CS 4604: Introduction to Database Management Systems

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Lecture #4: Subqueries in SQL
Announcements

- Project assignment 1 due today

- Homework 1 released today
  - due next Friday 2/8
  - SQL and Relational Algebra
Most popular embedded db in the world
– Iphone (iOS), Android, Chrome....

(Very) Easy to use: no need to set it up

Self-contained: data+schema

DB on your laptop: useful for testing, understanding....
Linear Notation for Relational Algebra

- Relational algebra expressions can become very long.

- Use linear notation to store results of intermediate expressions.
  - A relation name and a parenthesised list of attributes for that relation
  - Use Answer as the conventional name for the final result
  - The assignment symbol :=
  - Any expression in relational algebra on the right
Example of Linear Notation

- Name pairs of students who live at the same address.
- Normal expression:

\[
\pi_{S1.Name, S2.Name} \\
\sigma_{S1.Address = S2.Address} \\
(\rho_{S1}(Students) \times \rho_{S2}(Students))
\]
Example of Linear Notation

- **Normal expression:**

\[
\pi_{S1.Name,S2.Name} (\sigma_{S1.Address=S2.Address} (\rho_{S1}(Students) \times \rho_{S2}(Students)))
\]

- **Linear Notation:**

\[
\text{Pairs}(P1, N1, A1, P2, N2, A2) := \rho_{S1}(Students) \times \rho_{S2}(Students)
\]
\[
\text{Matched}(P1, N1, A1, P2, N2, A2) := \sigma_{A1=A2}(\text{Pairs}(P1, N1, A1, P2, N2, A2))
\]
\[
\text{Answer}(\text{Name1}, \text{Name2}) := \pi_{N1,N2}(\text{Matched}(P1, N1, A1, P2, N2, A2))
\]
Interpreting Queries Involving Multiple Relations

- SELECT A, B FROM R, S WHERE C;
- Nested loops:
  for each tuple t1 in R
    for each tuple t2 in S
      if the attributes in t1 and t2 satisfy C
        output the tuples involving attributes A and B
Interpreting Queries Involving Multiple Relations

- SELECT A, B FROM R, S WHERE C;
- Conversion to relational algebra:
  \[ \pi_{A,B}(\sigma_{C}(R \times S)) \]

  Compute R X S
  Apply selection operator \( \sigma() \) to R X S
  Project the result tuples to attributes A and B
Motivation for Subqueries

- Find the name of the professor who teaches “CS 4604.”

```sql
SELECT Name
FROM Professors, Teach
WHERE (PID = ProfessorPID) AND (Number = ‘4604’) AND (DeptName = ‘CS’);
```

- Do we need to take the natural join of two big relations just to get a relation with one tuple?
- Can we rewrite the query without using a join?
Nesting

- A query can be put inside another query
- Most commonly in the WHERE clause
- Sometimes in the FROM clause (depending on the software)
- This subquery is executed first (if possible)
Subquery Example

- Find the name of the professor who teaches “CS 4604.”

SELECT Name
FROM Professors
WHERE PID =
  (SELECT ProfessorPID
   FROM Teach
   WHERE (Number = 4604) AND (DeptName = 'CS'))
;

- When using =, the subquery must return a single tuple
Conditions Involving Relations

- SQL includes a number of operators that apply to a relation and produce a boolean result.
- These operators are very useful to apply on results of sub-queries.
Conditions Involving Relations

- Let R be a relation and t be a tuple with the same set of attributes.
  - \textbf{EXISTS} R is true if and only if R contains at least one tuple.
  - t \textbf{IN} R is true if and only if t equals a tuple in R.
  - t > \textbf{ALL} R is true if and only if R is unary (has one attribute) and t is greater than every value in R.
    - Can use any of the other five comparison operators.
    - If we use $\langle\rangle$, R need not be unary.
  - t > \textbf{ANY} R (which is unary) is true if and only if t is greater than at least one value in R.

- We can use \textbf{NOT} to negate \textbf{EXISTS}, \textbf{ALL}, and \textbf{ANY}.
Subqueries Using Conditions

- Find the departments of the courses taken by the student with name ‘Suri’.

```sql
SELECT DeptName
FROM Take
WHERE StudentPID IN
  ( SELECT PID
    FROM Students
    WHERE (Name = ‘Suri’) )
```

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Correlated vs Uncorrelated

- The previous subqueries did not depend on anything outside the subquery
  - ...and thus need to be executed just once.
  - These are called **uncorrelated**.

- A **correlated** subquery depends on data from the outer query
  - ... and thus has to be executed for each row of the outer table(s)
Correlated Subqueries

- Find course names that have been used for two or more courses.

```
SELECT CourseName
FROM Courses AS First
WHERE CourseName IN
  (SELECT CourseName
   FROM Courses
   WHERE (Number <> First.Number)
   AND (DeptName <> First.DeptName))
```

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Evaluating Correlated Subqueries

SELECT CourseName
FROM Courses AS First
WHERE CourseName IN
  (SELECT CourseName
   FROM Courses
   WHERE (Number <> First.Number)
   AND (DeptName <> First.DeptName)
  );

- Evaluate query by looping over tuples of First, and for each tuple evaluate the subquery.
- Scoping rules: an attribute in a subquery belongs to one of the tuple variables in that subquery’s FROM clause, or to the immediately surrounding subquery, and so on.
Subqueries in FROM clauses

- Can use a subquery as a relation in a FROM clause.
- We must give such a relation an alias using the AS keyword.
- Let us find different ways of writing the query “Find the names of Professors who have taught the student whose first name is ‘Suri’.”
- The old way:

  ```sql
  SELECT Professors.Name
  FROM Professors, Take, Teach, Students
  WHERE (Professors.PID = Teach.ProfessorPID)
    AND (Teach.CourseNumber = Take.CourseNumber)
    AND (Teach.DeptName = Take.DeptName)
    AND (Take.StudentPID = Student.PID)
    AND (Student.Name = 'Suri %');
  ```
“Find the names of (Professors who have taught (courses taken by (student with first name ‘Suri’))).”

```
SELECT Name
FROM Professors
WHERE PID IN
    (SELECT ProfessorPID
        FROM Teach
        WHERE (Number, DeptName) IN
            (SELECT Number, DeptName
                FROM Take, Students
                WHERE (StudentPID = PID) AND
                    (Students.Name = 'Suri%')));
```
Unrolling it further

- SELECT Name
  FROM Professors
  WHERE PID IN
    (SELECT ProfessorPID
     FROM Teach
     WHERE (Number, DeptName) IN
       (SELECT Number, DeptName
        FROM Take
        WHERE StudentPID IN
          (SELECT PID
           FROM Students
           WHERE Name = 'Suri %')));
Aggregate Operators

- **COUNT (*)**
- **COUNT ([DISTINCT] A)**
  - A is a column
- **SUM ([DISTINCT] A)**
- **AVG ([DISTINCT] A)**
- **MAX (A)**
- **MIN (A)**
- Count the number of sailors

  ```sql
  SELECT COUNT (*) FROM Sailors S
  ```
Find name and age of the oldest sailor(s)

\[
\text{SELECT } S.\text{sname}, \text{MAX} (S.\text{age}) \\
\text{FROM Sailors } S
\]

- This is illegal, but why?
  - Cannot combine a column with a value

\[
\text{SELECT } S.\text{sname}, S.\text{age} \\
\text{FROM Sailors } S \\
\text{WHERE } S.\text{age} = (\text{SELECT } \text{MAX} (S2.\text{age}) \text{ FROM Sailors } S2)
\]
So far, aggregate operators are applied to all (qualifying) tuples.
  – Can we apply them to each of several groups of tuples?

Example: find the age of the youngest sailor for each rating level.
  – In general, we don’t know how many rating levels exist, and what the rating values for these levels are!
  – Suppose we know that rating values go from 1 to 10; we can write 10 queries that look like this:

```
For i = 1, 2, ..., 10:

SELECT MIN (S.age)
FROM Sailors S
WHERE S.rating = i
```
Find the age of the youngest sailor for each rating level

```
SELECT S.rating, MIN (S.age) as age
FROM Sailors S
GROUP BY S.rating
```

(1) The sailors tuples are put into “same rating” groups.

(2) Compute the Minimum age for each rating group.
Find the age of the youngest sailor for each rating level that has at least 2 members

\[
\text{SELECT } S.\text{rating}, \text{MIN}(S.\text{age}) \text{ as minage} \\
\text{FROM Sailors } S \\
\text{GROUP BY } S.\text{rating} \\
\text{HAVING COUNT(*) > 1}
\]

1. The sailors tuples are put into “same rating” groups.

2. Eliminate groups that have < 2 members.

3. Compute the Minimum age for each rating group.
Queries With \textit{GROUP BY} and \textit{HAVING}

\begin{verbatim}
SELECT [DISTINCT] target-list
FROM relation-list
WHERE qualification
GROUP BY grouping-list
HAVING group-qualification
\end{verbatim}

- The \textit{target-list} contains (i) attribute names (ii) terms with aggregate operations (e.g., \texttt{AVG (S.age)}).
- The attribute list (e.g., \texttt{S.rating}) in \textit{target-list} must be in \textit{grouping-list}.
- The attributes in group-qualification must be in \textit{grouping-list}.