

A *leaf procedure* is one that doesn't call any other procedures.

A *non-leaf procedure* is one that does call another procedure.

Non-leaf procedures pose an additional, but simple, challenge; we make procedure calls by executing a jump-and-link instruction:

```
jal    procedure_0    # puts PC+4 into $ra for return
```

But, if `procedure_0` also makes a call, say

```
jal    procedure_1    # puts PC+4 into $ra for return
```

then the original return address just got overwritten... the effect is fascinating...

# Preserving the Return Address

Non-leaf procedures must back up the value of their return address before making a call to another procedure:

```
addi    $sp, $sp, -4    # make room on stack
sw      $ra, 0($sp)    # save return address
```

And they must restore the return address before they attempt to return:

```
lw      $ra, 0($sp)    # retrieve return address
addi    $sp, $sp, 4    # pop it off the stack
```

Failure to do this will almost certainly lead to a catastrophic runtime failure.

The safest way to do this is to back up the address immediately when the procedure is entered, and to restore it immediately before the return is executed. Of course, you must keep careful track of the stack pointer during all of this...

# Factorial: First Version

```
#####  
# Returns factorial of parameter.  
#  
# Pre:  
#   $a0 stores N  
# Post:  
#   $v0 stores N!  
#  
# Modifies: $t0, $t1, $v0, $a0  
#  
fac1:  
    li    $t0, 1           # check for base case  
    bgt   $a0, $t0, recurse  
    li    $v0, 1           # if so, set $v0  
    jr    $ra              #         and return  
  
recurse:  
    move  $t1, $a0         # save N  
    addi  $a0, $a0, -1     # calc N-1 for recursive call  
    jal   fac1             # calc (N-1)!  
    mul   $v0, $v0, $t1    # multiply that by N  
  
    jr   $ra              #         and return
```

**Unfortunately, fac1 falls into an infinite loop when it's called with any value larger than 1 for \$a0.**

# What went wrong?

## Recursion in MIPS 4

```
fac1:
    li    $t0, 1           # check for base case
    bgt   $a0, $t0, recurse
    li    $v0, 1           # if so, set $v0
    jr    $ra              # and return

recurse:
    move  $t1, $a0         # save N
    addi  $a0, $a0, -1     # calc N-1 for recursive call
    jal   fac1             # calc (N-1)!
    mul   $v0, $v0, $t1    # multiply that by N
    jr    $ra              # and return
```

Making the recursive call overwrites the original return address with the address of *what*?

And the effect of that is....?

An infinite loop in a pgm with no loops.

And the moral of that is....?

Back up `$ra` before a call in a non-leaf proc.

# Factorial: Second Version

## Recursion in MIPS 5

```
fac2:
    li    $t0, 1           # check for base case
    bgt   $a0, $t0, recurse
    li    $v0, 1           # if so, set return value
    jr    $ra              # and return

recurse:
    move  $t1, $a0         # save N
    addi  $a0, $a0, -1     # calc N-1 for recursive call
    addi  $sp, $sp, -4     # save return address on stack
    sw    $ra, ($sp)
    jal   fac2             # calc (N-1)!
    mul   $v0, $v0, $t1    # multiply that by N

    lw    $ra, ($sp)      # restore return address
    addi  $sp, $sp, 4
    jr    $ra              # and return
```

**Unfortunately, fac2 returns  
32 when called with \$a0 == 6.**

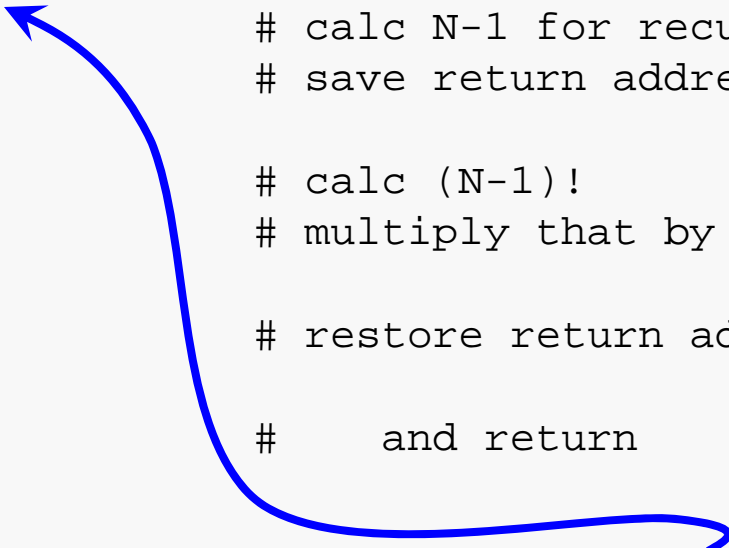
# What went wrong?

## Recursion in MIPS 6

```
fac2:
. . .

recurse:
    move    $t1, $a0           # save N
    addi   $a0, $a0, -1       # calc N-1 for recursive call
    addi   $sp, $sp, -4       # save return address on stack
    sw     $ra, ($sp)
    jal    fac2               # calc (N-1)!
    mul    $v0, $v0, $t1      # multiply that by N

    lw     $ra, ($sp)         # restore return address
    addi   $sp, $sp, 4
    jr     $ra                 # and return
```



During the recursive call, the previous contents of `$t1` and `$a0` are overwritten.

Moral: before making a call, back up your registers as necessary.

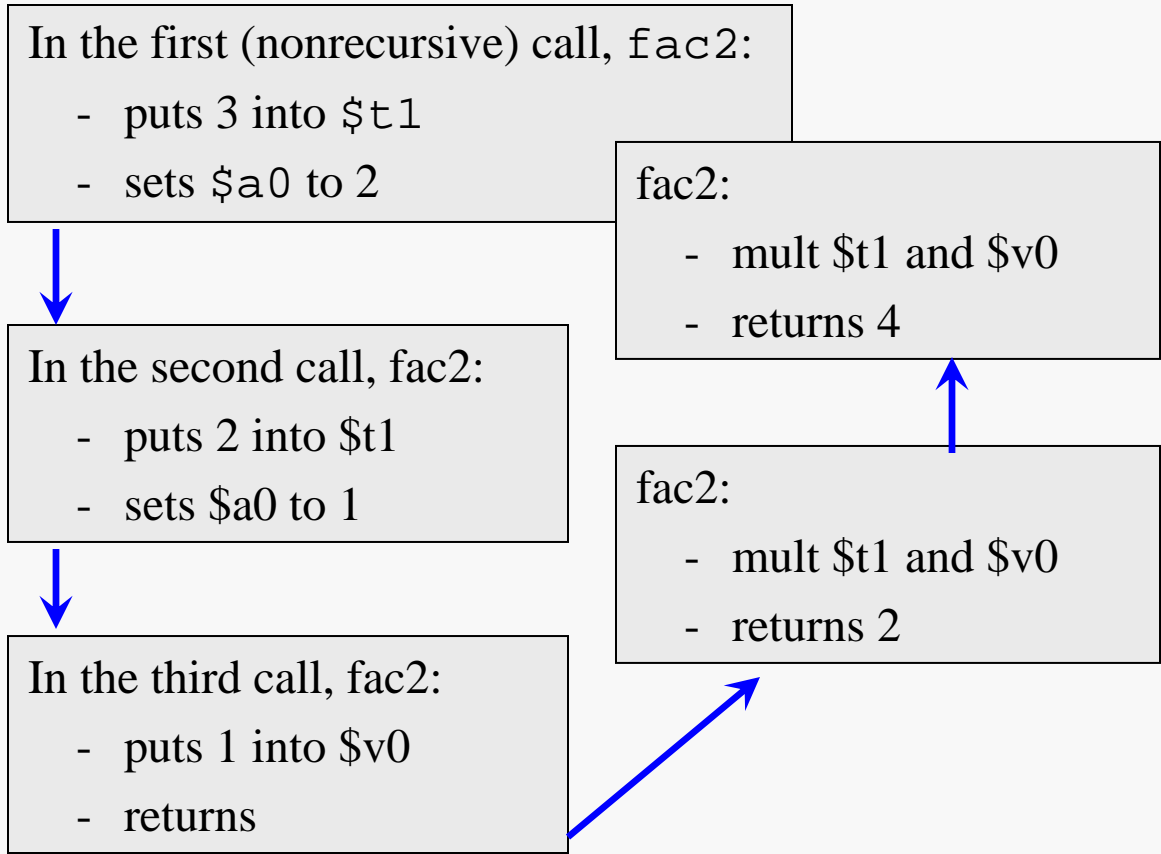
# What went wrong: Details

Let's say we call this with \$a0 set to 3:

```
fac2:
. . .

recurse:
    move    $t1, $a0
    addi   $a0, $a0, -1
    addi   $sp, $sp, -4
    sw     $ra, ($sp)
    jal    fac2
    mul    $v0, $v0, $t1

    lw     $ra, ($sp)
    addi   $sp, $sp, 4
    jr     $ra
```

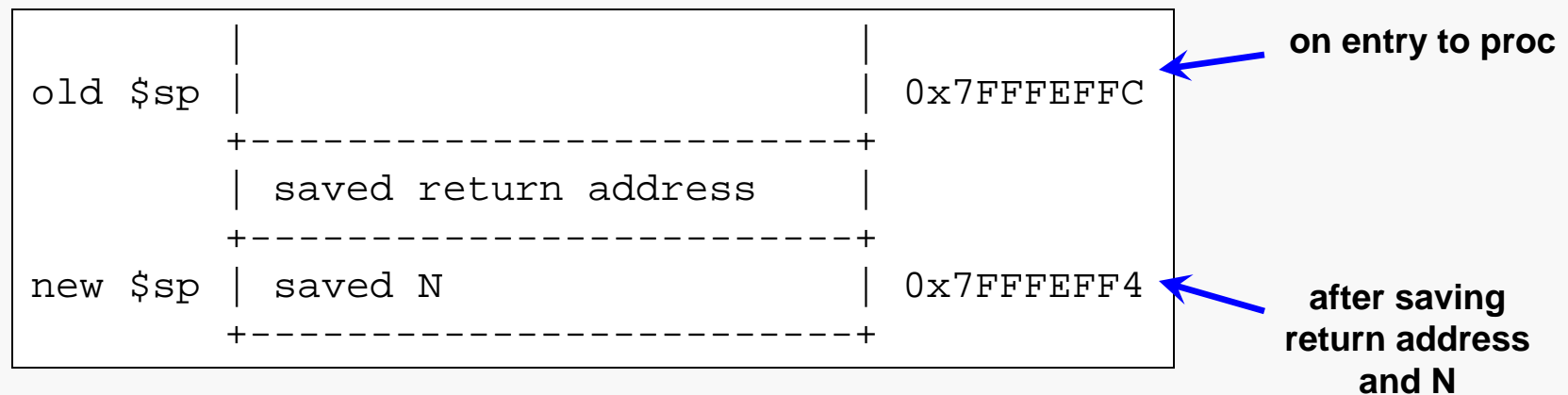


Moral: before making a call, back up your registers as necessary.

# Factorial: Stack Organization

In order to fix the execution of the recursive factorial procedure, we need to use the stack to save values that would otherwise be overwritten when a recursive call takes place.

Here's one idea for organizing the stack:





## Factorial: Third Version

## Recursion in MIPS 9

```
fac3:
    li    $t0, 1           # check for base case
    bgt   $a0, $t0, recurse
    li    $v0, 1           # if so, set return value
    jr    $ra              # and return

recurse:
    addi   $sp, $sp, -8    # make room on stack for
    sw     $ra, 4($sp)     # return address, and
    sw     $a0, 0($sp)     # N

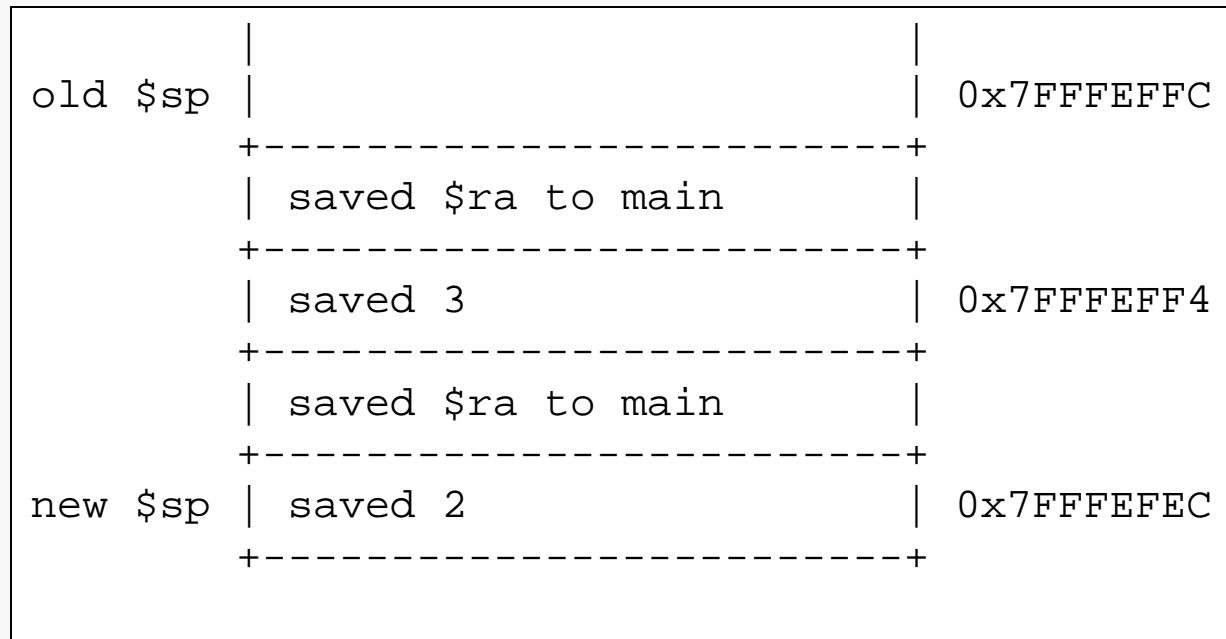
    addi   $a0, $a0, -1    # calc N-1 for recursive call
    jal    fac3            # calc (N-1)!

    lw     $t1, 0($sp)     # restore N from stack
    mul    $v0, $v0, $t1   # multiply (N-1)! by N

    lw     $ra, 4($sp)     # restore return address from
    #     stack
    addi   $sp, $sp, 8     # and restore stack pointer
    jr     $ra             # and return
```

# Factorial: Stack Trace

Say we call factorial(3):



**on first call  
(not recursive)**

**on second call  
(recursive)**

**Third call triggers base case and returns with \$v0 == 1**

**Saved value of N (2) is retrieved from stack and multiplied to \$v0; 2\*1 is returned to from second call.**

**Saved value of N (3) is retrieved from stack and multiplied to \$v0; 3\*2\*1 is returned from first call.**