A graph $G$ consists of a set $V$ of vertices and a set $E$ of pairs of distinct vertices from $V$. These pairs of vertices are called edges.

If the pairs of vertices are unordered, $G$ is an undirected graph. If the pairs of vertices are ordered, $G$ is a directed graph or digraph.
An undirected graph $G$, where:

\[ V = \{a, b, c, d, e, f, g, h, i\} \]

\[ E = \{\{a, b\}, \{a, c\}, \{b, e\}, \{b, h\}, \{b, i\}, \{c, d\}, \{c, e\}, \{e, f\}, \{e, g\}, \{h, i\}\} \]

$e = \{c, d\}$ is an edge, incident upon the vertices $c$ and $d$

Two vertices, $x$ and $y$, are adjacent if $\{x, y\}$ is an edge (in $E$).

A path in $G$ is a sequence of distinct vertices, each adjacent to the next.

A path is simple if no vertex occurs twice in the path.

A cycle in $G$ is a path in $G$, containing at least three vertices, such that the last vertex in the sequence is adjacent to the first vertex in the sequence.
A graph $G$ is **connected** if, given any two vertices $x$ and $y$ in $G$, there is a path in $G$ with first vertex $x$ and last vertex $y$.

The graph on the previous slide is connected.

If a graph $G$ is not connected, then we say that a maximal connected set of vertices is a **component** of $G$. 
Directed Graph Terminology

The terminology for directed graphs is only slightly different…

\[ e = (c, d) \] is an edge, from \( c \) to \( d \)

A directed path in a directed graph \( G \) is a sequence of distinct vertices, such that there is an edge from each vertex in the sequence to the next.

A directed graph \( G \) is weakly connected if, the undirected graph obtained by suppressing the directions on the edges of \( G \) is connected (according to the previous definition).

A directed graph \( G \) is strongly connected if, given any two vertices \( x \) and \( y \) in \( G \), there is a directed path in \( G \) from \( x \) to \( y \).
Adjacency Matrix Representation

A graph may be represented by a two-dimensional adjacency matrix:

If $G$ has $n = |V|$ vertices, let $M$ be an $n$ by $n$ matrix whose entries are defined by

$$m_{ij} = \begin{cases} 
1 & \text{if } (i, j) \text{ is an edge} \\
0 & \text{otherwise}
\end{cases}$$

$$M(G) = \begin{bmatrix}
0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \\
1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0
\end{bmatrix}$$
The adjacency matrix:

- \( \Theta(1) \) to determine existence of a specific edge
- \( \Theta(|V|^2) \) storage cost (cut cost by 75% or more by changing types)
- \( \Theta(|V|) \) for finding all vertices accessible from a specific vertex
- \( \Theta(1) \) to add or delete an edge
- Not easy to add or delete a vertex; better for static graph structure.
- Symmetric matrix for undirected graph; so half is redundant then.

\[
M(G) = \begin{bmatrix}
0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \\
1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0
\end{bmatrix}
\]
Adjacency Table Representation

A slightly different approach is to represent only the adjacent nodes in the structure:

```
0 |  1  2
1 |  4  7  8
2 |  0  3  4
3 |
4 |  1  6
5 |  4
6 |
7 |  1  8
8 |
```
The adjacency list structure is simply a linked version of the adjacency table:

Array of linked lists, where list nodes store node labels for neighbors.
Adjacency List Representation

The adjacency list structure:

- Worst case: $\Theta(|V|)$ to determine existence of a specific edge
- $\Theta(|V| + |E|)$ storage cost
- Worst case: $\Theta(|V|)$ for finding all neighbors of a specific vertex
- Worst case: $\Theta(|V|)$ to add or delete an edge
- Still not easy to add or delete a vertex; however, we can use a linked list in place of the array.

Note, for an undirected graph, the upper bound on the number of edges is:

$$|E| \leq |V|*(|V|-1)$$

So, the space comparison with the adjacency matrix scheme is not trivial.
An Adjacency Matrix Class

```java
public class AdjMatrix {

    private int numVertices;
    private boolean[] Marker;  // used for vertex marking
    private int[][] Edge;      // Edge[i][j] == 1 iff (i,j) exists

    public AdjMatrix(int numV) {...

    public boolean addEdge(int Src, int Trm) {...
    public boolean delEdge(int Src, int Trm) {...
    public boolean hasEdge(int Src, int Trm) {...

    public int firstNeighbor(int Src) {...
    public int nextNeighbor(int Src, int Prev) {...

    public boolean isMarked(int V) {...
    public boolean Mark(int V) {...
    public boolean unMark(int V) {...

    public void Clear() {...
        // erase edges and vertex marks
    public void Display() {...
```