

Design: a mod-8 Counter

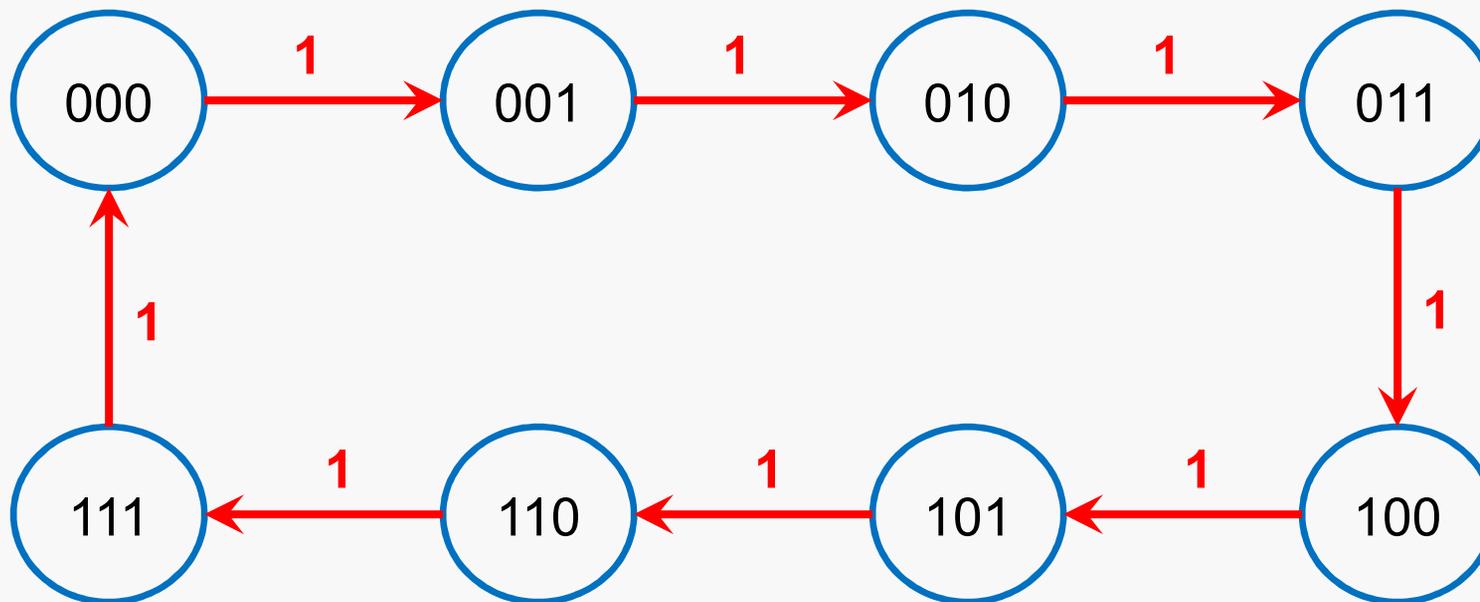
A mod-8 counter stores an integer value, and increments that value (say) on each clock tick, and wraps around to 0 if the previous stored value was 7.

So, the stored value follows a cycle:

000 001 010 011 100 101 110 111

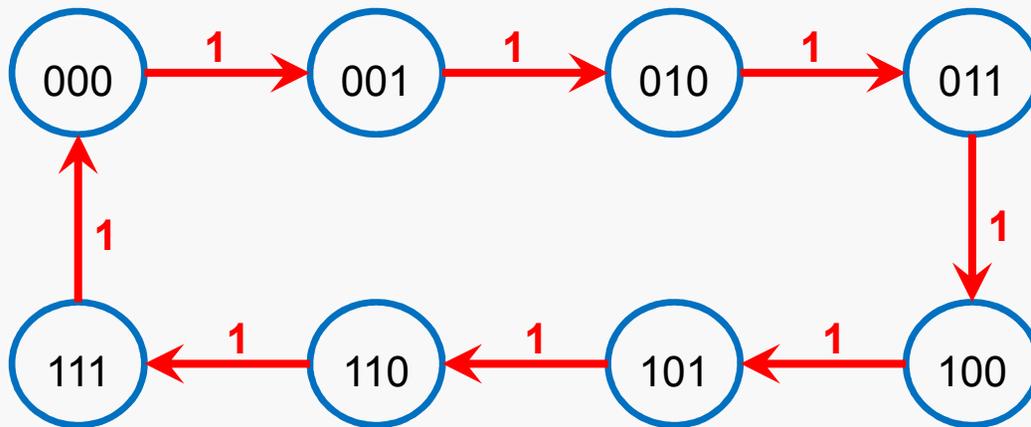
We need eight different states for our counter, one for each value from 0 to 7.

We can describe the operation by drawing a *state machine*. The nodes represent states and the edges represent transitions and are labeled with the input (clock in this case) that causes the transition to occur.



A state table summarizes the state machine and is useful in deriving equations later:

Current State			Input	Next State		
---	---	---	---	---	---	---
0	0	0	1	0	0	1
0	0	1	1	0	1	0
0	1	0	1	0	1	1
0	1	1	1	1	0	0
1	0	0	1	1	0	1
1	0	1	1	1	1	0
1	1	0	1	1	1	1
1	1	1	1	0	0	0



We will derive an equation for each of the next state functions:

Current State			Input	Next State		
C2	C1	C0		N2	N1	N0
0	0	0	1	0	0	1
0	0	1	1	0	1	0
0	1	0	1	0	1	1
0	1	1	1	1	0	0
1	0	0	1	1	0	1
1	0	1	1	1	1	0
1	1	0	1	1	1	1
1	1	1	1	0	0	0

$$N0 = \overline{C2} \cdot \overline{C1} \cdot \overline{C0} + \overline{C2} \cdot C1 \cdot \overline{C0} + C2 \cdot \overline{C1} \cdot \overline{C0} + C2 \cdot C1 \cdot \overline{C0} = \overline{C0}$$

$$N1 = \overline{C2} \cdot \overline{C1} \cdot C0 + \overline{C2} \cdot C1 \cdot \overline{C0} + C2 \cdot \overline{C1} \cdot C0 + C2 \cdot C1 \cdot \overline{C0} = \overline{C1} \cdot C0 + C1 \cdot \overline{C0}$$

$$\begin{aligned} N2 &= \overline{C2} \cdot C1 \cdot C0 + C2 \cdot \overline{C1} \cdot \overline{C0} + C2 \cdot \overline{C1} \cdot C0 + C2 \cdot C1 \cdot \overline{C0} \\ &= \overline{C2} \cdot C1 \cdot C0 + C2 \cdot \overline{C1} + C2 \cdot \overline{C0} \end{aligned}$$

Since each state is represented by a 3-bit integer, we can represent the states by using a collection of three flip-flops (more-or-less a mini-register).

We will implement the circuit using D flip-flops, which make for a simple translation from the state table because a D flip-flop simply accepts its input value.

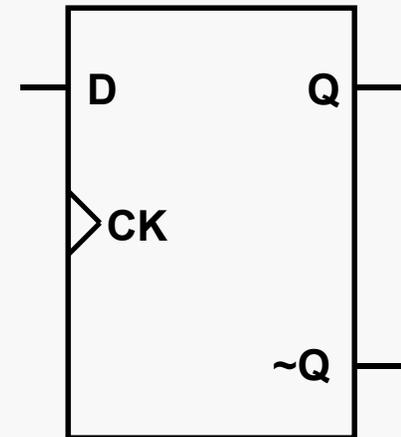
D	Q	$\sim Q$
0	0	0
0	1	0
1	0	1
1	1	1

So, we just need to feed each of the flip-flops the value of the appropriate next-state function equation derived earlier...

D Flip-flop

The D flip-flop takes one data input and updates its state Q , on a clock tick, according to the table:

D	Q	$\sim Q$
0	0	1
1	1	1

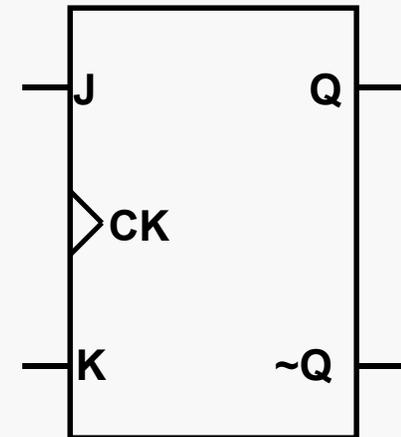


In the following Logisim diagrams, the D flip-flops update state on the falling edge (when the clock goes from high to low).

The JK flip-flop takes two data inputs and updates its state Q , on a clock tick, according to the table:

J	K	Q	$\sim Q$

0	0	no change	
0	1	0	1
1	0	1	0
1	1	opposite	

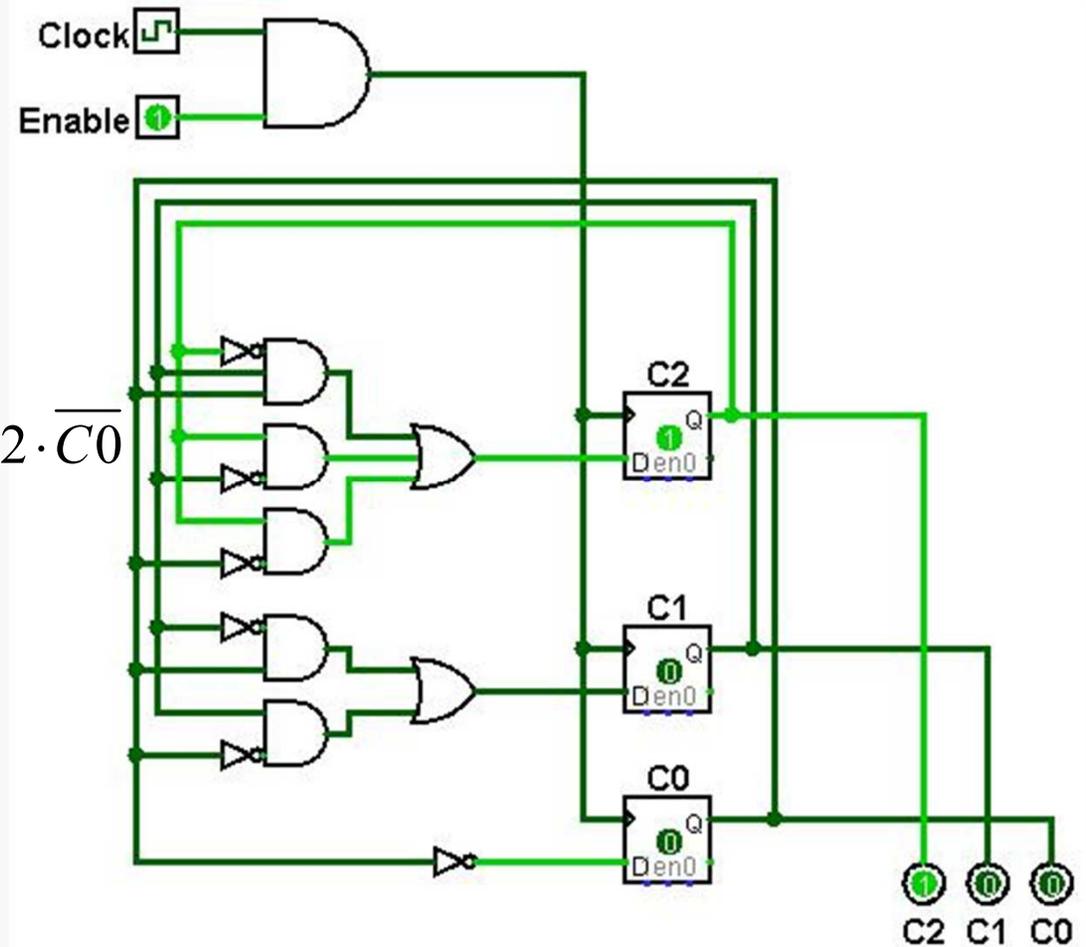


In the following Logisim diagrams, the JK flip-flops update state on the falling edge (when the clock goes from high to low).

$$N2 = \overline{C2} \cdot C1 \cdot C0 + C2 \cdot \overline{C1} + C2 \cdot \overline{C0}$$

$$N1 = \overline{C1} \cdot C0 + C1 \cdot \overline{C0}$$

$$N0 = \overline{C0}$$



We could also implement the circuit using JK flip-flops.

J	K	Q	~Q

0	0	no change	
0	1	0	1
1	0	1	0
1	1	opposite	

This will require a little more effort, since the inputs to the JK flip-flops cannot merely be set to equal the next state functions.

Design: Deriving JK Inputs

We must examine each current/next state pair and determine how/if the relevant flip-flop needs to change state:

			Flip-flop Inputs									
Current	State	Next	State	J2	K2	J1	K1	J0	K0			

0	0	0	0	0	0	0	0	1	1			
0	0	1	0	0	0	1	1	1	1			
0	1	0	0	0	0	0	0	1	1			
0	1	1	1	0	0	1	1	1	1			
1	0	0	1	0	0	0	0	1	1			
1	0	1	1	0	0	1	1	1	1			
1	1	0	1	0	0	0	0	1	1			
1	1	1	0	0	0	1	1	1	1			

For this simple circuit, we either want the JK flip-flop to hold state (both inputs 0) or to toggle state (both inputs 1).

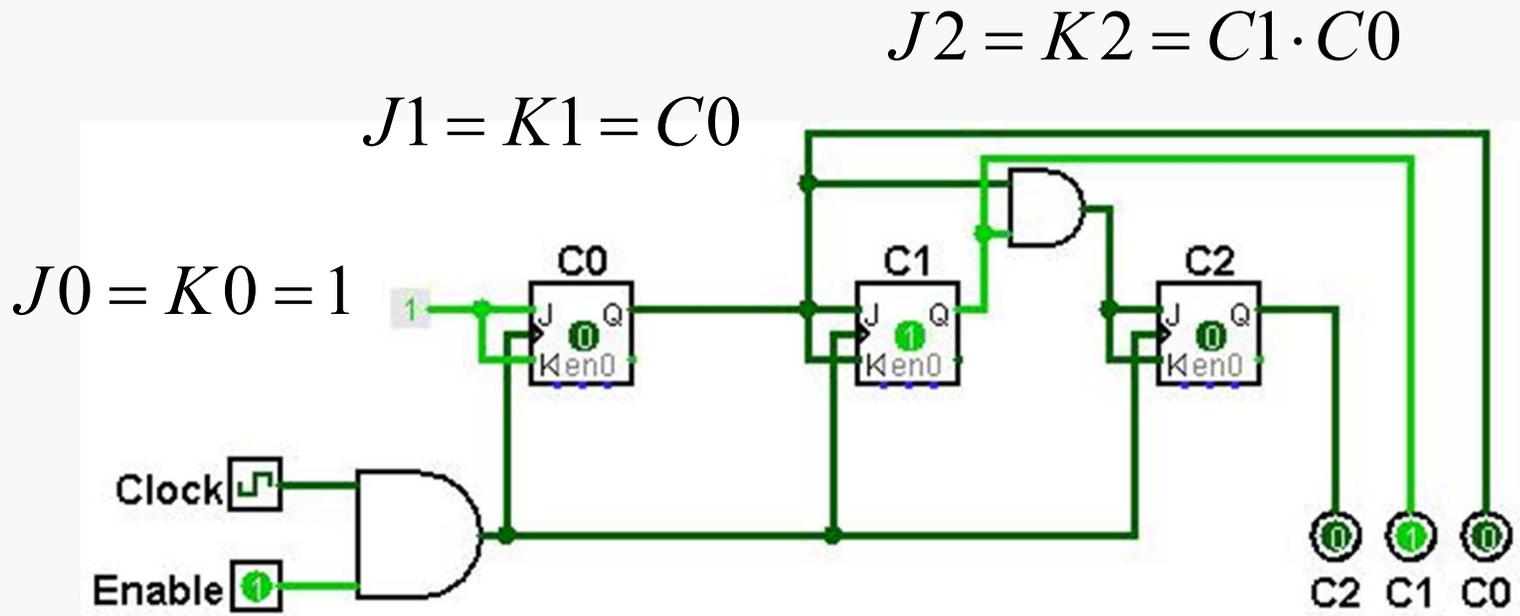
We can derive equations in the usual manner from the previous table:

Current State			Flip-flop Inputs						
C2	C1	C0	J2	K2	J1	K1	J0	K0	
0	0	0	0	0	0	0	0	1	1
0	0	1	0	0	1	1	1	1	1
0	1	0	0	0	0	0	0	1	1
0	1	1	1	1	1	1	1	1	1
1	0	0	0	0	0	0	0	1	1
1	0	1	0	0	1	1	1	1	1
1	1	0	0	0	0	0	0	1	1
1	1	1	1	1	1	1	1	1	1

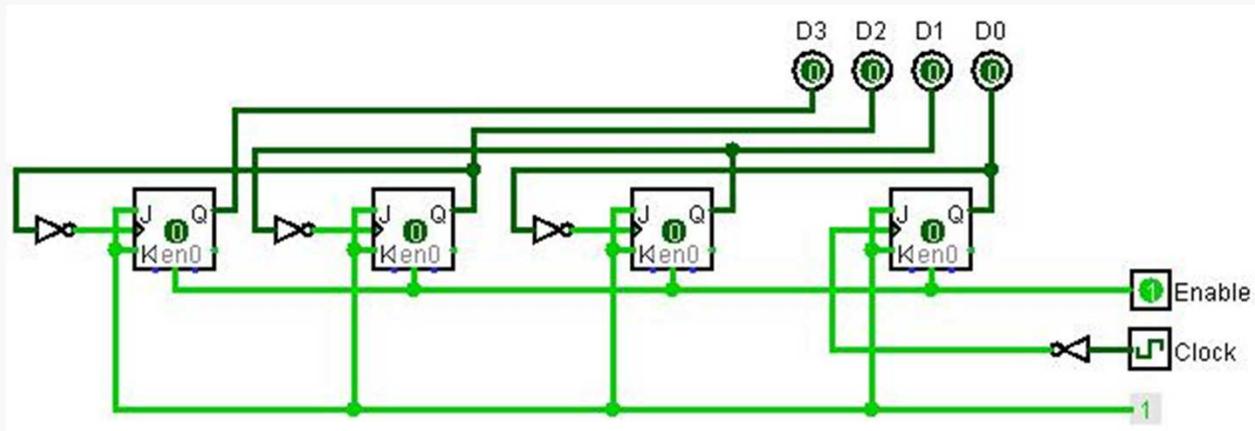
$$J0 = K0 = 1$$

$$J1 = K1 = \overline{C2} \cdot \overline{C1} \cdot C0 + \overline{C2} \cdot C1 \cdot C0 + C2 \cdot \overline{C1} \cdot C0 + C2 \cdot C1 \cdot C0 = C0$$

$$J2 = K2 = \overline{C2} \cdot C1 \cdot C0 + C2 \cdot C1 \cdot C0 = C1 \cdot C0$$



We can use JK flip-flops to implement a 4-bit counter:



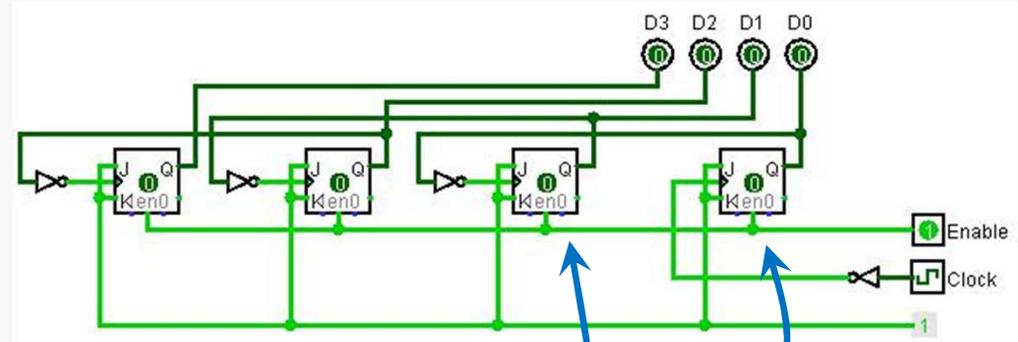
Note that the J and K inputs are all set to the fixed value 1, so the flip-flops "toggle".

As the clock signal runs, the circuit will cycle its outputs through the values 0000, 0001, 0010, . . . , 1111 and then repeat the pattern.

So, it counts clock ticks, modulo 16.

mod-16 Counter: first tick

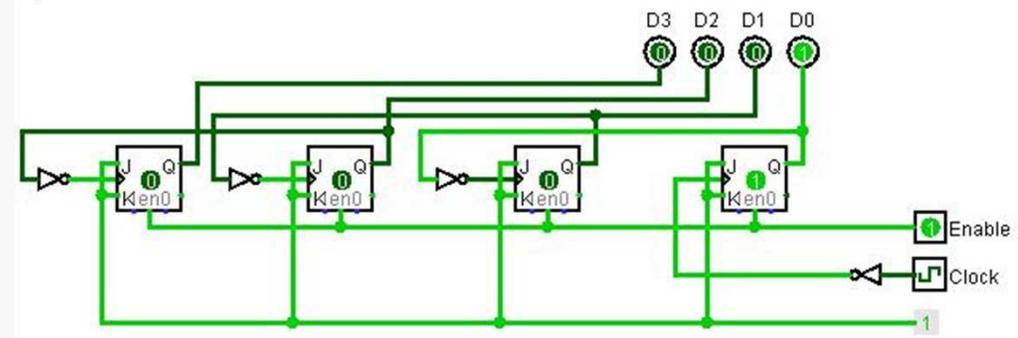
Suppose the counter is in the initial state shown below (output is 0000).



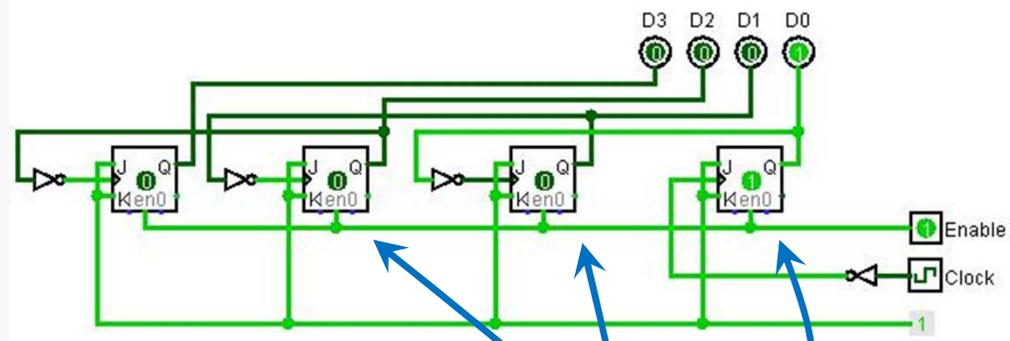
When the clock cycles from high to low:

- the right-most sees its (inverted) clock signal go from low to high, and so it toggles its state to 1
- the next flip-flop sees its clock signal go from high to low, and so it doesn't toggle
- and so, neither do the other flip-flops...

So, the output is 0001:



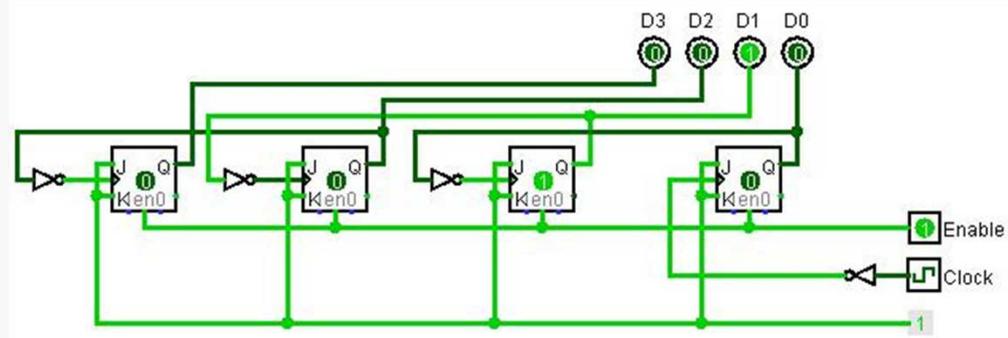
Suppose the counter is now in the state shown below (output is 0001).



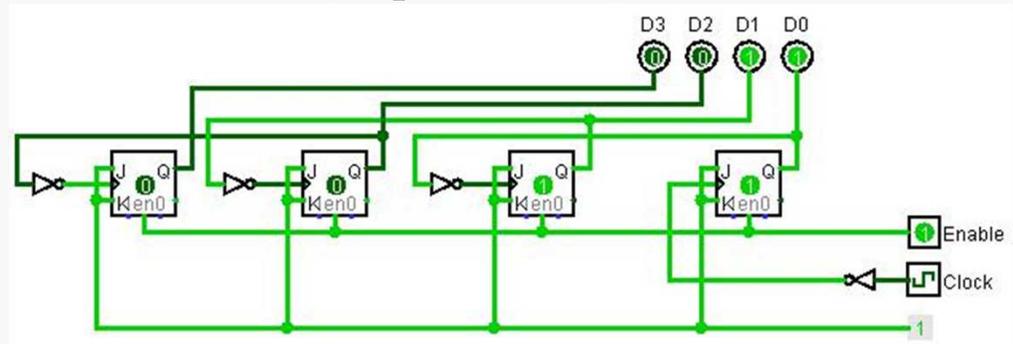
When the clock cycles from high to low (2nd cycle):

- the right-most sees its (inverted) clock signal go from low to high, and so it toggles its state to 0
- the next flip-flop sees its clock signal go from low to high, and so it toggles its state to 1
- the next flip-flop sees its clock signal go from high to low, so it doesn't toggle

So the output is 0010.



Suppose the counter is now in the state shown below (output is 0011).



When the clock cycles from high to low (4th cycle):

- the right-most sees its (inverted) clock signal go from low to high, and so it toggles its state to 0
- the next flip-flop sees its clock signal go from low to high, and so it toggles its state to 0
- the next flip-flop sees its clock signal go from low to high, and so it toggles its state to 1

So the output is 0100.

