run)
  - Action (what happens if the trigger runs)

Maintain a unary relation New_Courses which has the list of brand new courses

Triggers

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  - Event (activates the trigger)
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```
CREATE TRIGGER incr_count
AFTER INSERT ON Teach            // Event
REFERENCING NEW ROW AS new
FOR EACH ROW
WHEN (new.id NOT IN (SELECT ID FROM Courses))    // Condition
INSERT INTO New_Courses(id) VALUES(new.id);      // Action
```

OK, what could have been done?