

Data Definition and Modification in SQL

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Data Types in SQL

- Character strings:
 - CHAR(n): fixed-length string of n characters.
 - VARCHAR(n): string of length of up to n characters.
- Bit strings:
 - BIT(n): bit string of length n.
 - BIT VARYING(n): bit string of length upto n.
- BOOLEAN: possible values are TRUE, FALSE, and UNKNOWN (read Chapter 6.1.7).
- integers: INTEGER (INT), SHORTINT.
- floats: FLOAT (or REAL), DOUBLE PRECISION.
- fixed point numbers: DECIMAL(n, d): a number with n digits, with the decimal point d positions from the right.
- dates and times: DATE and TIME (read Chapter 6.1.5).

Creating and Deleting Tables

- A table is a relation that is physically stored in a database.
- A table is **persistent**; it exists indefinitely unless dropped or altered in some way.
- Creating a table: CREATE TABLE followed by the name of the relation and a parenthesized list of attribute names and their types.
- **CREATE TABLE** Students (PID VARCHAR(8), Name CHAR(20), Address VARCHAR(255));
- Deleting a table: **DROP TABLE** followed by the name of the table.

Modifying Table Schemas

- **ALTER TABLE** followed by the name of the relation followed by:
 - **ADD** followed by a column name and its data type.
 - Add date of birth (Dob) to Students:
 - ALTER TABLE Students ADD Dob DATE;
- **DROP** followed by a column name.

Null and Default Values

- SQL allows NULL for unknown attribute values. (Read Chapter 6.1.6, especially for how SQL treats comparisons using NULL).
- NULL not allowed in certain cases.
- We can specify a default value for an attribute using the DEFAULT keyword.
 - ALTER TABLE Students ADD Gender char(1) DEFAULT '?';

Inserting Data into a Table

- **INSERT INTO** R(A1,A2, . . . An) VALUES (v1, v2, . . . , vn).
 - (A1,A2, . . . ,An) can be a subset of R's schema.
 - Remaining attributes get NULL values.
 - Can omit names of attributes if we provide values for all attributes and list values in standard order.
- Insertion: Instead of VALUES, can use a SELECT statement.
 - Insert into the Professors table all professors who are mentioned in Teach but are not in Professors.

```
INSERT INTO Professors(PID)
    SELECT ProfessorPID
    FROM Teach
    WHERE ProfessorPID NOT IN
        (SELECT PID FROM Professors);
```

Deleting Data from a Table

- **DELETE FROM** R WHERE C.
- Every tuple satisfying the condition C is deleted from R.

Updating Data in a Table

- An update in SQL is a change to one of the tuples existing in the database.
- Example: change the name of a student so that every male student has 'Mr.' added to the name and every female student has 'Ms.' added to the name.

```
SET
```

 - **UPDATE** Students

```
    SET Name = 'Ms. ' || Name  
    WHERE Gender = 'F';
```
 - **UPDATE** Students

```
    SET Name = 'Mr. ' || Name  
    WHERE Gender = 'M';
```
- Can set multiple attributes in the SET clause, separated by commas.
- The WHERE clause can involve a subquery.

Loading Data: BULK

- Different RDBMs have different syntax.
- PostgreSQL: Use the `\copy 'filename' INTO TABLE tablename;` at the psql prompt
- File format:
 - Tab-delimited with columns in the same order as the attributes.
 - Use `\N` to indicate null values.
- Do not make assumptions about how the RDBMS will behave!
- Check to make sure your data is not corrupted.
- Do not delete the original files that contain the raw data.

Saving Data

- Use the `pg_dump` program:
 - `pg_dump -t table database`
- Use `man pg_dump` for more information.

Specific Project Guidelines

- We will create an account and a database for each student.
- A database for each project will be created.
 - The name of the database is the name of your project.
 - Only the members of each project will be able to access the database for their project.
- A webpage detailing how you can access the database is maintained.
- You can create as many tables within a database as you want.

General Project Guidelines

- The database schema is not something that should change often.
 - Think long and hard about your schema.
 - DROP may be better than ALTER TABLE.
- Do not delete the files containing raw data.
- Read documentation for the RDBMS you are using.

Constraints in Relational Algebra and SQL

Maintaining Integrity of Data

- Data is **dirty**.
- How does an application ensure that a database modification does not corrupt the tables?
- Two approaches:
 - Application programs check that database modifications are consistent.
 - Use the features provided by SQL.

Integrity Checking in SQL

- PRIMARY KEY and UNIQUE constraints.
- FOREIGN KEY constraints.
- Constraints on attributes and tuples.
- Triggers (schema-level constraints).

- How do we express these constraints?
- How do we check these constraints?
- What do we do when a constraint is violated?

Keys in SQL

- A set of attributes S is a key for a relation R if every pair of tuples in R **disagree on at least one attribute** in S .
- Select one key to be the **PRIMARY KEY**; declare other keys using **UNIQUE**.

Primary Keys in SQL

- Modify the schema of Students to declare PID to be the key.
 - CREATE TABLE Students(
 PID VARCHAR(8) PRIMARY KEY,
 Name CHAR(20), Address VARCHAR(255));
- What about Courses, which has two attributes in its key?
 - CREATE TABLE Courses(Number integer, DeptName:
 VARCHAR(8), CourseName VARCHAR(255), Classroom
 VARCHAR(30), Enrollment integer,
 PRIMARY KEY (Number, DeptName)
);

Effect of Declaring PRIMARY KEYS

- Two tuples in a relation cannot agree on all the attributes in the key. DBMS will reject any action that inserts or updates a tuple in violation of this rule.
- A tuple cannot have a NULL value in a key attribute.

Other Keys in SQL

- If a relation has other keys, declare them using the UNIQUE keyword.
- Use UNIQUE in exactly the same places as PRIMARY KEY.
- There are two differences between PRIMARY KEY and UNIQUE:
 - A table may have **only one PRIMARY KEY** but more than one set of attributes declared UNIQUE.
 - A tuple **may have NULL values in UNIQUE** attributes.

Enforcing Key Constraints

- Upon which actions should an RDBMS enforce a key constraint?
- Only tuple update and insertion.
- RDBMS searches the tuples in the table to find if any tuple exists that agrees with the new tuple on all attributes in the primary key.
- To speed this process, an RDBMS automatically creates an efficient search index on the primary key.
- User can instruct the RDBMS to create an index on one or more attributes (If interested see Chapter 8.3).

Foreign Key Constraints

- **Referential integrity constraint**: in the relation Teach (that “connects” Courses and Professors), if Teach relates a course to a professor, then a tuple corresponding to the professor **must** exist in Professors.
- How do we express such constraints in Relational Algebra?
- Consider the Teach(ProfessorPID, Number, DeptName) relation.

We want to require that every non-NULL value of ProfessorPID in Teach must be a valid ProfessorPID in Professors.

- **RA** $\pi_{\text{ProfessorPID}}(\text{Teach}) \subseteq \pi_{\text{PID}}(\text{Professors})$.

Foreign Key Constraints in SQL

- We want to require that every non-NULL value of ProfessorPID in Teach must be a valid ProfessorPID in Professors.
- In Teach, declare ProfessorPID to be a foreign key.
- CREATE TABLE Teach(ProfessorPID VARCHAR(8) REFERENCES Professor(PID), Name VARCHAR(30) ...);
- CREATE TABLE Teach(ProfessorPID VARCHAR(8), Name VARCHAR(30) ..., FOREIGN KEY ProfessorPID REFERENCES Professor(PID));
- If the foreign key has multiple attributes, use the second type of declaration.

Requirements for FOREIGN KEYS

- If a relation R declares that some of its attributes refer to foreign keys in another relation S, then these attributes **must** be declared UNIQUE or PRIMARY KEY in S.
- Values of the foreign key in R must appear in the referenced attributes of some tuple in S.

Enforcing Referential Integrity

- **Three** policies for maintaining referential integrity.
- Default policy: reject violating modifications.
- Cascade policy: mimic changes to the referenced attributes at the foreign key.
- Set-NULL policy: set appropriate attributes to NULL.

Default Policy for Enforcing Referential Integrity

- **Reject** violating modifications. There are **four cases**.
- Insert a new Teach tuple whose ProfessorPID is not NULL and is not the PID of any tuple in Professors.
- Update the ProfessorPID attribute in a tuple in Teach to a value that is not the PID value of any tuple in Professors.
- Delete a tuple in Professors whose PID value is the ProfessorPID value for one or more tuples in Teach.
- Update the PID value of a tuple in Professors when the old PID value is the value of ProfessorPID in one or more tuples of Teach.

Cascade Policy for Enforcing Referential Integrity

- Only applies to deletions of or updates to tuples in the referenced relation (e.g., Professors).
- If we delete a tuple in Professors, delete all tuples in Teach that refer to that tuple.
- If we update the PID value of a tuple in Professors from p1 to p2, update all value of ProfessorPID in Teach that are p1 to p2.

Set-NULL Policy for Enforcing Referential Integrity

- Also applies only to deletions of or updates to tuples in the referenced relation (e.g., Professors).
- If we delete a tuple in Professors, set the ProfessorPID attributes of all tuples in Teach that refer to the deleted tuple to NULL.
- If we update the PID value of a tuple in Professors from p1 to p2, set all values of ProfessorPID in Teach that are p1 to NULL

Specifying Referential Integrity Policies in SQL

- SQL allows the database designer to specify the policy for deletes and updates independently.
- Optionally follow the declaration of the foreign key with **ON DELETE** and/or **ON UPDATE** followed by the policy: **SET NULL** or **CASCADE**.
- Constraints can be circular, e.g., if there is a one-one mapping between two relations.
- In this case, SQL allows us to defer the checking of constraints. (Read Chapter 7.1.3).
- For your project, you do not have to consider deferring constraints.

Constraining Attributes and Tuples

- SQL also allows us to specify constraints on attributes in a relation and on tuples in a relation.
 - Disallow courses with a maximum enrollment greater than 100.
 - A chairperson of a department must teach at most one course every semester.
- How do we express such constraints in SQL?
- How can we change our minds about constraints?
- A simple constraint: NOT NULL
 - Declare an attribute to be NOT NULL after its type in a CREATE TABLE statement.
 - Effect is to disallow tuples in which this attribute is NULL.

Attribute-Based CHECK Constraints

- Disallow courses with a maximum enrollment greater than 100.
- This constraint only affects the value of a single attribute in each tuple.
- Follow the declaration of the Enrollment attribute with the CHECK keyword and a condition.
- CREATE TABLE Courses(...
 Enrollment INT CHECK (Enrollment <= 100) ...);
- The condition can be any condition that can appear in a WHERE clause.
- CHECK statement may use a subquery to mention other attributes of the same or other relations.
- An attribute-based CHECK constraint is checked **only when the value of that attribute changes.**

Tuple-Based CHECK Constraints

- Tuple-based CHECK constraints are checked whenever a tuple is inserted into or updated in a relation.
- Designer may add these constraints after the list of attributes in a CREATE TABLE statement.
- A chairperson of a department teach at most one course in any semester.

```
CREATE TABLE Teach(...  
    CHECK ProfessorPID NOT IN  
        ((SELECT ProfessorPID FROM Teach)  
         INTERSECT  
         (SELECT ChairmanPID FROM  
          Departments)  
        )  
    );
```

Modifying Constraints

- SQL allows constraints to be named.
- Use CONSTRAINT followed by the name of the constraint in front of PRIMARY KEY, UNIQUE, or CHECK.
- Can use constraint names in ALTER TABLE statements to delete constraints: say DROP CONSTRAINT followed by the name of the constraint.
- Can add constraints in an ALTER TABLE statement using ADD CONSTRAINT followed by an optional name followed by the (required) CHECK statement.

Triggers

- Trigger: procedure that starts automatically if specified changes occur to the DBMS
- A trigger has three parts:
 - **Event** (activates the trigger)
 - **Condition** (tests whether the triggers should run)
 - **Action** (what happens if the trigger runs)

```
CREATE TRIGGER incr_count AFTER INSERT ON Students // Event
WHEN (new.age < 18) // Condition
FOR EACH ROW
BEGIN // ACTION
    count := count + 1
END
```