

FP Foundations, Scheme (2)

In Text: Chapter 15

Functional programming

- LISP: John McCarthy 1958 MIT
 - List Processing => Symbolic Manipulation
- First functional programming language
 - Every version after the first has imperative features, but we will discuss the functional subset

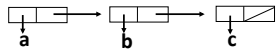
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LISP Data Types

- There are only two types of data objects in the original LISP
 - Atoms: symbols, numbers, strings,...
 - E.g., a, 100, "foo"
 - Lists: specified by delimitating elements within parentheses
 - Simple lists: elements are only atoms
 - E.g., (A B C D)
 - Nested lists: elements can be lists
 - E.g., (A (B C) D (E (F G)))

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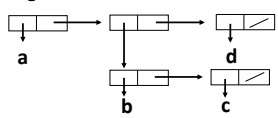
LISP Data Types

- Internally, lists are stored as **single-linked list** structures
 - Each node has two pointers: one to element, the other to next node in the list
 - Single atom:
 - atom
 - List of atoms: (a b c)
 - 

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LISP Data Types

– List containing list (a (b c) d)



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Scheme

- Scheme is a dialect of LISP, emerged from MIT in 1975
- Characteristics
 - simple syntax and semantics
 - small size
 - exclusive use of static scoping
 - treating functions as first-class entities
 - As first-class entities, Scheme functions can be the values of expressions, elements of lists, assigned to variables, and passed as parameters

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Interpreter

- Most Scheme implementations employ an interpreter that runs a "read-eval-print" loop
 - The interpreter repeatedly reads an expression from a standard input, evaluates the expression, and prints the resulting value

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Primitive Numeric Functions

- Primitive functions for the basic arithmetic operations:

+, -, *, /

- + and * can have zero or more parameters. If * is given no parameter, it returns 1; if + is given no parameter, it returns 0

- - and / can have two or more parameters

- Prefix notation

Expression	Value
42	42
(* 3 6)	18
(+ 1 2 3)	6
(sqrt 16)	4

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Numeric Predicate Functions

- Predicate functions return Boolean values (#T or #F): =, <>, >, <, >=, <=, EVEN?, ODD?, ZERO?

Expression	Value
(= 16 16)	#T
(even? 29)	#F
(> 10 (* 2 4))	
(zero? (-10(* 2 5)))	

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Type Checking

- Dynamic type checking
- Type predicate functions
 - (boolean? x) ; Is x a Boolean?
 - (char? x)
 - (string? x)
 - (symbol? x)
 - (number? x)
 - (pair? x)
 - (list? x)

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Lambda Expression

- E.g., lambda(x) (* x x) is a nameless function that returns the square of its given numeric parameter
- Such functions can be applied in the same ways as named functions
 - E.g., ((lambda(x) (* x x)) 7) = 49
- It allows us to pass function definitions as parameters

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"define"

- To bind a name to the value of a variable:
 - (define symbol expression)**
 - E.g., (define pi 3.14159)
 - E.g., (define two_pi (* 2 pi))
- To bind a function name to an expression:
 - (define (function_name parameters) expression)**
 -)**
 - E.g., (define (square x) (* x x))

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"define"

- To bind a function name to a lambda expression


```
(define function_name
  (lambda_expression)
)
```

 – E.g., (define square (lambda (x) (* x x)))

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Control Flow

- Simple conditional expressions can be written using if:
 - E.g. (if (< 2 3) 4 5) => 4
 - E.g., (if #f 2 3) => 3

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Control Flow

- It is modeled based on the evaluation control used in mathematical functions:

```
(COND
  (predicate_1 expression)
  (predicate_2 expression)
  ...
  (predicate_n expression)
  [ELSE expression]
)
```

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An Example

$$f(x) = \begin{cases} 1 & \text{if } x = 0 \\ x * f(x-1) & \text{if } x > 0 \end{cases}$$

```
(define (factorial x)
  (cond
    ((< x 0) #f)
    ((= x 0) 1)
    (#t (* x (factorial (- x 1)))) ; or else (...)
  )
)
```

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Bindings & Scopes

- Names can be bound to values by introducing a nested scope
- let takes two or more arguments:
 - The first argument is a list of pairs
 - In each pair, the first element is the name, while the second is the value/expression
 - Remaining arguments are evaluated in order
 - The value of the construct as a whole is the value of the final argument
 - E.g. (let ((a 3)) a)

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let Examples

- E.g., (let ((a 3)
 (b 4)
 (square (lambda (x) (* x x)))
 (plus +))
 (sqrt (plus (square a) (square b))))
- The scope of the bindings produced by let is its second and following arguments

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let Examples

- E.g., `(let ((a 3))
 (let ((a 4)
 (b a))
 (+ a b))) => ?`
- b takes the value of the outer a, because the defined names are visible "all at once" at the end of the declaration list

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let* Example

- let* makes sure that names become available "one at a time"
- E.g., `(let*((x 1) (y (+ x 1)))
 (+ x y)) => ?`

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Functions

- quote: identity function
 - When the function is given a parameter, it simply returns the parameter
 - E.g., `(quote A) => A`
`(quote (A B C)) => (A B C)`
- The common abbreviation of quote is apostrophe (`'`)
 - E.g., `'a => a`
`'(A B C) => (A B C)`

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List Functions

- car: returns the first element of a given list
 - E.g., `(car '(A B C)) => A`
`(car '((A B) C D)) => (A B)`
`(car 'A) => ?`
`(car '(A)) => ?`
`(car '()) => ?`

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List Functions

- cdr: returns the remainder of a given list after its car has been removed
 - E.g., `(cdr '(A B C)) => (B C)`
`(cdr '((A B) C D)) => (C D)`
`(cdr 'A) => ?`
`(cdr '(A)) => ?`
`(cdr '()) => ?`

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List Functions

- cons: concatenates an element with a list
- cons builds a list from its two arguments
 - The first can be either an atom or a list
 - The second is usually a list
 - E.g., `(cons 'A '()) => (A)`
`(cons 'A '(B C)) => (A B C)`
`(cons '() '(A B)) => ?`
`(cons '(A B) '(C D)) => ?`
 - How to compose a list (A B C) from A, B, and C?

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List Functions

- Note that cons can take two atoms as parameters, and return a dotted pair
 - E.g., (cons 'A 'B) => (A . B)
 - The dotted pair indicates that this cell contains two atoms, instead of an atom + a pointer or a pointer + a pointer

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More Predicate Functions

- The following returns #t if the symbolic atom is of the indicated type, and #f otherwise
 - E.g., (symbol? 'a) => #t
(symbol? '()) => #f
 - E.g., (number? '55) => #t
(number? 55) => #t
(number? 'a) => #f
 - E.g., (list? '(a)) => #t
 - E.g., (null? '()) => #t

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More Predicate Functions

- eq? returns true if two objects are equal through pointer comparison
 - Guaranteed to work on symbols
 - E.g., (eq? 'A 'A) => #T
(eq? 'A '(A B)) => #F
- equal? recursively compares two objects to determine if they are equal
 - The objects can be symbols, atoms, numbers, and lists

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How do we implement equal?

```
(define (atom? atm)
  (cond
    ((list? atm) (null? atm))
    (else #T)
  )
)

(define (equal? lis1 lis2)
  (cond
    ((atom? lis1) (eq? lis1 lis2))
    ((atom? lis2) #F)
    ((equal? (car lis1) (car lis2))
     (equal? (cdr lis1) (cdr lis2)))
    (else #F)
  )
)
```

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More Examples

```
(define (member? atm lis)
  (cond
    ((null? lis) #F)
    ((eq? atm (car lis)) #T)
    (else (member? atm (cdr lis)))
  )
)

(define (append lis1 lis2)
  (cond
    ((null? lis1) lis2)
    (else (cons (car lis1)
                 (append (cdr lis1) lis2)))
  )
)
```

What is returned for the following function?
(member? 'b '(a (b c)))

Is lis2 appended to lis1, or lis1 prepended to lis2?

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An example: apply-to-all function

```
(define (mapcar fctn lis)
  (cond
    ((null? lis) '())
    (else (cons (fctn (car lis))
                 (mapcar fctn (cdr lis)) ))
  )
)
```

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