Chapter 3



OS Organization

Design of OS

- Factors influencing *design* of OS
 - 1. Performance
 - 2. Protection/Security
 - 3. Correctness
 - 4. Maintainability
 - 5. Commercial factors
 - 6. Standard & Open Systems

(1) Performance

- Functionality v/s Performance
 - More resource abstraction
 - Higher levels of resource abstraction
- Coding OS w.r.t. Performance
 - Assembly => Fast execution
 - BUT Assembly => Debugging ???
- Others?



- OS MUST NOT allow one process to interfere with the operations of another process
 - File access
 - Memory space
 - Resources
- Therefore, need to implement strategies that support Isolation & Sharing
- Challenge is:
 - If OS implements a policy, how to prevent <u>application</u> from changing it



- Maintainability
 - Design and write systems to be maintainable
 - => Sacrifice performance
- Correctness
 - Does the OS meet the requirements?
 - Can we write valid set of requirements?

(5) Commercial influence

Commercial Influence

- ∠ DOS => IBM-PC
- UNIX => open platform
- Commercial influence
 - => machine nuances that hinder portability
 - Z UNIX => portable
 - MAC ???
 - Windows ???



- Early systems: User tied to ONE vendor
- Desire: User gets pieces from ANY set of vendors
 - => Need for Standards and Open Systems
- Open Systems
 - => Network of heterogeneous systems
 - =>Information flow [Big Endian v/s Little Endian]



- Open systems achieved through
 - Application integration => common interface
 - Portability => more applications among hardware platforms
 - Interoperability
 - Standardize remote access facilities
 - => All systems talk same language over the network
- ∠ POSIX = Open system
 - Standardize OS interfaces

Basic Functions of OS

- Device Management
- 2. Process / Resource Management
- 3. Memory Management
- 4. File Management

Device Management

- Allocation
- Need device drivers
 - Must be able to configure into OS without recompiling OS (no Source Code)

Process / Resource Management

- Process
 - Creating
 - Destroying

 - **Running**
- Resource
 - Isolation
 - Sharing

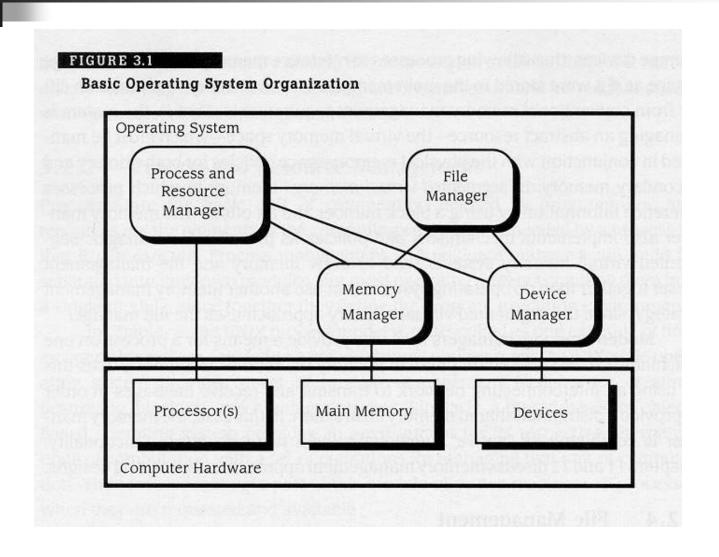
Memory Management

- Allocation & use of main memory
 - Isolation & Protection
 - Sharing
- Virtual Memory
 - Main memory & storage devices
 - Reference 'memory' on storage devices
- Segmented VM viable approach
 - Block & Offset

File Management

- Transfer from main memory to file
 - ∠ Code (VM)
 - ∠ Data (VM)
 - Editors
- Different file management strategies
 - Sequential
 - Indexed
 - Direct access
 - Networked

Basic OS Organization



Implementation Considerations

Process Modes

Kernels

Method of requesting system services

Processor Modes

- Supervisor mode
 - Z Can execute any instruction
- User mode
 - Subset of instructions

In UNIX:

What can root execute that application cannot?

z re-nice : OS call
...

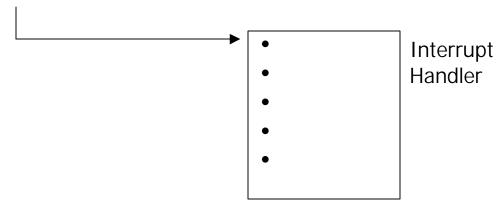
chown : OS call

∠ IOCTL (OS call) — if user interleaves output on printer

Memory accesses

Kernel

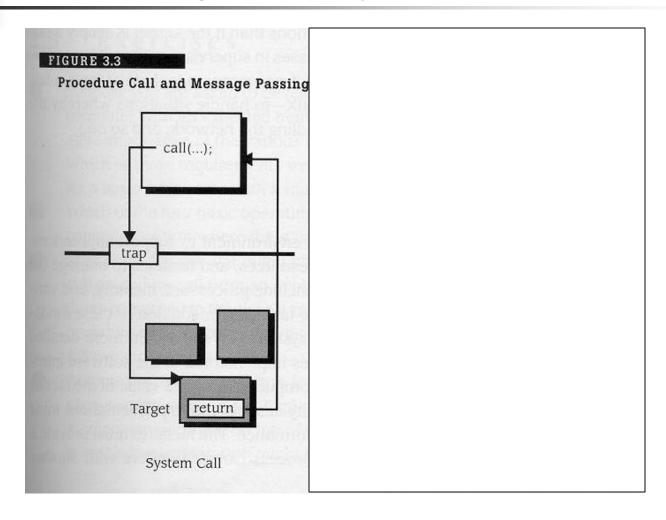
- Trusted part of the OS
- Executes in Supervisor mode
- Generally, memory resident
- OS <u>extension</u> run in User mode
 - Example: Drivers
- Kernel functions are invoked by "trap"



Requesting Service from OS

- System call
 - Process traps to OS Interrupt Handler
 - Supervisor mode set
 - Desired function executed
 - User mode set
 - Returns to application

Requesting Svc: System Call

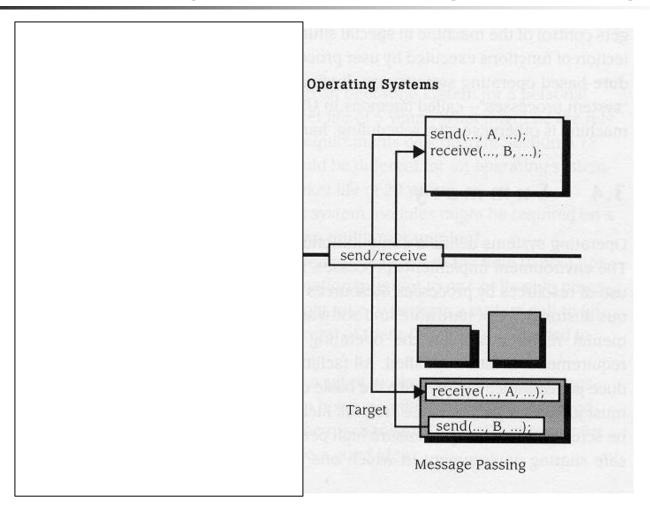


Message Passing

- User process constructs message indicating function (service) needed
- Invokes send to pass message to OS
- Process blocks
- OS receives message
- OS initiates Function execution
- Upon Function completion, OS Returns ("OK")
- Process un-blocks

Send and Receive analyze message for proper format, etc.

Requesting Svc: Message Passing



Message Passing...

System call are more efficient

BUT

they also unduly tie the Application to specifics of the OS